## Planar maps: bijections and applications

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## Motivations for bijections

- efficient manipulation of maps (random generation algo.)
- key ingredient to study distances in random maps
typical distances of order $n^{1 / 4}\left(\neq n^{1 / 2}\right.$ in random trees $)$ today's
Theo: [Le Gall, Miermont'13] $\quad M_{n}:=$ random quadrangulation $n$ faces Upon rescaling distances by $n^{1 / 4}, M_{n}$ converges to a continuum random metric space called the Brownian map
( $\mathbf{R k}$ : for random trees, rescaling by $n^{1 / 2}$, convergence to CRT) large tree



## The 2-point function

- Let $\mathcal{G}=\cup_{n} \mathcal{G}[n]$ be a family of maps (or trees, or graphs) where $n$ is a size-parameter (\# faces, \# edges, \# vertices,...)
- Let $\mathcal{G}^{\circ \circ}=$ family of objects from $\mathcal{G}$ with 2 marked vertices $v_{1}, v_{2}$
(or one marked vertex and one rooted edge, etc.)
let $\mathcal{G}_{d}^{\circ \circ}:=$ subfamily of $\mathcal{G}^{\circ \circ}$ where $\operatorname{dist}\left(v_{2}, v_{2}\right)=d$


The counting series $G_{d} \equiv G_{d}(t)$ of $\mathcal{G}_{d}^{\circ \circ}$ with respect to the size is called the 2 -point function of $\mathcal{G}$

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- Let $X_{n}:=\operatorname{dist}\left(v_{1}, v_{2}\right)$ in a random object from $\mathcal{G}^{\circ \circ}[n]$

Then $\mathbb{P}\left(X_{n}=d\right)=\frac{\left[t^{n}\right] G_{d}(t)}{\left[t^{n}\right] G^{\circ \circ}(t)}$

## The 2-point function of plane trees

Consider a random plane tree on $n$ edges with two marked corners


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\text { distance } d=3
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Lagrange inversion $\Rightarrow\left[t^{n}\right] G_{d}(t)=\frac{d+1}{n}\binom{2 n+2}{n+d}$

$$
\sim \frac{4^{n+1}}{n \sqrt{\pi}} x e^{-x^{2}} \quad \text { for } \frac{d}{\sqrt{n}} \rightarrow x
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and we have $\left[t^{n}\right] G^{\circ \circ}(t)=\operatorname{Cat}_{n} \cdot(2 n-1) \sim \frac{2 \cdot 4^{n}}{\sqrt{\pi n}}$

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\text { where } C(t)=1+t C(t)^{2} \\
\end{array} \\
& =C(t)^{2} E(t)^{d} & \begin{array}{l}
d \text { th power of series } E(t)=C(t)-1=\frac{1-2 t-\sqrt{1-4 t}}{2 t} \\
\end{array} \quad \begin{array}{l}
\text { with square-root singularity (also explains } \sqrt{n})
\end{array}
\end{array}
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## The Schaeffer bijection

bijection between: - vertex-pointed quadrangulations with $n$ faces - well-labelled trees with $n$ edges and min-label $=1$


Local rule in each face:


Crucial property: the label $\ell(v)$ of a vertex is its distance (in $Q$ ) from the pointed vertex

Let $\mathcal{G}=\cup_{n} \mathcal{G}[n]=$ family of quadrangulations, with $n=\#$ (faces)
Let $\mathcal{G}^{\circ \circ}=\{$ quadrangulations + marked vertex $v+$ marked edge $e\}$ $\mathcal{G}_{d}^{\circ \circ}:=\operatorname{subfamily}$ where $\operatorname{dist}(v, e)=d+1$


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Then $G_{d}(t)=$ generating function (by edges) of rooted well-labelled trees with root-vertex label $=d$ and min-label $=1$
Rk: Let $R_{i}(t)=$ GF of rooted well-labelled trees where root-label $=i$ and min-label $\geq 1$
Then $G_{d}(t)=R_{d}(t)-R_{d-1}(t)$

## An equation system for the $R_{i}(t)$

## [Bouttier, Di Francesco, Guitter'03]

$R_{i}(t)$ counts


$$
=\frac{1}{1-t\left(R_{i-1}(t)+R_{i}(t)+R_{i+1}(t)\right)}
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Hence the $R_{i}(t)$ are specified by (infinite!) equation system:

$$
R_{0}=0, \quad R_{i}(t)=1+t R_{i}(t) \cdot\left(R_{i-1}(t)+R_{i}(t)+R_{i+1}(t)\right) \text { for } i \geq 1
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$\mathbf{R k}$ : The series $R(t)=\lim _{i \rightarrow \infty} R_{i}(t)$ satisfies $R(t)=1+3 t R(t)^{2}$

$$
=\sum_{n \geq 0} 3^{n} \operatorname{Cat}_{n} t^{n}
$$

## Computing the $R_{i}(t)$ iteratively

We have $R_{1}(t)=\sum_{n \geq 0} \frac{2 \cdot 3^{n}(2 n)!}{n!(n+2)!} t^{n}=R-t R^{3}$
and for $i \geq 1$ we have $R_{i}=1+t R_{i} \cdot\left(R_{i-1}+R_{i}+R_{i+1}\right)$
$\Downarrow$

$$
R_{i+1}=\frac{R_{i}-1}{t R_{i}}-R_{i-1}-R_{i}
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$\Rightarrow$ compute $R_{2}, R_{3}, \ldots$ iteratively
each $R_{i}$ has a rational expression in $t$ and $R$ hence has a rational expression in $R$ (since $t=\frac{R-1}{3 R^{2}}$ )

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this rational expression does not take a nice form by a simple inspection
[Bouttier, Di Francesco, Guitter'03]
First step: ansatz $R_{i}(t)=R(t) \cdot\left(1-c(t) \cdot x(t)^{i}+O\left(x^{2 i}\right)\right)$ with $x(t)$ to be determined
Rk: should have $x(t)=\Theta(t)$ as $t \rightarrow 0$ since $R(t)-R_{i}(t)=\Theta\left(t^{i}\right)$
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Inject into equation $R_{i}=1+t R_{i} \cdot\left(R_{i-1}+R_{i}+R_{i+1}\right)$

$$
\begin{aligned}
R \cdot\left(1-c x^{i}\right) & =1+t R^{2} \cdot\left(1-c x^{i}\right)\left(3-c x^{i-1}-c x^{i}-c x^{i+1}\right)+O\left(x^{2 i}\right) \\
& \hat{\Downarrow} \epsilon=c x^{i} \\
R(1-\epsilon) & =1+t R^{2} \cdot(1-\epsilon)\left(3-\epsilon\left(x^{-1}+1+x\right)\right)+O\left(\epsilon^{2}\right)
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$$

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$$

extracting coefficient $[\epsilon]$ gives $\quad-R=t R^{2} \cdot\left(-3-x^{-1}-1-x\right)$

$$
\begin{gathered}
R-3 t R^{2}=t R^{2} \cdot\left(1+x^{-1}+x\right) \\
\mathbb{\Downarrow} \\
1+x+x^{-1}=\frac{1}{t R^{2}}
\end{gathered}
$$

## Expressing the $R_{i}(t)$ in terms of $x(t)$

[Bouttier, Di Francesco, Guitter'03]

- We have $1+x+x^{-1}=\frac{1}{t R^{2}}=\frac{3}{R-1}$
hence $R(t)$ is rational in terms of $x(t)$, we find $\quad R=\frac{x^{2}+4 x+1}{x^{2}+x+1}$
[Bouttier, Di Francesco, Guitter'03]
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hence $R(t)$ is rational in terms of $x(t)$, we find $\quad R=\frac{x^{2}+4 x+1}{x^{2}+x+1}$
- We then substitute in the expressions of $R_{1}, R_{2}, R_{3}, \ldots$ in terms of $R$
and recognize the explicit expression $\quad R_{i}=R \frac{\left(1-x^{i}\right)\left(1-x^{i+3}\right)}{\left(1-x^{i+1}\right)\left(1-x^{i+2}\right)}$
[Bouttier, Di Francesco, Guitter'03]
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$$
R_{i}=R \frac{\left(1-x^{i}\right)\left(1-x^{i+3}\right)}{\left(1-x^{i+1}\right)\left(1-x^{i+2}\right)}
$$

- To check that this guessed expression works
we have to check that this gives a power series for each $i \geq 0$,
(true since $\left.x(t)=t R(t)^{2} \cdot\left(1+x(t)+x(t)^{2}\right)\right)$
and that $R_{0}=0, \quad R_{i}=1+t R_{i} \cdot\left(R_{i-1}+R_{i}+R_{i+1}\right)$ for $i \geq 1$
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and that $R_{0}=0, \quad R_{i}=1+t R_{i} \cdot\left(R_{i-1}+R_{i}+R_{i+1}\right)$ for $i \geq 1$
We let $F(x, y)=R(x) \frac{(1-y)\left(1-y x^{3}\right)}{(1-y x)\left(1-y x^{2}\right)} \quad\left(y\right.$ plays the role of $\left.x^{i}\right)$ and check $F(t, x)=1+t(x) F(x, y) \cdot\left(F\left(x, y x^{-1}\right)+F(x, y)+F(x, y x)\right)$


## Exact expression

## [Bouttier, Di Francesco, Guitter'03]

The generating functions $R_{i} \equiv R_{i}(t)$ are expressed as

$$
R_{i}=R \frac{\left(1-x^{i}\right)\left(1-x^{i+3}\right)}{\left(1-x^{i+1}\right)\left(1-x^{i+2}\right)}
$$

with $R \equiv R(t)$ given by $R=1+3 t R^{2}$
and $x \equiv x(t)$ given by $x=t R^{2}\left(1+x+x^{2}\right)$

## References:

- first derivation in BDG'03: 'Geodesic distances in planar graphs'
- combinatorial derivations in
[Bouttier, Guitter'12]: 'planar maps and continued fractions'
(+ general determinant expressions for maps with bounded face-degrees)
[Guitter'17]: 'The distance-dependent two-point function of quadrangulations: a new derivation by direct recursion'


## Asymptotic considerations

- Two-point function of (plane) trees:

$$
\begin{gathered}
G_{d}(t)=\left(t R^{2}\right)^{d} \\
\text { with } R=1+t R^{2}=\frac{1-\sqrt{1-4 t}}{2 t}
\end{gathered}
$$


$d=5$
$G_{d}$ is the $d$ th power of a series having a square-root singularity
$\Rightarrow d / n^{1 / 2}$ converges in law (Rayleigh law, density $u \exp \left(-u^{2} / 2\right)$ )

- Two-point function of quadrangulations:

$$
G_{d}(g) \sim_{d \rightarrow \infty} a_{1} x^{d}+a_{2} x^{2 d}+\cdots
$$

where $x=x(t)$ has a quartic singularity
$\Rightarrow d / n^{1 / 4}$ converges to an explicit law
[BDG'03]
Convergence in the two cases "follows" from (proof by Hankel contour) [Banderier, Flajolet, Louchard, Schaeffer'03]: for $0<\alpha<1$,

$$
x(t) \underset{t \rightarrow 1}{\sim 1}-(1-t)^{\alpha} \Rightarrow\left[t^{n}\right] x^{u n^{\alpha}} \sim \frac{1}{2 \pi n} \int_{0}^{\infty} e^{-s} \operatorname{Im}\left(\exp \left(-u s^{\alpha} e^{i \pi \alpha}\right)\right) \mathrm{d} s
$$



vertex $v$ at edge-distance $2 k$ from the root-vertex

vertex $v$ at face-distance $k$ from the root-vertex

vertex $v$ at edge-distance $2 k$ from the root-vertex


Hence $R_{2 d+1}(t)=$ GF (by edges) of rooted maps + marked vertex $v$ such that $v$ is at face-distance $\leq d$ from root-vertex

vertex $v$ at edge-distance $2 k$ from the root-vertex

vertex $v$ at face-distance $k$ from the root-vertex

Hence $R_{2 d+1}(t)=$ GF (by edges) of rooted maps + marked vertex $v$ such that $v$ is at face-distance $\leq d$ from root-vertex

What about the 2-point function of maps for the edge-distance?
[Ambjørn-Budd'13]
a different bijection between quadrangulations and maps

'opposite’ Schaeffer rules


Hence $R_{d}(t)=$ GF (by edges) of rooted maps + marked vertex $v$ such that $v$ is at edge-distance $\leq d-1$ from root-vertex

## Distances from the meta-bijection $\Phi$ ?

Example for a 0-gonal source (pointed vertex $v_{0}$ )
accessible from $v_{0}$
accessible from $v_{0}$
no ccw cycle



Distances from the meta-bijection $\Phi$ ? inverse bijection can be done via growing a cactus from the mobile other way of doing the inverse bijection by labelling the white corners


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- every white corner of label $i \geq 2$ throws an edge to next corner $i-1$ in a ccw walk around the tree


## Distances from the meta-bijection $\Phi$ ?

 inverse bijection can be done via growing a cactus from the mobile other way of doing the inverse bijection by labelling the white corners

- every white corner of label $i \geq 2$ throws an edge to next corner $i-1$ in a ccw walk around the tree
- then every white corner of label 1 throws an edge to new created vertex
cf Schaeffer's bijection


## Distances from the meta-bijection $\Phi$ ?


white corner $c$

recover the orientation
edge $e$

## Distances from the meta-bijection $\Phi$ ?


white corner $c$ label of $c$

edge $e$
length of 'rightmost walk $P(e)$ starting from $e$

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white corner $c$ label of $c$

edge $e$
length of 'rightmost walk $P(e)$ starting from $e$
(at least geodesic length $L(e)$ )

## Distances from the meta-bijection $\Phi$ ?


white corner $c$ label of $c$

edge $e$
length of 'rightmost walk $P(e)$ starting from $e$ (at least geodesic length $L(e)$ )
[Addario-Berry\&Albenque'13]: for $G_{n}$ a random simple triangulation (or random simple quadrangulation) on $n$ vertices, and $e$ a random edge of $G_{n}$,

$$
\text { length }(P(e)) \sim L(e)
$$

(in their proof that $G_{n}$ converges to Brownian map)
‘Quasi’ 2-point function for simple quadrangulations

the 2-point function w.r.t. length of rightmost walk is

$$
G_{i}(s)=r_{i}(s)-r_{i-1}(s)
$$

where $r_{i}(s)=1+s \cdot r_{i-1}(s) r_{i}(s) r_{i+1}(s)$

similar expression for $r_{i}(s)$ as for $R_{i}(t)$ (cf [Bouttier,Guitter'10])

