

Logic Programming

Computational Model

Temur Kutsia

Research Institute for Symbolic Computation
Johannes Kepler University of Linz, Austria
`kutsia@risc.jku.at`

Contents

Preliminaries

Abstract Interpreter

Choice Points

Basic Notions

Term: Constant, variable, or compound term.

Compound Term: Functor, arguments

$$f(t_1, \dots, t_n)$$

Functor: Name, arity

$$f/n$$

Goal: Atom or compound term.

Logic Programs

Clause: Universally quantified logical sentence

$$A \leftarrow B_1, \dots, B_k, k \geq 0$$

A and B_i 's are goals.

Declarative reading: A is implied by the conjunction of the B_i 's.

Procedural reading: To answer the query A , answer the conjunctive query B_1, \dots, B_k .

Logic Program: Finite set of clauses.

Computation

Query: Existentially quantified conjunction

$\leftarrow A_1, \dots, A_n, n > 0$

A_i 's are goals.

Computation of a Logic Program P : finds an instance of a given query logically deducible from P .

How to Compute

- ▶ Start from initial query G .
- ▶ Computation terminates – success or failure.
- ▶ Computation does not terminate – no result.
- ▶ Output of a successful computation: the instance of G proved.
- ▶ A given query can have several successful computations with different output.

Abstract Interpreter

INPUT:

A logic program P and a query G .

OUTPUT:

$G\theta$, if this was the instance of G deduced from P , or *failure* if failure has occurred.

Abstract Interpreter

ALGORITHM:

Let *resolvent* be G

While *resolvent* is not empty **do**

1. **Choose** a goal A from *resolvent*.
2. **Choose** a renamed clause $A' \leftarrow B_1, \dots, B_n$ from P such that A and A' unify with an mgu θ (**exit** if no such goal and clause exist).
3. Remove A from and **add** B_1, \dots, B_n to *resolvent*.
4. Apply θ to *resolvent* and to G .

If *resolvent* is empty, **return** G , else **return** *failure*.

Choosing and Adding

Choosing and Adding:

- ▶ Left unspecified in the abstract interpreter.
- ▶ Must be resolved in a realization of the computational model.

Two Choices

Completely different nature.

Choice of a goal:

- ▶ Arbitrary.
- ▶ Does not affect computation.
- ▶ If there exists a successful computation by choosing one goal, then there is a successful computation by choosing any other goal.

Choice of a clause:

- ▶ Non-deterministic.
- ▶ Affects computation.
- ▶ Choosing one clause might lead to success, while choosing some other might lead to failure.

Adding Goal to Resolvent

Assume: Always the leftmost goal to be chosen

Then: Adding new goal to the beginning of the resolvent gives depth-first search.

Adding new goal to the end of the resolvent gives breadth-first search.

Prolog's Solution

- ▶ Choice of a goal: leftmost.
- ▶ Choice of a clause: Topmost.
- ▶ Adding new goal to the resolvent: At the beginning.