

A Bridge between Euclid and Buchberger

(An attempt to enhance Gröbner basis algorithm
by PRSs and GCDs
(Poly.Rem.Seq.) (Great.Com.Div.))

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Outline of Talk

- 1) Variable Elimination : **PRS-method** vs **GB-method**
(lexico.-order Gröbner Basis)
- 2) **2**-polynomial system : PRS-method \Rightarrow **lowest**($\langle G, H \rangle$)
- 3) **n**-polynomial system : **healthy** system \Rightarrow **Theorem 2**
- 4) **rectangular PRSs** \Rightarrow **extraneous factor** removal
- 5) Elimination of **Lead.Coeff**-vars \Rightarrow **LCtoW** polynomial

By Euclid, we mean following Two

- **Euclidean** Method for $G, H \in \mathbb{K}[x, \mathbf{u} = u_1, \dots, u_n]$
 $P_1 := G, P_2 := H$, where $\deg_x(G) \geq \deg_x(H) \geq 1$,
 $P_{i+1} := (\alpha_i P_{i-1} - Q_i P_i) / \beta_i, \alpha_i, \beta_i \in \mathbb{K}[\mathbf{u}]$,
 $\deg_x(P_{i+1}) < \deg_x(P_i), P_k \neq 0, P_{k+1} = 0$,
 α_i, β_i are so chosen that $P_{i+1} \in \mathbb{K}[x, \mathbf{u}]$.
- **Extended** Euclidean algorithm for $A_i, B_i \in \mathbb{K}[x, \mathbf{u}]$,
satisfying $A_i G + B_i H = P_i, (i = 2, 3, \dots, k-1)$,
 $A_{i+1} := (\alpha_i A_{i-1} - Q_i A_i) / \beta_i, A_1 = 1, A_2 = 0$,
 $B_{i+1} := (\alpha_i B_{i-1} - Q_i B_i) / \beta_i, B_1 = 0, B_2 = 1$.
 (A_i, B_i) is uniquely determined if we fix P_i .

History of Variable Elimination

(Sasaki's personal view)

(Given $\{F_1, \dots, F_{m+1}\} \in \mathbb{K}[\mathbf{x} = x_1, \dots, x_m, \mathbf{u}]$,
eliminate $\mathbf{x}_1, \dots, \mathbf{x}_m$ of $\{F_1, \dots, F_{m+1}\}$, if possible)

- Takakazu Seki (Japan) : multivariate **resultant, determinant** (with Tanaka et al.) (1674~1685)
discriminant (with Tatebe) (~1685)
- I. Newton
L. Euler : elimination method for 2-polynomial system (Newton:1707, Euler:1748)
- E. Bézout : variable elimination method (1764)
similar to that by Seki et al.
- J.J. Sylvester : determinant for uni-var elimi. (1840)
- F.S. Macaulay
A.L. Dixon
et al. : determinants for m -var elimination (Macaulay:1902,16, Dixon:1908)
monomials in \mathbf{x} , polynoms. in \mathbf{u}
encounter **extraneous factors**
- B. Buchberger (1965) : theory & algorithm of Gröbner basis
... monomials in both \mathbf{x} and \mathbf{u}
♠ gives **lowest-order resultant**
- J.E. Collins,
W.S. Brown : subresultant PRS algorithm: elimi.
main var. (Collins:1967, Brown:1978)
& extended PRS: $A_k G + B_k H = P_k$
- D. Lazard ('83) : matrix: all possible monos in column
apply Gaussian elimination to it
- D. Kapur et al. (1990s) : revival of **sparse resultants**
of Mcaulay, Dixon, et al.
extraneous factors still remain

Lead.-Mono. vs Lead.-Term Eliminations

(Ref. Knuth-Bendix (1967))

GB : ^(Monomial) **Mono Representat.** & **Lead.-mono Elimination**

$$F(x) = c_1 M_1(x) + \cdots + c_m M_m(x)$$

$$M(x) = x_1^{e_1} \cdots x_n^{e_n}, \quad M_1 \succ \cdots \succ M_m$$

$$\text{Spol}(F, F') = (c'_1 M'_1 / C) F - (c_1 M_1 / C) F'$$

$$\text{where } C = \text{gcd}(c_1 M_1, c'_1 M'_1)$$

PRS : ^(Recursive) **Recu. Representat.** & **Lead.-term Elimination**

$$F(X, u) := f_d(u) X^d + \cdots + f_0(u) X^0$$

$$\text{Elim}(F, F') := (f'_{d'} / \gamma) F - X^{d-d'} (f_d / \gamma) F'$$

$$\text{where } d \geq d', \quad \gamma = \text{gcd}(f_d(u), f'_{d'}(u))$$

Coefficients of Generators (^{new} name)

PRS: $P_k := \text{lastPRS}_x(G, H) = \mathbf{A}_k G + \mathbf{B}_k H :$

$$P_k \in \mathbb{K}[u] \Rightarrow \begin{cases} \deg_x(\mathbf{A}_k) < \deg_x(H) \\ \deg_x(\mathbf{B}_k) < \deg_x(G) \end{cases}$$

GB: $\hat{S} := \text{lowest}$ element of $\text{GB}(\{G, H\})$

$$= \mathbf{A} G + \mathbf{B} H :$$

$$\hat{S} \in \mathbb{K}[u] \Rightarrow \begin{cases} \deg_x(\tilde{\mathbf{A}}) \not< \deg_x(H) \\ \deg_x(\tilde{\mathbf{B}}) \not< \deg_x(G) \end{cases}$$

GB: in general, for $\text{GB}(\{F_1, \dots, F_{m+1}\}) :$

$$\hat{S} = \tilde{\mathbf{A}}_1 F_1 + \cdots + \tilde{\mathbf{A}}_{m+1} F_{m+1}$$

2-Pol. System : Compare GB vs PRS

(GBmethod(Mathematica) vs PRSmethod(GAL))
 (data by **Inaba**) (in **Sasaki** Lab.)

$$\text{Ex2017: } \begin{cases} G = X^6(u+2v+w) + X^4(u-2x-z) \\ \quad + X^2(v+3y-z) + (v+2w+y), \\ H = X^6(v-w+2x) - X^4(v+y-2z) \\ \quad + X^2(w-2x+y) + (u-v+2z). \end{cases}$$

Ex-6: $(G_6, H_6) := (G, H)$,

Ex-5: $(G_5, H_5) :=$ replace (z) by (w) in (G, H)

Ex-4: $(G_4, H_4) :=$ replace (y, z) by (v, w) in (G, H)

Ex-3: $(G_3, H_3) :=$ replace (x, y, z) by (u, v, w) in (G, H)

♠ **GB vs PRS : Lowest**($\langle G, H \rangle$) $\Leftrightarrow P_k$

(Table from S&Inaba (2017))

	GB (lex) time (msec)	sparsePRS with A'_k & B'_k			
		M-time	G-time	$\#(P'_k)$	$\#(P_k)$
Ex-3	46.33	78.0	5.27	65	28
Ex-4	12040.	218.	18.64	279	81
Ex-5	>90 min	749.	65.47	961	201
Ex-6	>90 min	2246.	224.8	2815	445

GB(lex) : reduced Gröbner Basis (lex. term-order)

M-time : \Leftarrow programed in Mathematica language

G-time : \Leftarrow programed in LISP language of GAL

$\#(P'_k, P_k)$: **#mono**(**with, without**) extran.-factor

A Relation between Two Eliminations

(assume $\text{ltm}(H) \not\parallel \text{ltm}(G)$)

Lemma 1

Let $\deg(G) \geq \deg(H) \geq 1$. Let E_1 be $\underline{\text{LtmElim}(G, H)}$.
 Let \widehat{E}_1 be the **lowest** polynom., obtained by decreasing **degree** of G to $\deg(E_1)$ by leading-monomial eliminatn, where only $\text{ltm}(G)$ & $\text{ltm}(H)$ are used in elimination.
 Then, E_1 is a **constant multiple** of \widehat{E}_1 .

Proof Both are lowest-order polynomials and unique. \square

We show a simple example

$$\begin{cases} G = \underline{x^4 \cdot (y+u)} + x^2 \cdot (y-2w) + (2u+w), \\ H = \underline{x^4 \cdot (y-w)} + x^2 \cdot (2y+u) + (u-2w). \end{cases}$$

LtmElim(G, H) gives E_1 as follows :

$$\text{lcf}(G) = y+u, \quad \text{lcf}(H) = y-w, \quad \gamma = \text{gcd}(y+u, y-w) = 1,$$

$$\text{LtmElim}(G, H) = \underline{(y-w)} \times G - \underline{(y+u)} \times H \implies E_1 :=$$

$$\underline{(y-w)}[x^2(y-2w) + (2u+w)] - \underline{(y+u)}[x^2(2y+u) + (u-2w)]$$

Leading-mono eliminations give \widehat{E}_1 as follows :

(we put $R_G := \text{rest}(G)$ and $R_H := \text{rest}(H)$)

$$G = \underline{x^4 y} + \underline{x^4 u} + R_G, \quad H = \underline{x^4 y} - \underline{x^4 w} + R_H,$$

$$\text{Spol}(G, H) = G - H = \underline{x^4 u} + \underline{x^4 w} + R_G - R_H \implies G_3,$$

$$\text{Spol}(G, G_3) = -\underline{x^4 y w} + \underline{x^4 u^2} - (y-u) R_G + (y+w) R_H$$

$$\xrightarrow{H} \underline{x^4 u^2} - \underline{x^4 w^2} - (y-u) R_G + (y+w) R_H$$

$$\xrightarrow{G_3} - (y-w) R_G + (y+u) R_H \implies \widehat{E}_1,$$

$$\text{Spol}(H, G_3) = \dots \xrightarrow{G} - (y-w) R_G + (y+u) R_H = \widehat{E}_1.$$

PRS-method Computes $\text{lowest}(\langle G, H \rangle)$ without Computing any S-polynomial

Theorem 1 (S&I 2017)

Let $G, H \in \mathbb{K}[X, u]$ be relatively prime, $P_k \in \mathbb{K}[u]$ be the last element of $\text{PRS}(G, H)$, $A_k, B_k \in \mathbb{K}[X, u]$ satisfy $A_k G + B_k H = P_k$ & **degree conditions** $\deg(A_k) < \deg(H)$, $\deg(B_k) < \deg(G)$. Then, we have $P_k / \text{gcd}(\text{cont}_X(A_k), \text{cont}_X(B_k)) = c \hat{S}$, $c \in \mathbb{K}$, where, \hat{S} is the lowest element of $\text{GB}(\{G, H\})$.

*) $\text{cont}(F) = \text{gcd}(f_d, \dots, f_0)$ for $F = f_d X^d + \dots + f_0$

Outline of Proof

- 1) Let $\tilde{A}G + \tilde{B}H = \hat{S} \iff$ Buchberger's method,
 $\deg(\tilde{A}) > \deg(H)$, $\deg(\tilde{B}) > \deg(G)$, in general.
- 2) Show that we **can decrease** $\deg(\tilde{A})$ and $\deg(\tilde{B})$.
Easy when $\gamma := \text{gcd}(\text{lcf}(G), \text{lcf}(H)) = 1$
 \Rightarrow **next screen** ($\text{lcf}(F)$: leading-coefficient)
- 3) else Show that factors of γ move to \tilde{A}, \tilde{B} as
 x_1 -elimination proceeds \Rightarrow **2-next screen**
(Lemma 1 \Rightarrow we treat A_i, B_i ($i \leq k$) instead of \tilde{A}_i, \tilde{B}_i)

Detail of Proof : Case of $\gamma = 1$

(for $\tilde{A}_{k+j}, \tilde{B}_{k+j}$ ($j \geq 1$))

Assuming $\deg(\tilde{A}G) = \deg(\tilde{B}H) \geq \deg(GH)$, consider **ltm** (= leading-term) of l.h.s. of $(*) \tilde{A}G + \tilde{B}H = \hat{S}$.
 $\gamma = 1 \Rightarrow q_A := \text{ltm}(\tilde{A})/\text{ltm}(H), q_B := \text{ltm}(\tilde{B})/\text{ltm}(G)$
 are polynoms. Put $\tilde{A} = q_A H + \tilde{A}', \tilde{B} = q_B G + \tilde{B}'$,
 where $\tilde{A}' = \text{rest}(\tilde{A}) - q_A \text{rest}(H)$ & $\tilde{B}' = (\dots)$, we see
 $q_A + q_B = 0, \deg(\tilde{A}') < \deg(\tilde{A}), \deg(\tilde{B}') < \deg(\tilde{B})$.
 Substituting these into $(*)$, we get $\tilde{A}'G + \tilde{B}'H = \hat{S}$.
 Repeating this, we attain the proof. \square

Detail of Proof : Case of $\gamma \neq 1$

(for A_i, B_i ($i \leq k$) rare detail is omitted)

Consider the formulas on **PRS** and related A_i (& B_i) :

$$P_{i+1} := (c_i/\gamma_i)P_{i-1} - (c_{i-1}/\gamma_i)X^{d_i}P_i, \quad i = 2, 3, \dots$$

$$A_{i+1} := (c_i/\gamma_i)A_{i-1} - \underline{(c_{i-1}/\gamma_i)X^{d_i}A_i}, \quad A_1 = 1, A_2 = 0$$

$$\gamma_i = \text{gcd}(c_{i-1}, c_i), \quad c_i = \text{lcf}(P_i), \quad d_i = \deg(P_{i-1}) - \deg(P_i)$$

Let $\hat{\gamma}$ be a factor of γ , and consider that $\hat{\gamma}$ is **contained**
 in c_{i-1} but **not** in $c_i \Rightarrow (c_{i-1}/\gamma_i)$ **contains** $\hat{\gamma}$.

This means that $\hat{\gamma}$ is **moved** to the leading-term of A_{i+1} ,
 because $\text{ltm}(A_{i+1}) = -\text{ltm}((c_{i-1}/\gamma_i)X^{d_i}A_i)$.

Since $\hat{\gamma} \rightarrow 1$ as $i \rightarrow k$, we attain the proof. \square

Main Target : Many-Polynom. System

$$\mathcal{F} := \{F_1(\mathbf{x}, \mathbf{u}), \dots, F_{m+1}(\mathbf{x}, \mathbf{u})\}, \quad m \geq 2$$

$$(\mathbf{x}) = (x_1, \dots, x_m), \quad (\mathbf{u}) = (u_1, \dots, u_n)$$

$$x_1 \succ \dots \succ x_m \quad \succ \quad u_1 \succ \dots \succ u_n$$

Coefficients of Generators (CofG in short)

$$A_1, \dots, A_{m+1} \in \mathbb{K}[\mathbf{x}, \mathbf{u}], \text{ satisfying,}$$

$$A_1 F_1 + \dots + A_{m+1} F_{m+1} = R \in \mathbb{K}[\mathbf{u}]$$

Coefficients of Generators in \mathbf{u} (CofGu)

$$(a_1, \dots, a_{m+1}) = (A_1, \dots, A_{m+1})|_{\mathbf{x} = \mathbf{0}}$$

Many-Pol. Systems are Complicated

- **ALL** variables (\mathbf{x} & \mathbf{u}) are eliminated
if $F_i = F_j + 1$ for some $i \neq j$
- **NONE** of x_1, \dots, x_m is eliminated
if $F_i = G(\mathbf{x})F'_i$ for $\forall i$
- At least one of x_1, \dots, x_m is **NOT** eliminated
if $F_i = aF_j + bF_k$ ($i \neq j \neq k$)
- and so on

We want to treat these systems **simply**
Pathological systems \Rightarrow **exceptional** cases.

Check Extran. Factor by Example

(we will use this **EXAMPLE** often)

$$\mathcal{F}_{2018} = \begin{cases} F_1 = x^4 \cdot (y+u) + x^2 \cdot (y-2w) + (2u+w), \\ F_2 = x^4 \cdot (yu) + x^2 \cdot (y+2w) + (3u-w), \\ F_3 = x^4 \cdot (y-u) + x^2 \cdot (2y+u) + (u-2w). \end{cases}$$

$$\begin{aligned} \widehat{S} = & 33 u^7 + 23 u^6 w - 126 u^6 - 55 u^5 w^2 - 343 u^5 w + 316 u^5 - 12 u^4 w^3 \\ & - 130 u^4 w^2 + 544 u^4 w - 202 u^4 + 32 u^3 w^4 + 218 u^3 w^3 + 548 u^3 w^2 \\ & - 128 u^3 w - 144 u^2 w^4 + 428 u^2 w^3 - 420 u^2 w^2 + 144 u w^4 - 256 u w^3 \\ & - 32 w^4. \end{aligned}$$

Is **Theorem 1** **EFFECTIVE** for \mathcal{F} ?

: **Wow, Extraneous Factor is Big!**

$$\begin{aligned} G_2 &:= \mathbf{res}_x(F_1, F_2), \quad G_3 := \mathbf{res}_x(F_1, F_3) \quad (\Leftarrow \text{eliminate } x) \\ &\Rightarrow H_3 := \mathbf{res}_y(G_2, G_3) \quad (\Leftarrow \text{eliminate } y) \end{aligned}$$

$$H_3 = \widehat{S} \times u^2 \times E_3, \quad \text{where}$$

$$\begin{aligned} E_3 = & 704 u^{12} + 1664 u^{11} w - 3568 u^{11} + 720 u^{10} w^2 - 2624 u^{10} w + 6932 u^{10} - 1136 u^9 w^3 \\ & + 16200 u^9 w^2 - 8 u^9 w - 6579 u^9 - 1084 u^8 w^4 + 22504 u^8 w^3 - 39387 u^8 w^2 \\ & - 12208 u^8 w + 192 u^7 w^5 - 983 u^7 w^4 - 11531 u^7 w^3 - 6351 u^7 w^2 + 667 u^6 w^6 \\ & - 12854 u^6 w^5 + 77287 u^6 w^4 - 28467 u^6 w^3 + 365 u^5 w^7 - 2337 u^5 w^6 + 58336 u^5 w^5 \\ & - 49039 u^5 w^4 + 87 u^4 w^8 + 4225 u^4 w^7 - 7134 u^4 w^6 - 22022 u^4 w^5 + 8 u^3 w^9 \\ & + 2267 u^3 w^8 - 1286 u^3 w^7 - 8044 u^3 w^6 + 336 u^2 w^9 + 10982 u^2 w^8 + 8882 u^2 w^7 \\ & + 3576 u w^9 + 23744 u w^8 + 6448 w^9. \end{aligned}$$

extraneous factor is $u^2 \times E_3$



Introduction of Healthy System

System \mathcal{F} is **Healthy** if

- 1) All the $\mathbf{x}_1, \dots, \mathbf{x}_m$ can be **eliminated**
- 2) None of $\mathbf{u}_1, \dots, \mathbf{u}_n$ can be **eliminated**
- 3) Such cases do **NOT occur** that

$$\text{GB}(\mathcal{F}) \cap \mathbb{K}[\mathbf{u}] = \{G_1, \dots, G_{l \geq 2}\},$$
 satisfying

$$\text{LMvars}(G_i) \cap \text{LMvars}(G_j) = \emptyset \text{ for } \forall(i \neq j);$$
 ($\text{LMvars}(G)$ = all variables in Lead-Monomial of G)
 ($\Leftrightarrow \mathbf{u}_1, \dots, \mathbf{u}_n$ are **distributed** into G_1, \dots, G_l)

Main Theorem for Many-Pol. Systems

Theorem 2 (S&I 2018)

If \mathcal{F} is healthy then $\text{GB}(\mathcal{F}) \cap \mathbb{K}[\mathbf{u}] = \{\hat{S}\}$

Outline of Proof

Suppose $\text{GB}(\mathcal{F}) \cap \mathbb{K}[\mathbf{u}] = \{S_1, \dots, S_{l \geq 2}\}, S_1 \prec \dots \prec S_l$.
 First, treat the case that each S_i satisfies Condition 2).
 Then, $\text{Spol}(S_1, S_2)$ is not zero, and of lower order than S_2 , contradicting the **reducedness** of $\text{GB}(\mathcal{F})$.

$\mathbf{u}_1, \dots, \mathbf{u}_n$ may be **distributed** among S_1, \dots, S_l .
 This case is not healthy by **Condition 3)**.



Introduction of RectAngular PRSs (rectPRSs, in short)

Triangular PRSs (conventional)

$$\begin{aligned}
 G_i &:= \text{lastPRS}_{\mathbf{x}_1}(F_1, F_i), & \dots, & \dots & (i \geq 2) \\
 G'_i &:= \text{lastPRS}_{\mathbf{x}_2}(G_2, G_i), & \dots & & (i \geq 3) \\
 & \vdots & & & \\
 G'''_{m+1} &:= \text{lastPRS}_{\mathbf{x}_m}(G''_m, G'''_{m+1})
 \end{aligned}$$

Rectangular PRSs (our method)

$$\begin{aligned}
 G_1 &:= \text{lastPRS}_{\mathbf{x}_1}(F_1, F_2), & \dots & G_{m+1} &:= \text{lastPRS}_{\mathbf{x}_1}(F_{m+1}, F_1) \\
 G'_1 &:= \text{lastPRS}_{\mathbf{x}_2}(G_1, G_2), & \dots, & G'_{m+1} &:= \text{lastPRS}_{\mathbf{x}_2}(G_{m+1}, G_1) \\
 & \vdots & & \vdots & \\
 G'''_1 &:= \text{lastPRS}_{\mathbf{x}_m}(G''_1, G'''_2), & \dots, & G'''_{m+1} &:= \text{lastPRS}_{\mathbf{x}_m}(G''_{m+1}, G'''_1)
 \end{aligned}$$



Remove Extrn.Factr by RectPRSs

(Eliminate $\mathbf{x}, \mathbf{y} \Rightarrow$ RectAngular PRSs)

$$(F_1, F_2, F_3) \Rightarrow (G_1, G_2, G_3) \Rightarrow (H_1, H_2, H_3)$$

Theorem 2 \Rightarrow Each H_i is a multiple of \hat{S}

$\text{gcd}(H_1, H_2, H_3)$ will be a small multiple of \hat{S}



$$\begin{aligned}
 H_1 &= 382239 u^{22} - 313632 u^{21}w - 3218292 u^{21} - 172611 u^{20}w^2 + \dots, \\
 H_2 &= 363 u^{21} - 4334 u^{20}w - 14190 u^{20} + 20453 u^{19}w^2 + \dots, \\
 H_3 &= -23232 u^{21} - 71104 u^{20}w + 206448 u^{20} - 23312 u^{19}w^2 + \dots
 \end{aligned}$$

$$\overline{H} := \text{gcd}(H_1, H_2, H_3) = u^2 \hat{S}$$

We want to remove u^2 further

$$(f_1^{(0)}, \dots, f_{m+1}^{(0)}) := (F_1(\mathbf{0}, \mathbf{u}), \dots, F_{m+1}(\mathbf{0}, \mathbf{u}))$$

(if some $f_i^{(0)} = 0$ then move **origin** of \mathbf{u})

$$(a_1, \dots, a_{m+1}) : \text{CofGs of } H_i \text{ in } \mathbf{u}$$

Proposition 1 (S&I 2018)

- 1) If $\bar{f} := \gcd(f_1^{(0)}, \dots, f_{m+1}^{(0)}) \notin \mathbb{K}$ then $\hat{S} = \text{lowest}(\text{GB}(\mathcal{F}))$ has \bar{f} as a factor.
- 2) If $\bar{a} := \gcd(a_1, \dots, a_{m+1}) \notin \mathbb{K}$ then \bar{a} is an **extraneous factor** of H_i .

Hint for the Proof

Consider $\bar{H}_i = a_1 F_1 + \dots + a_{m+1} F_{m+1} (\in \langle \mathcal{F} \rangle)$

Proposition 1 removes u^2 extraneous factor :

Prop. 1 tells that H_1, H_2, H_3 contain u^2, u^1, u^1 extran. factors, respectively.

while each H_i is not divisible by u^3 .

$$\Rightarrow \bar{H}/u^2 \text{ is } \underline{\text{irreducible}} \Rightarrow \bar{H}/u^2 = \hat{S}$$

HOW to Enhance Buchberger's Method

Let $\mathbf{GB}(\mathcal{F}) = \{\widehat{G}_1, \widehat{G}_2, \dots\}$, where $\widehat{G}_1 \prec \widehat{G}_2 \prec \dots$

Let \widetilde{G}_i be either a small multiple or LM-multiple of \widehat{G}_i
 (LM-multiple : $\text{lcm}(\widetilde{G}_i)$ is a multiple of $\text{lcm}(\widehat{G}_i)$)

==== **Plan** ====

- 1) Eliminate variables $x_1, \dots, x_m \implies$ obtain **rectPRSs**
 (Each element of rectPRSs $\in \langle \mathcal{F} \rangle$)
- 2) Remove extran.factor of **last Res** by Prop. 1 $\implies \widetilde{G}_1$
- 3) Trace rectPRSs **backwardsly** \implies calc. $\widetilde{G}_2, \widetilde{G}_3, \dots$
 (How to **calc?** \implies **next screen**)
- 4) Apply Buchberger's procedure to $\mathcal{F} \cup \{\widetilde{G}_1, \widetilde{G}_2, \dots\}$

Compare RectPRSs with $\mathbf{GB}(\mathcal{F}_{2018})$

$R_{3:(1,2)} = x^2 y^2 u + \dots$	$\widehat{G}_6 = 27 \text{ digits Coef } x^2 u w + \dots$
$R_{3:(1,2,3)} = 39 y^2 u^6 + \dots$	$\widehat{G}_5 = 26 \text{ digits Coef } x^2 w^2 + \dots$
$R_{4:(1,2,3)} = 1872 y u^{14} + \dots$	$\widehat{G}_4 = 24 \text{ digits Coef } y^2 w + \dots$
$R_{5:(1,2,3)} = 382239 u^{22} + \dots$	$\widehat{G}_3 = 27 \text{ digits Coef } y u + \dots$
	$\widehat{G}_2 = 48000 y w^8 - \dots + \dots$
	$\widehat{G}_1 = 33 u^7 + 23 u^6 w + \dots$
	(notice the coefficient sizes)



Tactics at Step-3 above

: **Eliminate Variables** in **LeadCoefficients**(R_{***})

Elimination of Similar Lead. Coeffs.

given : $R_1, \dots, R_l \in \mathbb{K}[x_m, u]$ ($l \geq 3$)

- i) Let $C_i := \text{LeadCoeff}(R_i)$, $C_1, \dots, C_l \in \mathbb{K}[u]$
- ii) Let $c_i := \text{lastPRS}_{u_1}(C_i, C_{i+1}) \in \mathbb{K}[u_2, \dots, u_n]$
($\text{lastPRS} = \text{last element of PRS}$)
- iii) Finally, let $\bar{c} \simeq \text{gcd}(c_1, \dots, c_l) \in \mathbb{K}[u_2, \dots, u_n]$

We have seen : $\#mn(\bar{H}) \ll \#mn(H_1), \dots, \#mn(H_3)$

We will see : $\#mn(\bar{c}) \ll \#mn(c_1), \dots, \#mn(c_3)$
($\#mn(P) = \#$ of monomials in P)

Important NOTE on \bar{c}

(often $\text{gcd}(c_1, \dots, c_l) \notin \langle \mathcal{F} \rangle$)

Let $\bar{c} = \alpha_1 c_1 + \dots + \alpha_l c_l$, $\alpha_j \in \mathbb{K}[u_2, \dots, u_n]$.

We compute \bar{c} as $\bar{c} = \hat{c} \text{gcd}(c_1, \dots, c_l)$,
where $\hat{c} \in \mathbb{K}[u_3, \dots, u_n]$ is determined
to make polynomials $\alpha_1, \dots, \alpha_l$ a.s.a.p.

Anyway, $\bar{c} \notin \langle \mathcal{F} \rangle$: How to Use \bar{c} ?

\spadesuit : from c_j, \bar{c} to Polynomials in $\langle \mathcal{F} \rangle$

Let $c_i = a_i C_i + b_i C_{i+1}$, $a_i, b_i \in \mathbb{K}[\mathbf{u}]$ (\Leftarrow Elim u_1)

Let $\bar{c} = \alpha_1 c_1 + \cdots + \alpha_l c_l$, $\alpha_1, \dots, \alpha_l \in \mathbb{K}[u_2, \dots, u_n]$

We define $\mathbf{LCtoW}(c_i) = W_i \stackrel{\text{def}}{=} a_i R_i + b_i R_{i+1}$:
 (LC to Whole-polynomial)

$\mathbf{LCtoW}(c_i) \in \langle \mathcal{F} \rangle$, s.t. $\mathbf{LeadCoef}(\mathbf{LCtoW}(c_i)) = c_i$

We define $\overline{\mathbf{LCtoW}}(\bar{c}) \stackrel{\text{def}}{=} \alpha_1 W_1 + \cdots + \alpha_l W_l$:

$\overline{\mathbf{LCtoW}}(\bar{c}) \in \langle \mathcal{F} \rangle$, s.t. $\mathbf{LeadCoef}(\overline{\mathbf{LCtoW}}(\bar{c})) = \bar{c}$

Let's Test above Scheme by \mathcal{F}_{2018}

As mentioned, we use **Spol** only in Buchberger-step

However, we use **Mreduction** indispensably
 (**Monomial** reduction)

Mreduce (F, G) : Mreduce F **fully** by G , i.e.,

$F \xrightarrow{G} R$: each term of R is **Mirreducible** by G

$F = QG + R$: $\text{quopol}(F, G) = Q$, $\text{rempol}(F, G) = R$

==== **in Testing** ====

1) : Use \mathbb{F}_p , $p = 1073738848$, to simplify coefficients

2) : **Mreduce** $\{R_1, R_2, R_3\}$ ($\{C_1, C_2, C_3\}$, too) by G_1

$R_i \xrightarrow{G_1} R'_i$, $C_i \xrightarrow{G_1} C'_i$ (' means Mred by G_1)

Computation of Second-Lowest \tilde{G}_2

given $R'_1, R'_2, R'_3 \in \mathbb{F}_p[y, u, w]$, $\deg_y(R'_i) = 1$ ($\forall i$)

$C'_i := \text{leadCoef}(R'_i) \in \mathbb{F}_p[u, w]$, $\deg_u(C'_i) = 6$

$\tilde{c}'_j := \text{lastPRS}_u(C'_j, C'_{j+1}) \in \mathbb{F}_p[w]$

$$\begin{cases} c'_1 = 182913124 w^{79} - 310233643 w^{78} + \dots + 301414704 w^{11}, \\ c'_2 = 504782002 w^{79} + 105447348 w^{78} + \dots + 465634055 w^{11}, \\ c'_3 = -242692664 w^{67} - 17207621 w^{66} + \dots + 211285272 w^{11}. \end{cases}$$

$\bar{c}' := \text{gcd}(c'_1, c'_2, c'_3)$ ($= \text{gcd}(c'_{j_1}, c'_{j_2 \neq j_1})$)

$$\begin{aligned} \bar{c}' = & w^{17} - 56371298 w^{16} + 138243860 w^{15} - 521121094 w^{14} \\ & - 96457750 w^{13} - 382429906 w^{12} - 247496825 w^{11}. \end{aligned}$$

LCtoW(c'_j) and $\overline{\text{LCtoW}}(\bar{c})$ for \tilde{G}_2

$$c'_i = a_i C'_1 + b_i C'_{i+1} \Rightarrow W'_i := \text{LCtoW}(c'_i) : \#mn(W'_i) = \mathbf{1016}$$

$$\bar{c} = \alpha_i c'_i + \beta_i c'_{i+1} \Rightarrow \overline{W}' := \overline{\text{LCtoW}}(\bar{c}) : \#mn(\overline{W}') = \mathbf{1686}$$

$$\text{We get } \overline{W}'' := \text{Mreduce}(\overline{W}', \hat{G}_1) \Rightarrow : \#mn(\overline{W}'') = \mathbf{61}$$

$$\begin{aligned} \overline{W}'' = & y \times (w^{17} - 56371298 w^{16} + \dots - 247496825 w^{11}) \\ & + u^6 \times (503315083 w^{14} + 511368115 w^{13} + \dots + 365540993 w^9) \\ & + u^5 \times (123032576 w^{15} + 461931391 w^{14} - \dots + 29125264 w^9) \\ & \vdots \qquad \qquad \qquad \vdots \qquad \qquad \qquad \vdots \\ & + u^0 \times (357912951 w^{18} + 304225978 w^{17} - \dots - 342880717 w^{12}). \end{aligned}$$

We see $\overline{W}'' = w^9 \hat{G}_2$, w^9 is extraneous.

How can we **remove** w^9 in \overline{W}'' ?

==== **We found that** ====

Although we have $a'_j R'_j + b'_j R'_{j+1} \propto w^9$, we have

$\text{redpol}(a'_j R'_j, w^9) = - \text{redpol}(b'_j R'_{j+1}, w^9)$
 both sides contain $w^0, w^1, \dots, w^7, w^8$ -terms

Proposition 2 (SSIK2020)

The w^j -terms ($\forall j \leq 8$) in the CofGs in u ,
 of \overline{W}'_j and \overline{W}''_j can be cut off.

Proof of Proposition 2

\overline{W}'_j is expressed by CofGs $a'_{j,1}, a'_{j,2}, a'_{j,3}$ as $\overline{W}'_j =$
 $\underline{C_{\text{ofG}}(\%P[1], \%P[2], \%P[3])} := a'_{j,1} \%P[1] + a'_{j,2} \%P[2] + a'_{j,3} \%P[3]$,
 where $\%P[i]$ is a system variable representing F_i .

Each $F_i(\mathbf{0}, u, w)$ has a nonzero w^0 -term, hence if we
 substitute $F_i(\mathbf{0}, u, w)$ for $\%P[i]$, $i \in \{1, 2, 3\}$, all the
 w^e -terms, $0 \leq \forall e \leq 8$, of $a'_{j,1}, a'_{j,2}, a'_{j,3}$ cancel.

Since $\overline{W}'_j = C_{\text{ofG}}(F_1(\mathbf{0}, u, w), F_2(\mathbf{0}, u, w), F_3(\mathbf{0}, u, w))$,
 this cancellation does **not** change \overline{W}'_j itself.

(The same reasoning is applicable to \overline{W}''_j , too.) □

Another Useful Technique

(We show technique by computing \tilde{G}_4
where $\hat{G}_4 = 17615 \cdots y^2 \underline{\underline{w}} + \cdots$)

given $R'_1, R'_2, R'_3 \in \mathbb{F}_p[y, u, w]$, $\deg_y(R'_i) = 2$,
 $\tilde{c}'_j := \text{lastPRS}_u(C'_j, C'_{j+1})$, $C'_j := \text{LCoef}(R'_j)$,
 compute $\bar{c}' := \text{gcd}(\tilde{c}'_1, \tilde{c}'_2, \tilde{c}'_3)$, then we obtained
 $\bar{c}' = -14400000 \underline{\underline{w^{14}}} + \cdots + 51678000 \underline{\underline{w^5}}$
 \bar{c}' is too higher-order than $\text{LCoef}_y(\hat{G}_4) = c w$.

We can decrease order(\bar{c}') easily

Our Method : $\tilde{c}'_4 := \text{gcd}(\bar{c}', \text{LCoef}_y(\tilde{G}_2))$
 $\Rightarrow \tilde{G}'_4 := \text{LCtoW}(\tilde{c}'_4) = \alpha_4 \overline{W}'_4 + \beta_4 \underline{\underline{y \tilde{G}_2}}$
 $\Rightarrow \tilde{G}_4 := \text{Mreduce}(\tilde{G}'_4, \tilde{G}_2, \tilde{G}_1)$, then

$$\begin{aligned} \tilde{G}_4 = & y^2 \times (260166204 w^2) \\ & + y^1 \times [u^6 \times (890901532 w^{16} + \cdots + 736495066 w - 263471195) \\ & \quad + u^5 \times (-360952533 w^{17} + \cdots - 539864510 w - 470888958) \\ & \quad \quad \quad \vdots \quad \quad \quad \vdots \quad \quad \quad \vdots \\ & + y^0 \times [u^6 \times (890901532 w^{16} + \cdots + 736495066 w - 263471195) \\ & \quad + u^5 \times (-360952533 w^{17} + \cdots - 539864510 w - 470888958) \\ & \quad \quad \quad \vdots \quad \quad \quad \vdots \quad \quad \quad \vdots \end{aligned}]]$$

How to Treat $m \gg 1$ case?

(mixed-triangular Elimination)

$$\{F_1, F_2, \dots, F_{m+1}\} \subset \mathbb{K}[x_1, \dots, x_m, u] \Rightarrow \\ \{F_1, F_2, F_3\} \cup \{F_1, F_2, F_4\} \cup \dots \cup \{F_1, F_2, F_{m+1}\}$$

$$\Rightarrow \text{rectPRS}_{x_1, x_2}(F_1, F_2, F_i) \quad (x_1, x_2 \text{ eliminated})$$

$$\Rightarrow \{G_{i,1}, G_{i,2}, G_{i,i}\} \subset \mathbb{K}[x_3, \dots, x_m, u]$$

$$\Rightarrow \hat{G}_i := \text{gcd}(G_{i,1}, G_{i,2}, G_{i,i}) \quad (i = 3, \dots, m+1)$$

$$\Rightarrow \{\hat{G}_3, \hat{G}_4, \dots, \hat{G}_{m+1}\} \subset \mathbb{K}[x_3, \dots, x_m, u]$$

Continue the above elimination

(We have NOT tested yet)

How to Treat **Non-Healthy** systems?

(various Computation-Branchings occur)

♠ 2-Polynomial (sub-)Systems

- if $(G, H) = (DG', DH')$, $D \notin \mathbb{K}$
then $\text{GB}(\{G, H\}) = D \times \text{GB}(\{G', H'\})$

♣ $(m+1)$ -Polynomial Systems

- separate $\{F_1, \dots, F_{m+1}\}$ into mutually disconnected systems
- separate $\text{GB}(G'(u') G''(u''))$ into $\text{GB}(G'(u'))$ & $\text{GB}(G''(u''))$, where $\text{LMvars}(G') \cap \text{LMvars}(G'') = \emptyset$: NOT yet

What is **Bridge** between E & B ?

Bridge = Coefficients of Generators



PRS : **LCtoW**-polynomial

GB : **Mreduce**-operation

♠ : collaboration is **UN**believably **NICE**

What is the **NEXT** Work ?

Develop Computational Techniques
for **big PRSs** & **Coef-of-Generators**

THANK YOU VERY MUCH
for YOUR ATTENTION

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