

## Initialization

```
Needs["Theorema`"]
```

```
SetGlobals[TraceLevel → 0];
```

```
SetGlobals[JuxtapositionForTimes → True];
```

```
Use[{Built-in["Numbers"], Built-in["Operators"], Built-in["Complex Numbers"],
      Built-in["Quantifiers"], Built-in["Tuples"], Built-in["Connectives"], Built-in["Sets"]}]}]
```

```
SetOptions[FlattenKB, ExpandFunctors → "replace"];
```

```
Off[General::"spell1"];
```

```
Unprotect[Or]; ClearAll[Or]; Protect[Or];
```

```
Off[Unset::"write"]
```

```
DeclareAsIdentifiers[ $\tilde{\alpha}$ ]
```

## Number Domains

## Natural Numbers

```
Definition["N", N = NaturalNumbers[]]
```

```
Definition["Natural Numbers", any[],
NaturalNumbers[] = Functor[N, any[x, y],
  s = ⟨⟩
  -----
   $\in_N[x] \Leftrightarrow \text{is-natural}[x]$ 
   $x >_N y \Leftrightarrow x > y$ 
   $0_N = 0$ 
   $x *_N y = x + y$ 
   $\text{lcm}_N[x, y] = \begin{cases} x & \Leftarrow x > y \\ y & \Leftarrow \text{True} \end{cases}$ 
  ]]
```

## Mma Auxiliary Functions

```
Clear[mymult];
```

```
Begin["System`"];
```

```
mymult[x_, y_] := x * y;
```

```
End[];
```

```
Built-in["Mult",  
  ℹMyMult → mymult]
```

```
Clear[myminus];
```

```
Begin["System`"];
```

```
myminus[x_, y_] := x - y;
```

```
End[];
```

```
Built-in["Minus",  
  ℹMyMinus → myminus]
```

```
Clear[isnumber];
```

```
Begin["System`"];
```

```
isnumber[ℹD | L | R | A | B | [x__] | x | xn] := False;
```

```
isnumber[ℹTimes[ℹD, m__] | ℹTimes[L, m__] | ℹTimes[R, m__] | ℹTimes[A, m__] |  
  ℹTimes[B, m__] | ℹTimes[ℹCeiling[m__], y__] | ℹTimes[x, y__]] := False;
```

```
isnumber[x_] :=  
  (x /. {ℹI → i, ℹE → e, ℹTimes → Times, ℹPlus → Plus, ℹMinus → Minus, ℹDivide → Divide, ℹPower → Power}) ∈  
  Complexes;
```

```
End[];
```

```
Built-in["IsNumber",  
  ℹIsNumber → isnumber]
```

```

Definition["Mma Auxiliary Functions", any[],
  MmaAuxF[] = Functor[N, any[x, y],
    s = ⟨⟩
    -----
    ∈ [x] ↔ mIsNumber[x]
    x N · y = mMyMult[x, y]
    x N - y = mMyMinus[x, y]
  ]]
```

### Reduction Field

```

Definition["Reduction Field", any[D],
  ReductionField[D] = Functor[⟨N, extends[D]⟩, any[x, y],
    s = ⟨⟩
    -----
    ∈ [x] ↔ ∈ [x]
    0 N = 0 D
    1 N = 1 D
    x N + y = x D + y
    -x N = -x D
    x N - y = x D - y
    x N * y = x D * y
    x N / y = x D / y
    x N > y ↔ x D > y
    x N > y ↔ ((x ≠ 0 D) ∧ (y = 0 D))
    rdm[x, y] (* the reduction multiplier of x modulo y *) = {
      x D / y ↔ x ≠ 0 D ∧ y ≠ 0 D}
      0 D ↔ otherwise
    }
    lcdr[x, y] (* the least common reducible of x and y *) = {
      1 D ↔ x ≠ 0 D ∧ y ≠ 0 D}
      0 D ↔ otherwise
    }
  ]]
```

### Rational Number Field

```

Definition["Q", Q = RationalNumberField[]]
```

```

Definition["Rational Number Field", any[],
RationalNumberField[] = Functor[N, any[x, y],
  s = ⟨⟩
  _____
  ∈N[x] ⇔ is-rational[x]
  xN > y ⇔ x > y
  0N = 0
  1N = 1
  xN + y = x + y
  -Nx = -x
  xN - y = x - y
  xN * y = x * y
  ∇Nx = 1 / x
  xN / y = x / y
  powerN[x, 0] = 1N
  powerN[x, y] = xN * powerN[x, y - 1]
]]

```

## Complex Numbers

IsComplex

```
Clear[IsComplex];
```

```
Begin["System`"];
```

```
IsComplex[x_] := (x /. {™I → i, ™E → e}) ∈ Complexes;
```

```
End[];
```

```
Built-in["IsComplex",
™IsComplex → IsComplex]
```

```
$TmaOperatorTranslations = DeleteCases[$TmaOperatorTranslations, tr_ /; MemberQ[tr, ™E | ™I, ∞]];
```

ComplexLt

```
Clear[ComplexLt];
```

```
Begin["System`"];
```

```
ComplexLt[x_, x_] := False;
```

```
ComplexLt[x_, y_] := OrderedQ[{x, y} /. {TM I → i, TME → e}];
```

```
End[];
```

```
Built-in["ComplexLt",  
TMComplexLt → ComplexLt]
```

ComplexNumberField

```
Definition["C", C = ComplexNumberField[]]
```

```
Definition["Complex Number Field", any[],  
ComplexNumberField[] = Functor[N, any[x, y, z],  
s = ⟨  
-----  
x ∈ [x] ⇔ TMIsComplex[x]  
x > y ⇔ TMComplexLt[y, x]  
0 = 0  
1 = 1  
x + y = x + y  
x * y = simplify[x * y]  
x * y * z = (x * y) * z  
-x = -x  
x - y = x - y  
∇ x = 1 / x  
x / y = x / y  
lcm[x, y] = { x ⇐ x > y  
          y ⇐ True  
power[x, y] = Power[x, y]  
exp[x] = Exp[x]  
simplify[x] = mma-simplify[x]  
PrettyPrint[x] = x  
⟩  
]]
```

ComplexNumbersMonoid

```
Definition["CM", CM = ComplexNumbers[]]
```

```

Definition ["Complex Numbers" , any[] ,
ComplexNumbers[] = Functor[N, any[x, y] ,
  s = ⟨ ⟩
  -----
  ∈N[x] ⇔ mIsComplex[x]
  x >N y ⇔ mComplexLt[y, x]
  []N = 0
  x *N y = x + y
  lcmN[x, y] = { x ⇐ x > y
                 y ⇐ True
  }
  ] ]

```

Word Monoids

```

Definition ["Finite Chain" , any[L] ,
FinChain[L] = Functor[C, any[x, y] ,
  s = ⟨ ⟩
  -----
  ∈C[x] ⇔ (x ∈ L)
  x >C y ⇔ ∃i=1,...,|L|} ∃j=i+1,...,|L|} ((x = Li) ∧ (y = Lj))
  ] ]

```

```

Definition ["Cartesian Product" , any[D, D̄, A, B] ,
CartesianProduct[D, D̄] = Functor[N, any[X, Y, X̄, Ȳ] ,
  s = ⟨ ⟩
  -----
  ∈N[X] ⇔ (is-tuple[X] ∧ (|⟨D, D̄⟩| = |X|) ∧ ∏j=1,...,|⟨D, D̄⟩|} ( ∈⟨D, D̄⟩j[Xj] ))
  []N = { []⟨D, D̄⟩j}j=1,...,|⟨D, D̄⟩|
  X >N Y ⇔ ( ∏k=1,...,|⟨D, D̄⟩|} ( ⋂l=1,...,k-1} { Xk >⟨D, D̄⟩k Yk} ) )
  X *N Y = { Xm *⟨D, D̄⟩m Ym}m=1,...,|⟨D, D̄⟩|
  lcmN[⟨X, X̄⟩, ⟨Y, Ȳ⟩] = { lcm⟨D, D̄⟩n[⟨X, X̄⟩n, ⟨Y, Ȳ⟩n]n=1,...,|⟨X, X̄⟩|
  ]
  A × B = CartesianProduct[A, B]

```

Definition["Tuples Monoid", any[A],  
 TuplesMonoid[A] = Functor[N, any[x, y, z,  $\bar{x}$ ,  $\bar{y}$ , X, Y],  
 $s = \langle \rangle$   
 $\in_N [X] \Leftrightarrow \left( \left( \text{is-tuple}[X] \wedge \left( \bigwedge_{i=1, \dots, |X|} \in_A [X_i] \right) \wedge (X_{|X|} \neq 0) \right) \vee (X = \langle \rangle) \right)$   
 $\emptyset_N = \langle \rangle$   
 $\left( \langle x, \bar{x} \rangle_N > \langle \rangle \right) \Leftrightarrow \text{True}$   
 $\left( \langle \rangle_N > \langle \bar{y} \rangle \right) \Leftrightarrow \text{False}$   
 $\left( \langle x, \bar{x} \rangle_N > \langle y, \bar{y} \rangle \right) \Leftrightarrow \left( \bigvee \left\{ \begin{array}{l} (x = y) \wedge \langle \bar{x} \rangle_N > \langle \bar{y} \rangle \\ x >_A y \end{array} \right\} \right)$   
 $x \ast_N \langle \rangle = x$   
 $\langle \rangle \ast_N y = y$   
 $\langle \bar{x}, x \rangle \ast_N \langle \bar{y}, y \rangle =$   

$$\left\{ \begin{array}{l} \left\langle \langle \bar{x}, x \rangle_{k_A} \ast \langle \bar{y}, y \rangle_{k_{k=1, \dots, |\langle \bar{x}, x \rangle|}} \right\rangle \Leftrightarrow (|\langle \bar{y}, y \rangle| = |\langle \bar{x}, x \rangle|) \\ \left\langle \langle \bar{x}, x \rangle_{k_A} \ast \langle \bar{y}, y \rangle_{k_{k=1, \dots, |\langle \bar{x}, x \rangle|}} \right\rangle \geq \left\langle \langle \bar{y}, y \rangle_{k_{k=|\langle \bar{x}, x \rangle|+1, \dots, |\langle \bar{y}, y \rangle|}} \right\rangle \Leftrightarrow |\langle \bar{y}, y \rangle| > |\langle \bar{x}, x \rangle| \\ \left\langle \langle \bar{x}, x \rangle_{k_A} \ast \langle \bar{y}, y \rangle_{k_{k=1, \dots, |\langle \bar{y}, y \rangle|}} \right\rangle \geq \left\langle \langle \bar{x}, x \rangle_{k_{k=|\langle \bar{y}, y \rangle|+1, \dots, |\langle \bar{x}, x \rangle|}} \right\rangle \Leftrightarrow \text{otherwise} \end{array} \right.$$
  
 $\text{lcm}_N[x, \langle \rangle] = x$   
 $\text{lcm}_N[\langle \rangle, y] = y$   
 $\text{lcm}_N[\langle x, \bar{x} \rangle, \langle y, \bar{y} \rangle] = \begin{cases} x - \text{lcm}_N[\langle \bar{x} \rangle, \langle \bar{y} \rangle] & \Leftrightarrow x >_A y \\ y - \text{lcm}_N[\langle \bar{x} \rangle, \langle \bar{y} \rangle] & \Leftrightarrow \text{otherwise} \end{cases}$   
 $\left. \right]$

### Cartesian Product

Definition["Cartesian Power", any[D, n],  
 CartesianPower[D, n] = CartesianProduct $\left[ D \mid_{i=1, \dots, n} \right]$

Definition ["Word Monoid", any[L],

LexWords[L] = Functor[W, any[u, v, w,  $\xi$ ,  $\eta$ ,  $\bar{\xi}$ ,  $\bar{\eta}$ ,  $\bar{\xi}$ ],

$s = \langle \rangle$

$$\in_W [w] \Leftrightarrow \left( \bigwedge_{i=1, \dots, |w|} \left\{ \begin{array}{l} \text{is-tuple}[w] \\ \forall \in [w_i] \end{array} \right. \right)$$

$\boxed{\square}_W (* \text{ empty word } *) = \langle \rangle$

$v * w (* \text{ multiplication of two words } *) = v \approx w$

$u * v * w (* \text{ multiplication of three words } *) = (u * v) * w$

$\langle \bar{\eta} \rangle_i \langle \bar{\xi}, \bar{\eta}, \bar{\xi} \rangle (* \text{ middle divisibility } *) \Leftrightarrow \text{True}$

$\langle \bar{\eta} \rangle_i \langle \bar{\xi} \rangle \Leftrightarrow \text{False}$

$\text{rquot}[\langle \rangle, v] (* \text{ right quotient } *) = \boxed{\square}_W$

$\text{rquot}[\langle \bar{\xi}, \bar{\eta} \rangle, \langle \bar{\eta} \rangle] = \boxed{\square}_W$

$\text{rquot}[w, v] = \text{rquot}\left[\left\langle w_i \mid_{i=1, \dots, |w|-1} \right\rangle, v\right] - w_{|w|}$

$\text{lquot}[w, v] (* \text{ left quotient } *) = \left\langle w_i \mid_{i=1, \dots, |w|-|v|-1} \text{rquot}[w, v] \right\rangle$

$\text{lcm}[\langle \bar{\eta}, \bar{\xi} \rangle, \langle \bar{\xi}, \bar{\xi} \rangle] (* \text{ the least common reducible } *) = \langle \bar{\eta}, \bar{\xi}, \bar{\xi} \rangle$

$\langle \eta, \bar{\eta} \rangle_w > \langle \rangle (* \text{ lexicographic ordering } *) \Leftrightarrow \text{True}$

$\langle \rangle_w > \langle \bar{\eta} \rangle \Leftrightarrow \text{False}$

$\langle \eta, \bar{\eta} \rangle_w > \langle \xi, \bar{\xi} \rangle \Leftrightarrow \left( \bigvee \left\{ \begin{array}{l} (\eta = \xi) \wedge \langle \bar{\eta} \rangle_w > \langle \bar{\xi} \rangle \\ \eta > \xi \end{array} \right. \right)$

Definition ["Word Monoid with Degree Ordering", any[L],

DegWords[L] = where [B = LexWords[L],

Functor[W, any[u, v, w,  $\eta$ ,  $\bar{\eta}$ ],

$s = \langle \rangle$

$\in_W [w] \Leftrightarrow \in_B [w]$

$\boxed{\square}_W = \boxed{\square}_B$

$v * w = v * w$

$u * v * w = u * v * w$

$\langle v \rangle_w \Leftrightarrow \langle v \rangle_B$

$\text{rquot}[w, v] = \text{rquot}[w, v]$

$\text{lquot}[w, v] = \text{lquot}[w, v]$

$\text{lcm}[w, v] = \text{lcm}[w, v]$

$\langle \eta, \bar{\eta} \rangle_w > \langle \rangle (* \text{ degree ordering } *) \Leftrightarrow \text{True}$

$\langle \rangle_w > \langle \bar{\eta} \rangle \Leftrightarrow \text{False}$

$\langle v \rangle_w \Leftrightarrow \left( \bigvee \left\{ \begin{array}{l} |v| > |w| \\ \bigwedge \left\{ \begin{array}{l} |v| = |w| \\ v > w \end{array} \right. \right. \right. \right)$

Free Module

Definition ["Free Module", any[K, B],

FreeModule[K, B] = Functor[V, any[c, d, x, y, ξ, η, A, x̄, ȳ],

s = ⟨⟩

$$\in_v[x] \Leftrightarrow \text{where} \left[ z = |x|, \bigwedge_{i=1, \dots, z} \left( \begin{array}{l} \text{is-tuple}[x_i] \\ |x_i| = 2 \\ \text{is-coeff}[(x_i)_1] \\ \text{is-bvec}[(x_i)_2] \\ \bigwedge_{i=1, \dots, z-1} (x_i)_2 \geq_B (x_{i+1})_2 \end{array} \right) \right]$$

$$\text{is-bvec}[\xi] \Leftrightarrow \in_B[\xi]$$

$$\text{is-coeff}[c] \Leftrightarrow \left( \in_K[c] \wedge c \neq 0_K \right)$$

$$\text{bvec}[\xi] = \left\langle \left\langle \frac{1}{K}, \xi \right\rangle \right\rangle$$

$$\text{to-bvec}[\langle\langle c, \xi \rangle\rangle] = \begin{cases} \xi & \leftarrow c = \frac{1}{K} \\ \left[ \frac{c}{B} \right] & \leftarrow \text{otherwise} \end{cases}$$

$$\text{to-bvec}[A] = \left[ \frac{A}{B} \right]$$

$$0_v = \langle \rangle$$

$$\langle \rangle +_v y = y$$

$$x +_v \langle \rangle = x$$

$$\langle\langle c, \xi \rangle, \bar{x} \rangle +_v \langle\langle d, \eta \rangle, \bar{y} \rangle = \begin{cases} \langle c, \xi \rangle - \left( \langle \bar{x} \rangle +_v \langle\langle d, \eta \rangle, \bar{y} \rangle \right) & \leftarrow \xi >_B \eta \\ \langle d, \eta \rangle - \left( \langle\langle c, \xi \rangle, \bar{x} \rangle +_v \langle \bar{y} \rangle \right) & \leftarrow \eta >_B \xi \\ \left\langle c +_K d, \xi \right\rangle - \left( \langle \bar{x} \rangle +_v \langle \bar{y} \rangle \right) & \leftarrow (\xi = \eta) \wedge c +_K d \neq 0_K \\ \langle \bar{x} \rangle +_v \langle \bar{y} \rangle & \leftarrow \text{otherwise} \end{cases}$$

$$V[+] [x, \bar{x}] = x +_v V[+] [\bar{x}] \quad \left. \vphantom{V[+] [x, \bar{x}]} \right] \right]$$

$$\bar{v} \langle \rangle = \langle \rangle$$

$$\bar{v} \langle\langle c, \xi \rangle, \bar{x} \rangle = \left\langle \frac{c}{K}, \xi \right\rangle - \left( \bar{v} \langle \bar{x} \rangle \right)$$

$$x \bar{v} y = x +_v \left( \bar{v} y \right)$$

$$0_v \bar{v} y (* \text{ scalar multiplication } *) = \langle \rangle$$

$$c \bar{v} \langle \rangle = \langle \rangle$$

$$c \bar{v} \langle\langle d, \eta \rangle, \bar{y} \rangle = \left\langle c +_K d, \eta \right\rangle - c \bar{v} \langle \bar{y} \rangle$$

$$y \bar{v} 0 = \langle \rangle$$

$$\langle \rangle \bar{v} c = \langle \rangle$$

$$\langle\langle c, \xi \rangle, \bar{x} \rangle \bar{v} d = \left\langle c +_K d, \xi \right\rangle - \langle \bar{x} \rangle \bar{v} d$$

$$\text{coord}_v[\langle \rangle, x] (* \text{ coordinate wrt a basis } *) = \langle \rangle$$

$$\text{coord}_v[\langle\langle c, \xi \rangle, \bar{x} \rangle, \xi] = c$$

$$\text{coord}_v[\langle\langle c, \xi \rangle, x, \bar{x} \rangle, \eta] = \text{coord}_v[\langle x, \bar{x} \rangle, \eta]$$

$$\text{coord}_v[\langle\langle c, \xi \rangle\rangle, \eta] = 0_K$$

$$\left( \langle \rangle >_v y \right) (* \text{ ordering } *) \Leftrightarrow \text{False}$$

$$\left( \langle\langle c, \xi \rangle, \bar{x} \rangle >_v \langle \rangle \right) \Leftrightarrow \text{True}$$

$$\left( \langle\langle c, \xi \rangle, \bar{x} \rangle >_v \langle\langle d, \eta \rangle, \bar{y} \rangle \right) \Leftrightarrow \bigvee \left( \begin{array}{l} \xi >_B \eta \\ \bigwedge \left\{ \begin{array}{l} \xi = \eta \\ c >_K d \end{array} \right\} \\ \bigwedge \left\{ \begin{array}{l} \xi = \eta \\ c = d \\ \langle \bar{x} \rangle >_v \langle \bar{y} \rangle \end{array} \right\} \end{array} \right)$$

## Monoid Algebra

```

Definition["Monoid Algebra", any[K, W],
MonoidAlgebra[K, W] = where[V = FreeModule[K, W],
  Functor[P, any[c, d, f, g, ξ, η, f̄, m̄, n̄],
    s = ⟨⟩
    -----
    εP[f] ⇔ εV[f]
    is-bvecP[ξ] ⇔ is-bvecV[ξ]
    is-coeffP[c] ⇔ is-coeffV[c]
    0P = 0V
    P[+][f̄] = V[+][f̄]
    fP = fV
    fP · g = fV · g
    ewP[ξ] = bvecV[ξ]
    c · f = c · f
    coordP[f, ξ] = coordV[f, ξ]
    -----
    1P = ⟨⟨1K, 0W⟩⟩
    ⟨⟩ *P g (* multiplication of two polynomials *) = ⟨⟩
    f *P ⟨⟩ = ⟨⟩
    ⟨⟨c, ξ⟩, m̄⟩ *P ⟨⟨d, η⟩, n̄⟩ = ⟨⟨c *K d, ξ *W η⟩⟩ *P ⟨⟨c, ξ⟩⟩ *P ⟨⟨n̄⟩⟩ *P ⟨⟨m̄⟩⟩ *P ⟨⟨d, η⟩, n̄⟩
    P[*][f, f̄] = f *P P[*][f̄]
    (⟨⟩ >P g) (* ordering *) ⇔ False
    (⟨⟨c, ξ⟩, m̄⟩ >P ⟨⟩) ⇔ True
    (⟨⟨c, ξ⟩, m̄⟩ >P ⟨⟨d, η⟩, n̄⟩) ⇔ ⎧ ⎨ ⎩
      ⎧ ξ >W η
      ⋀ { ξ = η
          c >K d
        }
      ⎧ ξ = η
      ⋀ { c = d
          ⟨m̄⟩ >P ⟨n̄⟩
        }
    ⎫ ⎬ ⎭
    rrdmP(⟨⟩, g) (* right reduction multiplier *) = 0P
    rrdmP(⟨⟨c, ξ⟩, m̄⟩, ⟨⟩) = 0P
    rrdmP(⟨⟨c, ξ⟩, m̄⟩, ⟨⟨d, η⟩, n̄⟩) = ⎧ ⎨ ⎩
      ⎧ ⟨⟨rdmK[c, d], rquotW[ξ, η]⟩⟩ ⇔ rdmK[c, d] ≠ 0 ∧ ηW | ξ
      0P ⇔ otherwise
    ⎫ ⎬ ⎭
    lrdmP(⟨⟩, f) (* left reduction multiplier *) = 0P
    lrdmP(⟨⟨c, ξ⟩, m̄⟩, ⟨⟩) = 0P
    lrdmP(⟨⟨c, ξ⟩, m̄⟩, ⟨⟨d, η⟩, n̄⟩) = ⎧ ⎨ ⎩
      ⎧ ⟨⟨1K, lquotW[ξ, η]⟩⟩ ⇔ rdmK[c, d] ≠ 0 ∧ ηW | ξ
      0P ⇔ otherwise
    ⎫ ⎬ ⎭
    lcrdP(⟨⟨c, ξ⟩, m̄⟩, ⟨⟨d, η⟩, n̄⟩) (* least common reducible *) = ⟨⟨lcrdK[c, d], lcmW[ξ, η]⟩⟩
  ]]]

```

Definition["Groebner Extension", any[P],

GB[P] = Functor[⟨G, extends[P]⟩, any[c, f, g, k, l, r, t, ξ, lhs1, lhs2, rhs1, rhs2,  $\bar{f}$ ,  $\bar{m}$ , R, S],

$$\frac{s = \langle \rangle}{f \underset{G}{>} g \Leftrightarrow f \underset{P}{>} g}$$

$$\frac{0 = 0}{\underset{G}{0} = \underset{P}{0}}$$

$$\frac{1 = 1}{\underset{G}{1} = \underset{P}{1}}$$

$$f \underset{G}{+} g \Leftrightarrow f \underset{P}{+} g$$

$$G[*][f, \bar{f}] = P[*][f, \bar{f}]$$

$$\text{hredp}[f, g] (* \text{ head reduction modulo polynomial } *) = f \underset{P}{\leftarrow} \text{lrdm}[f, g] \underset{P}{*} g \underset{P}{*} \text{rrdm}[f, g]$$

$$\text{hredp}[f, g, l, r] = f \underset{P}{\leftarrow} l \underset{P}{*} g \underset{P}{*} r$$

$$\text{hred}[f, S] (* \text{ head reduction modulo system } *) = \text{where}[q = S[f],$$

$$\quad \text{hredp}[f, q_2, q_1, q_3]]$$

$$\text{sred}[\langle \rangle, S, t] (* \text{ step reduction modulo system at given position } *) = 0 \underset{P}$$

$$\text{sred}[\langle \langle c, \xi \rangle, \bar{m} \rangle, S, 1] = \text{hred}[\langle \langle c, \xi \rangle, \bar{m} \rangle, S]$$

$$\text{sred}[\langle \langle c, \xi \rangle, \bar{m} \rangle, S, t] = \langle \langle c, \xi \rangle \rangle \underset{G}{+} \text{sred}[\langle \bar{m} \rangle, S, t - 1]$$

$$\text{sred}[f, S] (* \text{ step reduction modulo system } *) = \text{cred}[f, S, 1]$$

$$\text{cred}[f, S, k] (* \text{ complete reduction modulo system at given position } *) =$$

$$\text{where}[f1 = \text{sred}[f, S, k],$$

$$\quad \left\{ \begin{array}{l} \text{cred}[f, S, k+1] \leftarrow (f1 = f) \wedge (k < |f|) \\ f1 \leftarrow \text{otherwise} \end{array} \right\} \quad ]]$$

$$\text{tred}[f, S] (* \text{ total reduction modulo system } *) = \text{where}[q = \text{sred}[f, S],$$

$$\quad \left\{ \begin{array}{l} \text{tred}[q, S] \leftarrow q \neq f \\ f \leftarrow \text{otherwise} \end{array} \right\}]$$

$$\text{spol}[f, g] (* \text{ the S-polynomial of f and g } *) = \text{where}[L = \text{lcrd}[f, g],$$

$$\quad \text{hredp}[L, f] \underset{G}{-} \text{hredp}[L, g]]$$

$$\text{spolred}[\langle f, g \rangle, S] = \text{tred}[\text{spol}[f, g], S]$$

$$\text{Gb}[R, S] = \text{where}[\text{pairs} = \left\langle \langle R_i, R_j \rangle \mid \begin{array}{l} i=1, \dots, |R| \\ j=1, \dots, |R| \\ R_i \neq R_j \end{array} \right\rangle,$$

$$\quad \text{Gb}[R, \text{pairs}, S]]$$

$$\text{Gb}[R, \langle \rangle, S] = R$$

$$\text{Gb}[R, \langle \langle f, g \rangle, \bar{m} \rangle, S] = \text{where}[h = \text{spolred}[\langle f, g \rangle, S],$$

$$\quad \left\{ \begin{array}{l} \text{Gb}[R, \langle \bar{m} \rangle, S] \leftarrow h = 0 \underset{P} \\ \text{Gb}[R - h, ((h \rightarrow R) \times \langle \bar{m} \rangle) \times (R \leftarrow h), S] \leftarrow \text{otherwise} \end{array} \right\}$$

$$\quad ]]$$

Definition["Theorema general functions", any[F, l, r],

$$l \rightarrow F = \left\langle \langle l, F_i \rangle \mid \begin{array}{l} i=1, \dots, |F| \end{array} \right\rangle$$

$$F \leftarrow r = \left\langle \langle F_i, r \rangle \mid \begin{array}{l} i=1, \dots, |F| \end{array} \right\rangle$$

## Exponential Polynomials

## MakeBoxes&amp;MakeExpressions (Formatting&amp;Parsing)

```
MakeExpression[RowBox[{SubscriptBox["∂", Q_], p_}], StandardForm] :=
  MakeExpression[RowBox[{RowBox[{Q, "[", "™Der", "]" }], "[", p, "]" }], StandardForm];
```

```
MakeExpression[RowBox[{SubscriptBox["∫", Q_], p_}], StandardForm] :=
  MakeExpression[RowBox[{RowBox[{Q, "[", "™Int", "]" }], "[", p, "]" }], StandardForm];
```

```
MakeBoxes[Q_["™Der"][p_], StandardForm] :=
  RowBox[{SubscriptBox["∂", MakeBoxes[Q, StandardForm]], MakeBoxes[p, StandardForm]}];
```

```
MakeBoxes[Q_["™Int"][p_], StandardForm] :=
  RowBox[{SubscriptBox["∫", MakeBoxes[Q, StandardForm]], MakeBoxes[p, StandardForm]}];
```

```
MakeExpression[RowBox[{"[" , x_ , "]" }], StandardForm] :=
  MakeExpression[RowBox[{"™Ceiling", "[" , x , "]" }], StandardForm];
```

```
MakeBoxes["™Ceiling"[x_], StandardForm] := RowBox[{"[" , MakeBoxes[x, StandardForm], "]" }];
```

```
MakeExpression[RowBox[{"[" , x_ , "]" }], StandardForm] :=
  MakeExpression[RowBox[{"™Floor", "[" , x , "]" }], StandardForm];
```

```
MakeBoxes["™Floor"[x_], StandardForm] := RowBox[{"[" , MakeBoxes[x, StandardForm], "]" }];
```

## Built-ins

```
Begin["System`"];
```

```
MmaSimplify[p_] := FullSimplify[p];
```

```
End[];
```

```
Built-in["MmaSimplify",
  mma-simplify → MmaSimplify]
```

## Exponential Polynomials

```

Definition["Exp Polys", any[K],
ExpPolys[K] = where [P = MonoidAlgebra[K, N × CM],
  Functor[E, any[c, d, m, p, q, λ, ξ, m̄, n̄],
    s = ⟨⟩
    -----
     $\bar{\in}_E [p] \Leftrightarrow \in_P [p]$ 
     $0_E = 0_P$ 
     $1_E = 1_P$ 
     $p >_E q \Leftrightarrow p >_P q$ 
     $p +_E q = p +_P q$ 
     $\bar{p}_E = \bar{p}_P$ 
     $p \bar{p}_E = p \bar{p}_P$ 
     $c \bar{p}_E = c \bar{p}_P$ 
     $p \star_E q = p \star_P q$ 
     $ew_E [p] = ew_P [p]$ 
     $coord_E [p, \xi] = coord_P [p, \xi]$ 
     $is-bvec_E [\xi] \Leftrightarrow is-bvec_P [\xi]$ 
     $is-coeff_E [c] \Leftrightarrow is-coeff_P [c]$ 
    -----
     $lcrd_E [\langle \langle c, \lambda \rangle \rangle, \langle \langle d, \xi \rangle \rangle] = \left\langle \left\langle lcrd_K [c, d], lcm_{N \times CM} [\lambda, \xi] \right\rangle \right\rangle$ 
     $is-scal_E \left[ \left\langle \left\langle c, \begin{smallmatrix} \square \\ N \times CM \end{smallmatrix} \right\rangle \right\rangle \right] \Leftrightarrow is-coeff_P [c]$ 
     $is-scal_E [p] \Leftrightarrow False$ 
     $to-scal_{\mathcal{F}} [\langle \rangle] = 0_K$ 
     $to-scal_E \left[ \left\langle \left\langle c, \begin{smallmatrix} \square \\ N \times CM \end{smallmatrix} \right\rangle \right\rangle \right] = c$ 
     $eval_E [\langle \rangle, d] = 0_K$ 
     $eval_E [\langle \langle c, \langle \rangle \rangle \rangle, d] = c$ 
     $eval_E [\langle \langle c, \langle 0, \lambda \rangle \rangle \rangle, d] = c \star_K \exp_K [\lambda \star_K d]$ 
     $eval_E [\langle \langle c, \langle m, \lambda \rangle \rangle \rangle, \bar{m}, d] = c \star_K power_K [d, m] \star_K \exp_K [\lambda \star_K d] + eval_E [\langle \bar{m} \rangle, d]$ 
     $\partial_E \langle \rangle = 0_P$ 
     $\partial_E \langle \langle c, \langle \rangle \rangle \rangle = 0_P$ 
     $\partial_E \langle \langle c, \langle 0, \lambda \rangle \rangle \rangle, \bar{n} = \lambda \cdot \langle \langle c, \langle 0, \lambda \rangle \rangle \rangle + \partial_E \langle \bar{n} \rangle$ 
     $\partial_E \langle \langle c, \langle m, \lambda \rangle \rangle \rangle, \bar{n} = \left\langle \left\langle c \star_K m, \langle m-1, \lambda \rangle \right\rangle \right\rangle + \lambda \cdot \langle \langle c, \langle m, \lambda \rangle \rangle \rangle + \partial_E \langle \bar{n} \rangle$ 
     $\int_E \langle \rangle = 0_P$ 
     $\int_E \langle \langle c, \langle \rangle \rangle \rangle = \langle \langle c, \langle 1, 0 \rangle \rangle \rangle$ 
     $\int_E \langle \langle c, \langle m, 0 \rangle \rangle \rangle, \bar{n} = \left\langle \left\langle c /_K (m+1), \langle m+1, 0 \rangle \right\rangle \right\rangle + \int_E \langle \bar{n} \rangle$ 
     $\int_E \langle \langle c, \langle 0, \lambda \rangle \rangle \rangle, \bar{n} = \left( \left\langle \left\langle c /_K \lambda, \langle 0, \lambda \rangle \right\rangle \right\rangle - \left\langle \left\langle c /_K \lambda, \langle 0, 0 \rangle \right\rangle \right\rangle \right) + \int_E \langle \bar{n} \rangle$ 
     $\int_E \langle \langle c, \langle m, \lambda \rangle \rangle \rangle, \bar{n} = \left( \left\langle \left\langle c /_K \lambda, \langle m, \lambda \rangle \right\rangle \right\rangle - \int_E m \cdot \left\langle \left\langle c /_K \lambda, \langle m-1, \lambda \rangle \right\rangle \right\rangle \right) + \int_E \langle \bar{n} \rangle$ 
    PrettyPrint_E [p] =  $\text{mPlus} \left[ \text{PrettyPrintM}_E [p_i] \mid_{i=1, \dots, |p|} \right]$ 
    PrettyPrint_E [\langle m \rangle] = PrettyPrintM_E [m]
    PrettyPrint_E [\langle \rangle] = 0
    PrettyPrintM_E [\langle c, \langle m, \lambda \rangle \rangle] =  $\text{mTimes} [c, \text{mTimes} [x^m, \text{mE}^{\lambda x}]]$ 

```

## Built-ins

```
Begin["System`"];
```

```
Clear[CoeffQ, Mon, Coeff, SplitPoly, FuncOrder];
```

```
CoeffQ[expr_] := FreeQ[expr, x];
```

```
Mon[t1_, t2___] := Mon[t1 * t2];
```

```
Mon[Times[c_?CoeffQ, t___] := Mon[t];
```

```
Mon[Times[c_, t___] := Mon[c] * Mon[t];
```

```
Mon[Power[e, Plus[λ_?NumericQ, μ_]] := Power[e, μ];
```

```
Mon[c_?CoeffQ] := 1;
```

```
Mon[s_] := s;
```

```
Coeff[t1_, t2___] := Coeff[t1 * t2];
```

```
Coeff[Times[c_, t___] := Coeff[c] * Coeff[t];
```

```
Coeff[ePlus[λ_?NumericQ, μ_] := eλ;
```

```
Coeff[s_?CoeffQ] := s;
```

```
Coeff[s_] := 1;
```

```
SplitPoly[c_?CoeffQ] := {{c, mma[1]}};
```

```
SplitPoly[Plus[t1_, t2___] := Join[SplitPoly[t1], Flatten[SplitPoly /@ {t2}, 1]];
```

```
SplitPoly[0] := {};
```

```
SplitPoly[s_] := {{FullSimplify[Coeff[s]], mma[Mon[s]]}};
```

```
SplitPoly[t1_, t2____] := SplitPoly /@ {t1, t2};
```

```
MySplitPoly[A_] := SplitPoly[Expand[A]] /. List -> TMTuple;
```

```
FuncOrder[p_, p_] := False;
```

```
FuncOrder[p_, q_] := OrderedQ[{p, q}];
```

```
FuncQ[mma[t____]] := True;
```

```
FuncQ[s_] := False;
```

```
Clear[ToMma];
```

```
ToMma[⟨m_, r____⟩] := Plus[ToMma[m], ToMma[r]];
```

```
ToMma[⟨c_, mma[m_]⟩] := Times[c, m];
```

```
ToMma[t1_, t2____] := Plus[ToMma[t1], ToMma[t2]];
```

```
ToMma[⟨⟩] := 0;
```

```
ToMma[] := 0;
```

```
ToMma[c_] := c;
```

```
Clear[MmaPlus, MmaMinus, MmaTimes, MmaDif, MmaInt];
```

```
MmaPlus[p_, q_] := MySplitPoly[ToMma[p] + ToMma[q]];
```

```
MmaMinus[p_] := MySplitPoly[-ToMma[p]];
```

```
MmaMinus[p_, q_] := MySplitPoly[ToMma[p] - ToMma[q]];
```

```
MmaTimes[p_, q_] := MySplitPoly[ToMma[p] * ToMma[q]];
```

```
MmaDif[p_] := MySplitPoly[Expand[D[ToMma[p], x]]];
```

```
MmaInt[p_] := MySplitPoly[Expand[TrigToExp[Integrate[ToMma[p], {x, 0, x}]]]];
```

```
MmaSubs[p_, c_] := Subs[ToMma[p], c];
```

```
Subs[p_, c_] := p /. x → c;
```

```
MmaSimplify[p_] := FullSimplify[p];
```

```
SetOptions[TransformResultFromMma, PlusNegativeAsMinus → True,  
TimesNegPowerAsFraction → False, PowerRationalAsRoot → False, NegPowerAsReciprocal → False];
```

```
TmaTRFM[p_] := TransformResultFromMma[p /. {TMTimes → Times, TMPlus → Plus, TMMinus → Minus, TMPower → Power}];
```

```
End[];
```

```
Built-in["FuncOrder",  
func-less → FuncOrder]
```

```
Built-in["FuncQ",  
is-func → FuncQ]
```

```
Built-in["MmaDiff",  
mma-diff → MmaDiff]
```

```
Built-in["MmaInt",  
mma-int → MmaInt]
```

```
Built-in["ToMma",  
TMtoMma → ToMma]
```

```
Built-in["MmaPlus",  
mma-plus → MmaPlus]
```

```
Built-in["MmaMinus",  
mma-minus → MmaMinus]
```

```
Built-in["MmaTimes",  
mma-times → MmaTimes]
```

```
Built-in["MmaSubs",  
mma-sub → MmaSubs]
```

```
Built-in["MmaSimplify",  
mma-simplify → MmaSimplify]
```

```
Built-in["TmaTRFM",  
tma-trfm → TmaTRFM]
```

```
Built-in["TmaSplit",  
tma-split → MySplitPoly]
```

```
mytimes
```

```
Clear[mytimes];
```

```
mytimes /: MakeBoxes[mytimes[x_], StandardForm] := MakeBoxes[x, StandardForm];
```

```
mytimes /: MakeBoxes[mytimes[x_, y_], StandardForm] :=  
RowBox[{MakeBoxes[x, StandardForm], MakeBoxes[y, StandardForm]}];
```

```
mytimes /: MakeBoxes[mytimes[x_, y_, z_], StandardForm] := RowBox[  
{MakeBoxes[x, StandardForm], MakeBoxes[y, StandardForm], MakeBoxes[mytimes[z], StandardForm]}];
```

```
mytimes[1, x_] := x;
```

```
mytimes[-1, x_] := -x;
```

```
mytimes[x_, 1] := x;
```

```
mytimes[x_, -1] := -x;
```

## Mathematica Functions

```

Definition["Mma Functions", any[K],
MmaFct[K] = Functor[E, any[c, d, l, p, q, ξ],
s = ⟨⟩

$$\underset{E}{\in}[p] \Leftrightarrow \bigwedge_{i=1, \dots, |p|} \left( \text{is-tuple}[p_i] \wedge \left( \begin{array}{l} \text{is-tuple}[p_i] \\ |p_i| = 2 \\ \text{is-coeff}[(p_i)_1] \\ \text{is-bvec}[(p_i)_2] \end{array} \right) \wedge \bigwedge_{i=1, \dots, |p|-1} \text{func-less}[(p_{i+1})_2, (p_i)_2] \right)$$

0 = ⟨⟩
 $\underset{E}{1} = \langle \langle 1, \text{mma}[1] \rangle \rangle$ 
 $\underset{E}{p} > \underset{E}{q} \Leftrightarrow \text{func-less}[q, p]$ 
 $\underset{E}{p} + \underset{E}{q} = \text{mma-plus}[p, q]$ 
 $\underset{E}{p} - \underset{E}{q} = \text{mma-minus}[p, q]$ 
 $\underset{E}{p} \cdot \underset{E}{q} = \text{mma-times}[p, q]$ 
 $\underset{E}{p} * \underset{E}{q} = \text{mma-times}[p, q]$ 
 $\underset{E}{\text{is-bvec}}[\xi] \Leftrightarrow \text{is-func}[\xi]$ 
 $\underset{E}{\text{is-coeff}}[c] \Leftrightarrow \left( \underset{K}{\in}[c] \wedge c \neq 0 \right)$ 
 $\underset{E}{\text{is-scal}}[\langle \langle c, \text{mma}[1] \rangle \rangle] \Leftrightarrow \underset{E}{\text{is-coeff}}[c]$ 
 $\underset{E}{\text{is-scal}}[p] \Leftrightarrow \text{False}$ 
 $\underset{E}{\text{to-scal}}[\langle \rangle] = 0$ 
 $\underset{E}{\text{to-scal}}[\langle \langle c, \text{mma}[1] \rangle \rangle] = c$ 
 $\underset{E}{\text{eval}}[p, d] = \text{mma-subs}[p, d]$ 
 $\underset{E}{\partial} p = \text{mma-diff}[p]$ 
 $\underset{E}{\int} p = \text{mma-int}[p]$ 
 $\underset{E}{\text{wronmin}}[l] = \text{tma-wronmin}[l]$ 
 $\underset{E}{\text{simplify}}[p] = \text{mma-simplify}[p]$ 
]]

```

## Integro-Differential Operators

```

Definition["Basis", any[F],
Basis[F] = Functor[B, any[ξ, η],
s = ⟨⟩

$$\underset{B}{\in}[\langle "[]" , \xi \rangle] \Leftrightarrow \underset{F}{\text{is-bvec}}[\xi]$$


$$\langle "[]" , \xi \rangle \underset{B}{>} \langle "[]" , \eta \rangle \Leftrightarrow \xi \underset{F}{>} \eta$$

]]

```





Definition["Groebner Basis Normal Form", any[G, S],

GBNF[G, S] = Functor[C, any[a, b, c,  $\bar{a}$ ,  $\mathcal{B}$ , F],

$s = \langle \rangle$

$\frac{\text{NF}[a]}{c} = \text{tred}_G[a, S]$

$\frac{0}{c} = 0$

$\frac{1}{c} = \frac{1}{c}$

$C[+][\bar{a}] = G[+][\bar{a}]$

$\bar{c}a = \bar{c}a$

$a\bar{c}b = a\bar{c}b$

$C[*][\bar{a}] = G[*][\bar{a}]$

$c : a = c : a$

$\text{ew}[a] = \text{ew}[a]$

$\text{EvalMat}[\mathcal{B}, F] = \text{EvalMat}_G[\mathcal{B}, F]$

]]

Definition["Quotient Algebra", any[G],

QuotAlg[G] = Functor[Q, any[a, b, c,  $\bar{a}$ ,  $\mathcal{B}$ , F],

$s = \langle \rangle$

$\frac{\in[a]}{c} \Leftrightarrow \left( \frac{\text{NF}[a]}{c} = a \right)$

$\frac{0}{c} = 0$

$\frac{1}{c} = \frac{1}{c}$

$Q[+][\bar{a}] = \text{NF}_G[G[+][\bar{a}]]$

$\bar{c}a = \text{NF}_G[\bar{c}a]$

$a\bar{c}b = \text{NF}_G[a\bar{c}b]$

$Q[*][\bar{a}] = \text{NF}_G[G[*][\bar{a}]]$

$c : a = \text{NF}_G[c : a]$

$\text{ew}[a] = \text{ew}[a]$

$\text{EvalMat}[\mathcal{B}, F] = \text{EvalMat}_G[\mathcal{B}, F]$

]]

Definition["FreeIntDiffOp", any[ $\mathcal{F}$ , K],

FreeIntDiffOp[ $\mathcal{F}$ , K] =

where  $\mathcal{M} = \text{MonoidAlgebra}[K, \text{DegWords}["\partial" \mapsto (" \int" \mapsto (\text{DoubleBasis}[\text{Basis}[\mathcal{F}], \text{CharBasis}[K]]) )]]$ ,

Functor[ $\mathcal{A}$ , any[b, c, d, f, g,  $\alpha$ ,  $\beta$ ,  $\xi$ , w,  $\bar{f}$ ,  $\bar{g}$ ,  $\bar{m}$ ,  $\bar{w}$ ,  $\mathcal{B}$ , F, J],

```

s = ⟨ ⟩
-----
∈A[f] ⇔ ∈M[f̄]
0 = 0
-----
A[+] [f̄] = M[+] [f̄]
-f̄ = -Mf
f -g = f -Mg
c · f = c ·Mf
-----
A[*] [f̄] = M[*] [f̄]
is-bvec[b] ⇔ is-bvecM[b]
is-scal[⟨⟨c, ⟨⟩⟩⟩] ⇔ ( ∈K[c] )
is-scal[b] ⇔ False
lrdm[f, g] = lrdmM[f, g]
-----
rrdm[f, g] = rrdmM[f, g]
-----
f > g ⇔ f >M g
lcm[f, g] = lcmDegWords["∂"↦("↦(DoubleBasis[Basis[F], CharBasis[K]]))][f, g]
lcrd[⟨⟨⟨c, f⟩, f̄⟩, ⟨⟨⟨d, g⟩, ḡ⟩⟩] (* least common reducible *) = ⟨⟨lcrdK[c, d], lcmA[f, g]⟩⟩
1 = ⟨⟨1K, ⟨⟩⟩⟩
ew[⟨⟩] = ⟨⟩
ew[⟨⟨⟨⟩, g⟩⟩] = ⟨⟩
ew[⟨⟨⟨⟨⟨c, w⟩⟩, ⟨⟨⟨α, β, J⟩⟩, w̄⟩, g⟩, m̄⟩] = (α * c)M ew[⟨⟨⟨⟨⟨1, w⟩⟩, ⟨⟨1, β, J⟩⟩⟩, g⟩] +MonoidAlgebra[K, F]
  ew[⟨⟨⟨w̄⟩, g⟩⟩] +MonoidAlgebra[K, F] ew[⟨m̄⟩]
ew[⟨⟨⟨⟨⟨c, w⟩⟩, ξ⟩, w̄⟩, g⟩, m̄⟩] = c ·M ew[⟨⟨⟨⟨⟨1, w⟩⟩, ξ⟩, g⟩] +MonoidAlgebra[K, F] ew[⟨⟨⟨w̄⟩, g⟩⟩] +MonoidAlgebra[K,
  ew[⟨m̄⟩]
to-scal[⟨⟩] = 0
to-scal[⟨⟨c,   ⟩⟩] = c
op-to-fun[⟨⟩] (* integro-diff operators to functions *) = ⟨⟩
op-to-fun[⟨⟨c, ⟨⟩⟩⟩] = c · 1F
op-to-fun[⟨⟨c, ⟨"[]" , ξ⟩⟩, m̄⟩] = ⟨c, ξ⟩ - op-to-funA[⟨m̄⟩]
fun-to-op[⟨⟩] (* functions to integro-diff operators *) = ⟨⟩
fun-to-op[⟨⟨c, ξ⟩, m̄⟩] = { ⟨⟨to-scalF[⟨⟨c, ξ⟩⟩, ⟨⟩⟩⟩ + fun-to-opA[⟨m̄⟩] ⇐ is-scalA[⟨⟨c, ξ⟩⟩]
  ⟨⟨c, ⟨"[]" , ξ⟩⟩⟩ + fun-to-opM[⟨m̄⟩] ⇐ otherwise
}
⟨⟩ ◊ f (* the action of int-diff word on function *) = fun-to-opA[f]
w̄, "∂" ◊ ⟨⟨c, f⟩⟩ = c · w̄ ◊ ewA[∂Ff]
w̄, "J" ◊ ⟨⟨c, f⟩⟩ = c · w̄ ◊ ewA[∫Ff]
w̄, ⟨"[]" , d⟩ ◊ ⟨⟨c, f⟩⟩ = c · w̄ ◊ ewA[eval[f, d]]
w̄, ⟨"[]" , b⟩ ◊ ⟨⟨c, f⟩⟩ = c · w̄ ◊ ewA[b * f]
⟨⟩ ◊ f (* the action of int-diff op on function *) = ⟨⟩
⟨c, ξ⟩, m̄ ◊ f = c · ξ ◊ f + ⟨m̄⟩ ◊ f
EvalMat[B, F] = ⟨⟨to-scalA[Bi ◊ op-to-funA[Fj]]j=1,...,|F| | i=1,...,|B|⟩⟩
]]

```

```

Definition ["PrettyPrinting for IntDiffPol", any[ $\mathcal{F}$ , K],
PrettyPrintingIDP[ $\mathcal{F}$ , K] = Functor[P, any[b, c, m, n, p, r, w,  $\alpha$ ,  $\beta$ ,  $\lambda$ ,  $\gamma$ ,  $\eta$ , J,  $\bar{\eta}$ ,  $\bar{\xi}$ ,  $\bar{J}$ ],

s = <>
PrettyPrintW[<>] = ""

PrettyPrintW[< $\lambda$ ,  $\eta$ ,  $\bar{\eta}$ >] = {
  Times[ $\lambda^2$ , PrettyPrintW[< $\bar{\eta}$ >]] <-  $\lambda = \eta$ 
  myTimes[PrettyPrintL[ $\lambda$ ], PrettyPrintW[< $\eta$ ,  $\bar{\eta}$ >]] <- otherwise
}
PrettyPrintW[< $\eta$ >] = PrettyPrintL[ $\eta$ ]
PrettyPrintL[<"[">, <>, <>, J>] = PrettyPrintJ[J]
PrettyPrintL[<"[">, <>,  $\beta$ , J>] = myTimes[Power[u,  $\beta$ ], PrettyPrintJ[J]]
PrettyPrintL[<"[">, < $\alpha$ , <>, J>] = myTimes[Power[u[0],  $\alpha$ ], PrettyPrintJ[J]]
PrettyPrintL[<"[">, < $\alpha$ ,  $\beta$ , J>] = myTimes[Power[u[0],  $\alpha$ ], Power[u,  $\beta$ ], PrettyPrintJ[J]]
PrettyPrintJ[<>] = ""
PrettyPrintJ[<<b,  $\gamma$ ,  $\bar{J}$ >] = myTimes[Times["[">, PrettyPrint[b]], Power[u,  $\gamma$ ], PrettyPrintJ[< $\bar{J}$ >]]]
PrettyPrintL[" $\partial$ "] = " $\partial$ "
PrettyPrintL["[">] = "[">
PrettyPrintL[<>] = ""
PrettyPrintL[ $\eta$ ] =  $\eta$ 
PrettyPrintM[<c, w>] = Times[PrettyPrint[c], PrettyPrintW[w]]
PrettyPrint[<<<c,  $\eta$ ,  $\bar{\xi}$ ,  $\bar{\eta}$ >>] = { PrettyPrint[<<<c,  $\eta$ ,  $\bar{\xi}$ ,  $\bar{\eta}$ >>_i] | i=1,...,|<<c,  $\eta$ ,  $\bar{\xi}$ ,  $\bar{\eta}$ >>| }
PrettyPrint[p] = Plus[PrettyPrintM[p_i] | i=1,...,|p|]
PrettyPrint[<m>] = PrettyPrintM[m]
PrettyPrint[<>] = 0

```

```

Definition ["IntDiffOp", any[ $\mathcal{F}$ , K],
IntDiffOp[ $\mathcal{F}$ , K] = where[A = FreeIntDiffOp[ $\mathcal{F}$ , K],
  QuotAlg[GBNF[GB[A], GreenSystem[ $\mathcal{F}$ , K]]]]
]

```

## Integro-Differential Polynomials

```

Definition["Differential Polynomials", any[ $\mathcal{F}$ , K],
DiffPolys[ $\mathcal{F}$ , K] = where [M = MonoidAlgebra[ $\mathcal{F}$ , TuplesMonoid[N]],
Functor[V, any[c, d, f, k, m, n, w, x, y,  $\lambda$ ,  $\xi$ ,  $\eta$ , A,  $\bar{\eta}$ ,  $\bar{m}$ ,  $\bar{x}$ ,  $\bar{y}$ ],

s = <>
-----
 $\xi_V[x] \Leftrightarrow \xi_M[x]$ 
-----
is-bvec[ $\xi$ ]  $\Leftrightarrow$  is-bvec[ $\xi$ ]
-----
is-coeff[c]  $\Leftrightarrow$  is-coeff[c]
-----
bvec[ $\xi$ ] = bvec[ $\xi$ ]
-----
to-exp[<>] = <>
-----
to-exp[<<c,  $\xi$ >,  $\bar{x}$ >] =  $\begin{cases} \left\langle \left\langle c \cdot \frac{1}{\mathcal{F}}, \xi \right\rangle \right\rangle_V + \text{to-exp}[\langle \bar{x} \rangle] & \Leftarrow \text{is-coeff}[\langle \langle c, \langle 0, 0 \rangle \rangle] \\ \left\langle \left\langle \langle c, \langle 0, 0 \rangle \rangle \cdot \frac{1}{\mathcal{F}}, \xi \right\rangle \right\rangle_V + \text{to-exp}[\langle \bar{x} \rangle] & \Leftarrow \text{otherwise} \end{cases}$ 
-----
0 = 0
V M
-----
 $\frac{1}{V} = \left\langle \left\langle \frac{1}{\mathcal{F}}, \langle \rangle \right\rangle \right\rangle$ 
-----
x + y = x + y
V M
-----
V[+] [x,  $\bar{x}$ ] = x + M[+] [ $\bar{x}$ ]
-----
 $\bar{V}x = \bar{M}x$ 
-----
x  $\bar{V}y = x \bar{M}y$ 
-----
c  $\bar{V}y$  (* scalar multiplication *) = c  $\bar{M}y$ 
-----
x * y = x * y
V M
-----
D[<>] = <> ]]]
-----
D[<math>\bar{m}, 0] = D[<math>\bar{m}]
-----
D[<k>] = <<k, <k-1, 1>>
-----
D[<math>\bar{m}, n] = D[<math>\bar{m}] *  $\left\langle \left\langle \frac{1}{K}, \left( 0 \mid_{i=1, \dots, |\langle \bar{m} \rangle|} , n \right) \right\rangle \right\rangle_{\text{MonoidAlgebra}[K, \text{TuplesMonoid}[N]]} + \left\langle \left\langle n, \langle \bar{m}, n-1, 1 \rangle \right\rangle \right\rangle_{\text{MonoidAlgebra}[K, \text{TuplesMonoid}[N]]}$ 
-----
 $\partial_V \langle \rangle = \langle \rangle$ 
-----
 $\partial_V \langle \langle c, \xi \rangle, \bar{x} \rangle = \begin{cases} \partial_{\mathcal{F}} c \cdot \left\langle \left\langle \frac{1}{\mathcal{F}}, \xi \right\rangle \right\rangle_V + c \cdot \text{to-exp}[\text{D}[\xi]]_V + \partial_V \langle \bar{x} \rangle & \Leftarrow \partial_{\mathcal{F}} c \neq 0 \\ c \cdot \text{to-exp}[\text{D}[\xi]]_V + \partial_V \langle \bar{x} \rangle & \Leftarrow \text{otherwise} \end{cases}$ 
-----
PrettyPrint[x] =  $\text{mPlus}[\text{PrettyPrintM}[x_i] \mid_{i=1, \dots, |x|}]$ 
-----
PrettyPrint[<math>\bar{m}>] = PrettyPrintM[m]
-----
PrettyPrint[<>] = 0
-----
PrettyPrintM[<c, w>] = tma-trfm[ $\text{mTimes}[\text{PrettyPrintC}[c], \text{PrettyPrintW}[w]]$ ]
-----
PrettyPrintC[c] = PrettyPrint[c]
V F
-----
PrettyPrintW[<>] = 1
-----
PrettyPrintW[<math>\eta, \bar{\eta}>] =  $\text{mPower}[u, \langle \eta, \bar{\eta} \rangle]$ 

```

```

Definition["Term Monoid for IDP", any[N,  $\mathcal{F}$ ],
TermMonoid[ $\mathcal{F}$ , N] =
CartesianProduct[TuplesMonoid[N], TuplesMonoid[N], TuplesMonoid[ $\mathcal{F} \times \text{TuplesMonoid}[N]$ ]]

```

```

Definition["Integro-Differential Polys", any[K,  $\mathcal{F}$ ],

```

```

IntDiffPolys[ $\mathcal{F}$ , K] = where [M = FreeModule[ $\mathcal{F}$ , TermMonoid[ $\mathcal{F}$ , N]],
  Functor[V, any[a, b, c, f, g, j, k, p, w,  $\alpha$ ,  $\beta$ ,  $\gamma$ ,  $\eta$ ,  $\bar{\eta}$ , A, J,  $\bar{b}$ ,  $\bar{\alpha}$ ,  $\bar{\beta}$ ,  $\bar{\gamma}$ ,  $\bar{J}$ ,  $\bar{a}$ ,  $\bar{\alpha}$ ,  $\bar{\gamma}$ ,  $\bar{J}$ ,  $\bar{\bar{J}}$ ],
    s = <>
     $\in_V[\langle\langle f, J, \bar{J} \rangle\rangle] \Leftrightarrow \left( \in_M[\langle\langle f, J, \bar{J} \rangle\rangle] \right)$ 
     $f >_V g \Leftrightarrow f >_M g$ 
     $\text{is-bvec}[\langle\langle c, f \rangle\rangle] \Leftrightarrow \left( \in_{\mathcal{F}}[c] \wedge \text{is-bvec}[f] \right)$ 
     $\text{is-coeff}[c] \Leftrightarrow \text{is-coeff}[c]$ 
    ew[f] = bvec[f]
     $\text{lcrd}[\langle\langle a, J \rangle\rangle, \langle\langle b, \bar{J} \rangle\rangle] = \left\langle \left\langle \text{lcrd}[a, b], \text{lcm}_{\text{TermMonoid}[\mathcal{F}, N]}[J, \bar{J}] \right\rangle \right\rangle$ 
     $\text{is-const}[\langle\langle c, \langle a, \langle \rangle, J \rangle\rangle] \Leftrightarrow \left( \alpha \neq_{\text{TuplesMonoid}[N]} \bigvee J \neq_{\text{TuplesMonoid}[\mathcal{F} \times \text{TuplesMonoid}[N]]} \right)$ 
     $\text{is-const}[A] \Leftrightarrow \text{False}$ 
     $\text{is-scal}[\langle\langle c, \text{TermMonoid}[\mathcal{F}, N] \rangle\rangle] \Leftrightarrow \text{is-coeff}[c]$ 
     $\text{is-scal}[\langle\langle c, \langle \rangle \rangle] \Leftrightarrow \text{is-coeff}[c]$ 
     $\text{is-scal}[a] \Leftrightarrow \text{False}$ 
     $\text{to-scal}[\langle \rangle] = 0_K$ 
     $\text{to-scal}[\langle\langle c, \text{TermMonoid}[\mathcal{F}, N] \rangle\rangle] = c$ 
     $\text{to-scal}[\langle\langle c, \langle \rangle \rangle] = c$ 
    0 = <>
     $\frac{1}{V} = \left\langle \left\langle \frac{1}{\mathcal{F}}, \langle \rangle, \langle \rangle, \langle \rangle \right\rangle \right\rangle$ 
     $\bar{v}f = \frac{f}{M}$ 
     $f +_V g = f +_M g$ 
     $f \bar{v} g = f \bar{v} g$ 
     $\langle \rangle \cdot_V f = \langle \rangle$ 
     $c \bar{v} f = c \cdot_M f$ 
     $\text{eval}[\langle\langle f, \langle a, \beta, \langle \rangle \rangle \rangle, 0] = \begin{cases} \left\langle \left\langle \text{eval}[f, 0] \cdot \frac{1}{\mathcal{F}}, \left\langle \alpha \cdot_{\text{TuplesMonoid}[N]} \beta, \langle \rangle, \langle \rangle \right\rangle \right\rangle \right\rangle & \Leftrightarrow \text{is-coeff}[\text{eval}[f, 0]] \\ \left\langle \left\langle \text{eval}[f, 0], \left\langle \alpha \cdot_{\text{TuplesMonoid}[N]} \beta, \langle \rangle, \langle \rangle \right\rangle \right\rangle \right\rangle & \Leftrightarrow \text{eval}[f, 0] \neq 0_K \\ \langle \rangle & \Leftrightarrow \text{otherwise} \end{cases}$ 
     $\text{eval}[\langle\langle f, \langle a, \beta, J \rangle \rangle, 0] = \langle \rangle$ 
     $\text{to-idp}[\langle \rangle] = \langle \rangle$ 
     $\text{to-idp}[\langle\langle c, J, \bar{J} \rangle\rangle] = \langle\langle c, \langle \rangle, \langle \rangle, J \rangle\rangle \bar{v} \text{to-idp}[\langle\bar{J} \rangle]$ 
     $\text{dp-to-idp}[\langle \rangle] = \langle \rangle$ 
     $\text{dp-to-idp}[\langle\langle b, \gamma, \bar{J} \rangle\rangle] = \langle\langle b, \langle \rangle, \gamma, \langle \rangle \rangle \bar{v} \text{dp-to-idp}[\langle\bar{J} \rangle]$ 
     $\text{ls-to-idp}[\langle \rangle] = \langle \rangle$ 
     $\text{ls-to-idp}[\langle\langle J, \bar{J} \rangle\rangle] = \text{to-idp}[J] \bar{v} \text{ls-to-idp}[\langle\bar{J} \rangle]$ 
    exp[ $\langle \rangle$ ] = <>
     $\text{exp}[a] = \left\langle \left\langle \frac{1}{\mathcal{F}}, a \right\rangle \right\rangle$ 
     $a \times_V \langle \rangle = \langle \rangle$ 
     $a \times_V \langle\langle c, \langle a, \bar{a} \rangle \rangle, \bar{a} \rangle = \left\langle \left\langle c, \langle a, \alpha, \bar{a} \rangle \right\rangle \right\rangle \text{FreeModule}[\mathcal{F}, \text{FreeModule}[\mathcal{F}, \text{TuplesMonoid}[N]]] \bar{v} a \times_V \langle \bar{a} \rangle$ 
     $a \circ_V \langle \rangle = \langle \rangle$ 
     $\langle a \rangle \circ_V \langle\langle c, \langle a, \beta, \langle \rangle \rangle \rangle, \bar{a} \rangle = \langle\langle c, \langle a, \beta, \langle a \rangle \rangle \rangle \bar{v} \langle a \rangle \circ_V \langle \bar{a} \rangle$ 

```

$$\langle \mathbf{a} \rangle_{\mathbb{V}} \circ \langle \langle \mathbf{c}, \langle \alpha, \beta, \langle \langle \mathbf{b}, \gamma \rangle, \bar{\mathbf{J}} \rangle \rangle, \bar{\mathbf{a}} \rangle = \langle \langle \mathbf{c}, \langle \alpha, \beta, \langle \mathbf{a}, \langle \mathbf{b}, \gamma \rangle, \bar{\mathbf{J}} \rangle \rangle \rangle + \langle \mathbf{a} \rangle_{\mathbb{V}} \langle \bar{\mathbf{a}} \rangle$$

$$\langle \rangle_{\mathbb{V}} \circ \alpha = \exp_{\mathbb{V}}[\alpha]$$

$$\mathbf{a} \circ \langle \rangle = \exp_{\mathbb{V}}[\mathbf{a}]$$

$$\langle \mathbf{a}, \bar{\mathbf{J}} \rangle_{\mathbb{V}} \circ \langle \alpha, \bar{\mathbf{J}} \rangle (* \text{ shuffle product } *) = \mathbf{a} \times_{\mathbb{V}} \left( \langle \bar{\mathbf{J}} \rangle_{\mathbb{V}} \circ \langle \alpha, \bar{\mathbf{J}} \rangle \right) + \alpha \times_{\mathbb{V}} \left( \langle \bar{\mathbf{J}} \rangle_{\mathbb{V}} \circ \langle \mathbf{a}, \bar{\mathbf{J}} \rangle \right)$$

$$\langle \rangle * \bar{\mathbf{J}} = \langle \rangle$$

$$\bar{\mathbf{J}} * \langle \rangle = \langle \rangle$$

$$\langle \langle \mathbf{b}, \langle \alpha, \beta, \langle \rangle \rangle \rangle, \bar{\mathbf{J}} \rangle * \langle \langle \tilde{\mathbf{b}}, \langle \tilde{\alpha}, \tilde{\beta}, \tilde{\mathbf{J}} \rangle \rangle, \tilde{\mathbf{J}} \rangle = \left( \left\langle \left\langle \mathbf{b} *_{\mathcal{F}} \tilde{\mathbf{b}}, \left\langle \alpha *_{\text{TuplesMonoid}[\mathcal{N}]} \tilde{\alpha}, \beta *_{\text{TuplesMonoid}[\mathcal{N}]} \tilde{\beta}, \tilde{\mathbf{J}} \right\rangle \right\rangle \right\rangle + \langle \langle \mathbf{b}, \langle \alpha, \beta, \langle \rangle \rangle \rangle *_{\mathbb{V}} \langle \tilde{\mathbf{J}} \rangle + \langle \bar{\mathbf{J}} \rangle *_{\mathbb{V}} \langle \langle \tilde{\mathbf{b}}, \langle \tilde{\alpha}, \tilde{\beta}, \tilde{\mathbf{J}} \rangle \rangle, \tilde{\mathbf{J}} \rangle \right)$$

$$\langle \langle \mathbf{b}, \langle \alpha, \beta, \bar{\mathbf{J}} \rangle \rangle, \bar{\mathbf{J}} \rangle * \langle \langle \tilde{\mathbf{b}}, \langle \tilde{\alpha}, \tilde{\beta}, \tilde{\mathbf{J}} \rangle \rangle, \tilde{\mathbf{J}} \rangle = \left( \left\langle \left\langle \langle \mathbf{b}, \langle \alpha, \beta, \langle \rangle \rangle \rangle *_{\mathbb{V}} \langle \langle \tilde{\mathbf{b}}, \langle \tilde{\alpha}, \tilde{\beta}, \langle \rangle \rangle \rangle \right\rangle \right\rangle *_{\mathbb{V}} \text{to-idp}[\bar{\mathbf{J}} \circ \tilde{\mathbf{J}}] + \langle \langle \mathbf{b}, \langle \alpha, \beta, \bar{\mathbf{J}} \rangle \rangle *_{\mathbb{V}} \langle \tilde{\mathbf{J}} \rangle + \langle \bar{\mathbf{J}} \rangle *_{\mathbb{V}} \langle \langle \tilde{\mathbf{b}}, \langle \tilde{\alpha}, \tilde{\beta}, \tilde{\mathbf{J}} \rangle \rangle, \tilde{\mathbf{J}} \rangle \right)$$

$$\exp\text{-basis}[\langle \rangle] = \langle \rangle$$

$$\exp\text{-basis}[\langle \langle \mathbf{c}, \langle \alpha, \beta, \langle \rangle \rangle \rangle, \bar{\mathbf{J}} \rangle] = \langle \langle \mathbf{c}, \langle \alpha, \beta, \langle \rangle \rangle \rangle + \exp\text{-basis}[\langle \bar{\mathbf{J}} \rangle]$$

$$\exp\text{-basis}[\langle \langle \mathbf{c}, \langle \alpha, \beta, \langle \langle \langle \rangle, \gamma \rangle \rangle \rangle, \bar{\mathbf{J}} \rangle \rangle] = \langle \rangle$$

$$\exp\text{-basis}[\langle \langle \mathbf{c}, \langle \alpha, \beta, \langle \langle \mathbf{f}, \langle 0, 1 \rangle \rangle \rangle, \bar{\mathbf{J}} \rangle \rangle \rangle] =$$

$$\mathbf{c} \cdot \exp\text{-basis}[\langle \langle \frac{1}{\mathcal{F}}, \langle \alpha, \beta, \langle \bar{\mathbf{J}} \rangle \rangle \rangle \rangle] *_{\mathbb{V}} \left( \langle \langle \langle \mathbf{f}, \langle \langle \rangle, \langle 1 \rangle \rangle \rangle \rangle \rangle_{\mathbb{V}} \int_{\mathbb{V}} \langle \langle \partial_{\mathcal{F}} \mathbf{f}, \langle \langle \rangle, \langle 1 \rangle \rangle \rangle \rangle \rangle_{\mathbb{V}} \right)$$

$$\text{eval}_{\mathbb{V}}[\langle \langle \mathbf{f}, \langle \langle \rangle, \langle 1 \rangle \rangle \rangle, 0 \rangle]$$

$$\exp\text{-basis}[\langle \langle \mathbf{c}, \langle \alpha, \beta, \langle \langle \mathbf{f}, \langle \bar{\alpha}, \mathbf{k}, 1 \rangle \rangle \rangle, \bar{\mathbf{J}} \rangle \rangle \rangle] =$$

$$\left( \frac{1}{\mathbf{k}} (\mathbf{k} + 1) \cdot \mathbf{c} \right) \cdot \exp\text{-basis}[\langle \langle \frac{1}{\mathcal{F}}, \langle \alpha, \beta, \langle \bar{\mathbf{J}} \rangle \rangle \rangle \rangle] *_{\mathbb{V}} \left( \left\langle \left\langle \langle \mathbf{f}, \langle \langle \rangle, \langle \bar{\alpha}, \mathbf{k} + 1 \rangle \rangle \rangle \right\rangle \right\rangle_{\mathbb{V}} \int_{\mathbb{V}} \left\langle \partial_{\mathbb{V}} \left\langle \frac{1}{\mathcal{F}}, \langle \langle \rangle, \langle \bar{\alpha} \rangle, \langle \rangle \right\rangle \right\rangle *_{\mathbb{V}} \right)$$

$$\left\langle \left\langle \frac{1}{\mathcal{F}}, \langle \langle \rangle, \left\langle 0 \begin{array}{c} | \\ i=0, \dots, |\bar{\alpha}|-1 \end{array}, \mathbf{k} + 1 \right\rangle, \langle \rangle \right\rangle \right\rangle \right)$$

$$\bar{\mathbb{V}} \text{eval}_{\mathbb{V}}[\langle \langle \frac{1}{\mathcal{F}}, \langle \langle \rangle, \langle \bar{\alpha}, \mathbf{k} + 1 \rangle \rangle \rangle, 0 \rangle]$$

$$\exp\text{-basis}[\langle \langle \mathbf{c}, \langle \alpha, \beta, \langle \langle \langle \mathbf{b}, \mathbf{a} \rangle \rangle, \gamma \rangle \rangle \rangle \rangle] = \langle \langle \mathbf{c} \cdot \mathbf{b}, \langle \alpha, \beta, \langle \langle \langle \frac{1}{\mathbf{k}}, \mathbf{a} \rangle \rangle, \gamma \rangle \rangle \rangle \rangle$$

$$\exp\text{-basis}[\langle \langle \mathbf{c}, \langle \alpha, \beta, \langle \langle \langle \mathbf{b}, \mathbf{a} \rangle, \bar{\mathbf{a}} \rangle, \gamma \rangle, \bar{\mathbf{J}} \rangle \rangle, \tilde{\mathbf{J}} \rangle \rangle] = \left( \left\langle \left\langle \left\langle \frac{1}{\mathbf{k}}, \mathbf{a} \right\rangle \right\rangle, \gamma \right\rangle \right) \circ \exp\text{-basis}[\langle \langle \mathbf{c} \cdot \mathbf{b}, \langle \alpha, \beta, \langle \bar{\mathbf{J}} \rangle \rangle \rangle \rangle] + \exp\text{-basis}[\langle \langle \mathbf{c}, \langle \alpha, \beta, \langle \langle \langle \bar{\mathbf{a}} \rangle, \gamma \rangle \rangle, \bar{\mathbf{J}} \rangle \rangle \rangle] + \exp\text{-basis}[\langle \tilde{\mathbf{J}} \rangle]$$

$$\partial_{\mathbb{V}} \langle \rangle = \langle \rangle$$

$$\partial_{\mathbb{V}} \langle \mathbf{f}, \langle \langle \rangle, \langle \rangle, \langle \rangle \rangle = \begin{cases} \langle \rangle & \Leftarrow \text{is-scal}[\mathbf{f}] \\ \partial_{\mathcal{F}} \mathbf{f} \cdot \frac{1}{\mathbb{V}} & \Leftarrow \text{otherwise} \end{cases}$$

$$\partial_{\mathbb{V}} \langle \mathbf{f}, \langle \alpha, \beta, \langle \rangle \rangle = \begin{cases} \text{dp-to-idp}[\partial_{\text{DiffPolys}[\mathcal{F}, \mathbf{K}]}[\langle \frac{1}{\mathcal{F}}, \beta \rangle]] *_{\mathbb{V}} \langle \langle \mathbf{f}, \langle \alpha, \langle \rangle, \langle \rangle \rangle \rangle & \Leftarrow \text{is-scal}[\mathbf{f}] \\ \partial_{\mathcal{F}} \mathbf{f} \cdot \langle \frac{1}{\mathcal{F}}, \langle \alpha, \beta, \langle \rangle \rangle \rangle + \text{dp-to-idp}[\partial_{\text{DiffPolys}[\mathcal{F}, \mathbf{K}]}[\langle \frac{1}{\mathcal{F}}, \beta \rangle]] *_{\mathbb{V}} \langle \langle \mathbf{f}, \langle \alpha, \langle \rangle, \langle \rangle \rangle \rangle & \Leftarrow \text{otherwise} \end{cases}$$

$$\partial_{\mathbb{V}} \langle \mathbf{f}, \langle \alpha, \beta, \langle \langle \mathbf{b}, \gamma \rangle, \bar{\mathbf{J}} \rangle \rangle = \begin{cases} \text{dp-to-idp}[\partial_{\text{DiffPolys}[\mathcal{F}, \mathbf{K}]}[\langle \frac{1}{\mathcal{F}}, \beta \rangle]] *_{\mathbb{V}} \langle \langle \mathbf{f}, \langle \alpha, \langle \rangle, \langle \langle \mathbf{b}, \gamma \rangle, \bar{\mathbf{J}} \rangle \rangle \rangle + & \Leftarrow \text{is-scal}[\mathbf{f}] \\ \left\langle \left\langle \mathbf{f} *_{\mathcal{F}} \mathbf{b}, \left\langle \alpha, \beta *_{\text{TuplesMonoid}[\mathcal{N}]} \gamma, \langle \bar{\mathbf{J}} \rangle \right\rangle \right\rangle \right\rangle & \\ \left( \langle \langle \partial_{\mathcal{F}} \mathbf{f}, \langle \alpha, \beta, \langle \langle \mathbf{b}, \gamma \rangle, \bar{\mathbf{J}} \rangle \rangle \rangle \rangle + \text{dp-to-idp}[\partial_{\text{DiffPolys}[\mathcal{F}, \mathbf{K}]}[\langle \frac{1}{\mathcal{F}}, \beta \rangle]] *_{\mathbb{V}} & \Leftarrow \text{otherwise} \\ \langle \langle \mathbf{f}, \langle \alpha, \langle \rangle, \langle \langle \mathbf{b}, \gamma \rangle, \bar{\mathbf{J}} \rangle \rangle \rangle \rangle + \left\langle \left\langle \mathbf{f} *_{\mathcal{F}} \mathbf{b}, \left\langle \alpha, \beta *_{\text{TuplesMonoid}[\mathcal{N}]} \gamma, \langle \bar{\mathbf{J}} \rangle \right\rangle \right\rangle \right\rangle & \end{cases}$$

$$\partial_{\mathbb{V}} \langle \langle \mathbf{f}, \langle \alpha, \beta, \bar{\mathbf{J}} \rangle \rangle, \tilde{\mathbf{J}} \rangle = \partial_{\mathbb{V}} \langle \mathbf{f}, \langle \alpha, \beta, \bar{\mathbf{J}} \rangle \rangle + \partial_{\mathbb{V}} \langle \tilde{\mathbf{J}} \rangle$$

$$\int_{\mathbb{V}} \langle \rangle = \langle \rangle$$

$$\int_V \langle \langle \mathbf{f}, \langle \alpha, \langle \rangle, \langle \rangle \rangle \rangle \rangle = \left( \int_{\mathcal{F}} \mathbf{f} \right) \dot{\vee} \left\langle \left\langle \frac{1}{\mathcal{F}}, \langle \alpha, \langle \rangle, \langle \rangle \right\rangle \right\rangle$$

$$\int_V \langle \langle \mathbf{f}, \langle \alpha, \langle \rangle, \langle \langle \mathbf{b}, \gamma \rangle, \bar{\mathcal{J}} \rangle \rangle \rangle = (* \text{ quasiconstant } *) \left( \int_{\mathcal{F}} \mathbf{f} \right) \dot{\vee} \left\langle \left\langle \frac{1}{\mathcal{F}}, \langle \alpha, \langle \rangle, \langle \langle \mathbf{b}, \gamma \rangle, \bar{\mathcal{J}} \rangle \right\rangle \right\rangle \bar{\vee}$$

$$\left( \int_V \left\langle \left\langle \mathbf{b} * \int_{\mathcal{F}} \mathbf{f}, \langle \alpha, \gamma, \langle \bar{\mathcal{J}} \rangle \right\rangle \right\rangle \right)$$

$$\int_V \langle \langle \mathbf{f}, \langle \alpha, \langle 0, 1 \rangle, \langle \rangle \rangle \rangle = (* \text{ quasilinear base case } *)$$

$$\left( \langle \langle \mathbf{f}, \langle \alpha, \langle 1 \rangle, \langle \rangle \rangle \rangle \right) \bar{\vee} \int_V \left( \partial_{\mathcal{F}} \mathbf{f} \right) \dot{\vee} \left\langle \left\langle \frac{1}{\mathcal{F}}, \langle \alpha, \langle 1 \rangle, \langle \rangle \right\rangle \right\rangle \right) \bar{\vee} \text{eval}_V[\langle \langle \mathbf{f}, \langle \alpha, \langle 1 \rangle, \langle \rangle \rangle \rangle, 0]$$

$$\int_V \langle \langle \mathbf{f}, \langle \alpha, \langle 0, 1 \rangle, \langle \langle \mathbf{b}, \gamma \rangle, \bar{\mathcal{J}} \rangle \rangle \rangle = (* \text{ quasilinear base case } *)$$

$$\left( \langle \langle \mathbf{f}, \langle \alpha, \langle 1 \rangle, \langle \langle \mathbf{b}, \gamma \rangle, \bar{\mathcal{J}} \rangle \rangle \rangle \right) \bar{\vee} \int_V \left( \partial_{\mathcal{F}} \mathbf{f} \right) \dot{\vee} \left\langle \left\langle \frac{1}{\mathcal{F}}, \langle \alpha, \langle 1 \rangle, \langle \langle \mathbf{b}, \gamma \rangle, \bar{\mathcal{J}} \rangle \right\rangle \right\rangle \right) \bar{\vee} \int_V \left\langle \left\langle \mathbf{f} * \mathbf{b}, \langle \alpha, \langle 1 \rangle \right. \right. \left. \left. \begin{matrix} * \\ \text{TuplesMonoid}[N] \end{matrix} \right. \right. \left. \left. \gamma, \langle \bar{\mathcal{J}} \rangle \right\rangle \right\rangle$$

$$\int_V \langle \langle \mathbf{f}, \langle \alpha, \langle \bar{\alpha}, k, 1 \rangle, \langle \rangle \rangle \rangle = (* \text{ quasilinear } *)$$

$$\left( \frac{1}{k} \int_{\mathcal{F}} \left( \frac{1}{k+1} \right) \dot{\vee} \left( \langle \langle \mathbf{f}, \langle \alpha, \langle \bar{\alpha}, k+1 \rangle, \langle \rangle \rangle \rangle \right) \bar{\vee} \int_V \left( \partial_V \langle \langle \mathbf{f}, \langle \alpha, \langle \bar{\alpha} \rangle, \langle \rangle \rangle \rangle * \left\langle \left\langle \frac{1}{\mathcal{F}}, \langle \rangle, \left\langle 0 \begin{matrix} | \\ i=0, \dots, \langle \bar{\alpha} \rangle - 1 \end{matrix} \right\rangle, k+1 \right\rangle, \langle \rangle \right\rangle \right) \right)$$

$$\text{eval}_V[\langle \langle \mathbf{f}, \langle \alpha, \langle \bar{\alpha}, k+1 \rangle, \langle \rangle \rangle \rangle, 0]$$

$$\int_V \langle \langle \mathbf{f}, \langle \alpha, \langle \bar{\alpha}, k, 1 \rangle, \bar{\mathcal{J}} \rangle \rangle \rangle = (* \text{ quasilinear } *)$$

$$\left( \frac{1}{k} \int_{\mathcal{F}} \left( \frac{1}{k+1} \right) \dot{\vee} \left( \langle \langle \mathbf{f}, \langle \alpha, \langle \bar{\alpha}, k+1 \rangle, \bar{\mathcal{J}} \rangle \rangle \right) \bar{\vee} \int_V \left( \partial_V \langle \langle \mathbf{f}, \langle \alpha, \langle \bar{\alpha} \rangle, \bar{\mathcal{J}} \rangle \rangle * \left\langle \left\langle \frac{1}{\mathcal{F}}, \langle \rangle, \left\langle 0 \begin{matrix} | \\ i=0, \dots, \langle \bar{\alpha} \rangle - 1 \end{matrix} \right\rangle, k+1 \right\rangle, \langle \rangle \right\rangle \right) \right)$$

$$\int_V \langle \langle \mathbf{f}, \langle \alpha, \beta, \langle \bar{\mathcal{J}} \rangle \rangle \rangle = (* \text{ functional } *) \left\langle \left\langle \frac{1}{\mathcal{F}}, \langle \alpha, \langle \rangle, \langle \langle \mathbf{f}, \beta \rangle, \bar{\mathcal{J}} \rangle \right\rangle \right\rangle$$

$$\int_V \langle \bar{\mathcal{J}} \rangle = \int_V \langle \bar{\mathcal{J}} \rangle + \int_V \langle \bar{\mathcal{J}} \rangle$$

$$\text{PrettyPrint}_V[\mathbf{p}] = \text{mPlus} \left[ \text{PrettyPrintM}[\mathbf{p}_i] \begin{matrix} | \\ i=1, \dots, |\mathbf{p}| \end{matrix} \right]$$

$$\text{PrettyPrintV}[\mathbf{p}] = \text{mPlus} \left[ \text{PrettyPrintMV}[\mathbf{p}_i] \begin{matrix} | \\ i=1, \dots, |\mathbf{p}| \end{matrix} \right]$$

$$\text{PrettyPrint}[\langle \mathbf{p} \rangle] = \text{PrettyPrintM}[\mathbf{p}]$$

$$\text{PrettyPrintV}[\langle \mathbf{p} \rangle] = \text{PrettyPrintMV}[\mathbf{p}]$$

$$\text{PrettyPrint}[\langle \rangle] = 0$$

$$\text{PrettyPrintV}[\langle \rangle] = 0$$

$$\text{PrettyPrintM}[\langle \mathbf{c}, \mathbf{w} \rangle] = \text{tma-trfm} \left[ \text{mTimes} \left[ \text{PrettyPrint}[\mathbf{c}], \text{PrettyPrintW}[\mathbf{w}] \right] \right]$$

$$\text{PrettyPrintMV}[\langle \mathbf{c}, \mathbf{w} \rangle] = \text{tma-trfm} \left[ \text{mTimes} \left[ \text{PrettyPrint}[\mathbf{c}], \text{PrettyPrintWV}[\mathbf{w}] \right] \right]$$

$$\text{PrettyPrintW}[\langle \rangle] = 1$$

$$\text{PrettyPrintWV}[\langle \rangle] = 1$$

$$\text{PrettyPrintW}[\langle \langle \rangle, \langle \rangle, \bar{\mathcal{J}} \rangle] = \text{PrettyPrintL}[\bar{\mathcal{J}}]$$

$$\text{PrettyPrintWV}[\langle \langle \rangle, \langle \rangle, \bar{\mathcal{J}} \rangle] = \text{PrettyPrintLV}[\bar{\mathcal{J}}]$$

$$\text{PrettyPrintW}[\langle \langle \rangle, \beta, \bar{\mathcal{J}} \rangle] = \text{mytimes} \left[ \text{mPower}[\mathbf{u}, \beta], \text{PrettyPrintL}[\bar{\mathcal{J}}] \right]$$

$$\text{PrettyPrintWV}[\langle \langle \rangle, \beta, \bar{\mathcal{J}} \rangle] = \text{mytimes} \left[ \text{mPower}[\mathbf{v}, \beta], \text{PrettyPrintLV}[\bar{\mathcal{J}}] \right]$$

$$\text{PrettyPrintW}[\langle \alpha, \langle \rangle, \bar{\mathcal{J}} \rangle] = \text{mytimes} \left[ \text{mPower}[\mathbf{u}[0], \alpha], \text{PrettyPrintL}[\bar{\mathcal{J}}] \right]$$

$$\text{PrettyPrintWV}[\langle \alpha, \langle \rangle, \bar{\mathcal{J}} \rangle] = \text{mytimes} \left[ \text{mPower}[\mathbf{v}[0], \alpha], \text{PrettyPrintLV}[\bar{\mathcal{J}}] \right]$$

$$\text{PrettyPrintW}[\langle \alpha, \beta, \bar{\mathcal{J}} \rangle] = \text{mytimes} \left[ \text{mPower}[\mathbf{u}[0], \alpha], \text{mPower}[\mathbf{u}, \beta], \text{PrettyPrintL}[\bar{\mathcal{J}}] \right]$$

$$\text{PrettyPrintWV}[\langle \alpha, \beta, \bar{\mathcal{J}} \rangle] = \text{mytimes} \left[ \text{mPower}[\mathbf{v}[0], \alpha], \text{mPower}[\mathbf{v}, \beta], \text{PrettyPrintLV}[\bar{\mathcal{J}}] \right]$$

$$\text{PrettyPrintL}[\langle \rangle] = ""$$

$$\text{PrettyPrintLV}[\langle \rangle] = ""$$

$$\text{PrettyPrintL}[\langle \langle \mathbf{b}, \gamma \rangle, \bar{\mathcal{J}} \rangle] = \text{mytimes} \left[ \text{mytimes} \left[ \text{" "}, \text{PrettyPrint}[\mathbf{b}] \right], \text{mPower}[\mathbf{u}, \gamma], \text{PrettyPrintL}[\langle \bar{\mathcal{J}} \rangle] \right]$$

$$\text{PrettyPrintLV}[\langle \langle \mathbf{b}, \gamma \rangle, \bar{\mathcal{J}} \rangle] = \text{mytimes} \left[ \text{mytimes} \left[ \text{" "}, \text{PrettyPrint}[\mathbf{b}] \right], \text{mPower}[\mathbf{v}, \gamma], \text{PrettyPrintLV}[\langle \bar{\mathcal{J}} \rangle] \right]$$

```

Definition ["PrettyPrinting", any[ $\mathcal{F}$ , K],
PrettyPrinting[ $\mathcal{F}$ , K] = Functor[P, any[b, c, m, n, p, r, w,  $\alpha$ ,  $\beta$ ,  $\lambda$ ,  $\gamma$ ,  $\eta$ , J, L,  $\bar{\eta}$ ,  $\bar{\xi}$ ,  $\bar{J}$ ],

s = <>
PrettyPrintW[<>] = ""
PrettyPrintW[<math>\lambda, \eta, \bar{\eta}>] = mytimes[PrettyPrintL[ $\lambda$ ], PrettyPrintW[<math>\eta, \bar{\eta}>]]
PrettyPrintW[<math>\eta>] = PrettyPrintL[ $\eta$ ]
PrettyPrintF[<<>, <>, <>>] = ""
PrettyPrintF[<math>\alpha, <>, <>>] = mPower[v[0],  $\alpha$ ]
PrettyPrintF[<<>,  $\beta$ , <>>] = mytimes[mPower[v,  $\beta$ ]]
PrettyPrintF[<math>\alpha, \beta, <>>] = mytimes[mPower[v[0],  $\alpha$ ], mPower[v,  $\beta$ ]]
PrettyPrintF[<<>, <>, J] = PrettyPrintJ[J]
PrettyPrintF[<math>\alpha, \beta, J>] = mytimes[mPower[v[0],  $\alpha$ ], mPower[v,  $\beta$ ], PrettyPrintJ[J]]
PrettyPrintL[<math>\lceil \rceil, <<c,  $\alpha, \beta, J$ >>] = mytimes[IntDiffPolys[Exp,K] PrettyPrint[c], P PrettyPrintF[<math>\alpha, \beta, J>]]
PrettyPrintL[<math>\lfloor \rfloor, 0] = "E"
PrettyPrintJ[<>] = ""
PrettyPrintJ[<<math>\langle b, \gamma, \bar{J} \rangle>] = mytimes[mytimes[f, IntDiffPolys[Exp,K] PrettyPrintV[b]], mPower[v,  $\gamma$ ], P PrettyPrintJ[<math>\langle \bar{J} \rangle]]
PrettyPrintL[" $\partial$ "] = " $\partial$ "
PrettyPrintL[" $\int$ "] = " $\int$ "
PrettyPrintL[<>] = ""
PrettyPrintL[ $\eta$ ] =  $\eta$ 
PrettyPrintM[<>] = 0
PrettyPrintM[<c, <>>] = c
PrettyPrintM[<c, w>] = mytimes[K PrettyPrint[c], P PrettyPrintW[w]]
PrettyPrintRule[<m, n>] = P PrettyPrint[m]  $\rightarrow$  P PrettyPrint[n]
PrettyPrintRules[L] = <math>\langle P PrettyPrintRule[Li] | i=1, \dots, |L| >
PrettyPrint[<<<math>\langle c, \eta, \bar{\xi} \rangle, \bar{\eta} \rangle>] = <math>\langle P PrettyPrint[<<<math>\langle c, \eta, \bar{\xi} \rangle, \bar{\eta} \rangle> | i=1, \dots, |<<<math>\langle c, \eta, \bar{\xi} \rangle, \bar{\eta} \rangle| >
PrettyPrint[p] = mPlus[P PrettyPrintM[pi] | i=1, \dots, |p|]
PrettyPrint[<m>] = P PrettyPrintM[m]
PrettyPrint[<>] = 0
]]

```

Translator IDP

Translator Exponential Polynomials

```

Use[Built-in["Tuples"], Built-in["Quantifiers"],
Built-in["Connectives"], Built-in["Mult"], Built-in["Minus"], Built-in["IsNumber"]]

```

```

Definition["Parsing Exp", any[ $\mathcal{F}$ ],
  Parsing[ $\mathcal{F}$ ] = Functor[P, any[c, m, n, r,  $\lambda$ ,  $\bar{m}$ ,  $\bar{n}$ ],
    s = <>
    ParsePol[] = <>
    ParsePol[ $\mathbb{T}^m$ Tuple[ $\bar{m}$ ]] =  $\left( \text{ParsePol}[\langle \bar{m} \rangle_i] \mid_{i=1, \dots, |\langle \bar{m} \rangle|} \right)$ 
    ParsePol[ $\mathbb{T}^m$ Plus[]] = <>
    ParsePol[ $\mathbb{T}^m$ Plus[m,  $\bar{n}$ ]] =  $\text{ParsePol}[m] \times \text{ParsePol}[\mathbb{T}^m$ Plus[ $\bar{n}$ ]]
    ParsePol[ $\mathbb{T}^m$ Times[ $\mathbb{T}^m$ Minus[c],  $\bar{m}$ ]] =  $\text{ParseTerm}[\mathbb{T}^m$ Times[ $\mathbb{T}^m$ Minus[c],  $\bar{m}$ ]]
    ParsePol[ $\mathbb{T}^m$ Times[c,  $\bar{m}$ ]] =  $\text{ParseTerm}[\mathbb{T}^m$ Times[c,  $\bar{m}$ ]]
    ParsePol[ $\mathbb{T}^m$ Minus[ $\mathbb{T}^m$ Times[ $\mathbb{T}^m$ Minus[c],  $\bar{m}$ ]]] =  $\text{ParseTerm}[\mathbb{T}^m$ Times[c,  $\bar{m}$ ]]
    ParsePol[ $\mathbb{T}^m$ Minus[ $\mathbb{T}^m$ Times[c,  $\bar{m}$ ]]] =  $\text{ParseTerm}[\mathbb{T}^m$ Times[ $\mathbb{T}^m$ Minus[c],  $\bar{m}$ ]]
    ParsePol[ $\mathbb{T}^m$ Minus[m, n]] =  $\text{ParsePol}[m] \times \text{ParsePol}[\mathbb{T}^m$ Minus[n]]
    ParsePol[ $\mathbb{T}^m$ Minus[m, n, r,  $\bar{m}$ ]] =  $\text{ParsePol}[m] \times \text{ParsePol}[\mathbb{T}^m$ Minus[ $\mathbb{T}^m$ Minus[n], r,  $\bar{m}$ ]]
    ParsePol[ $\mathbb{T}^m$ Minus[c]] =  $\begin{cases} \langle \langle -c, \langle \rangle \rangle \rangle & \Leftarrow \mathbb{T}^m$ IsNumber[c] \\ \text{ParseTerm}[\mathbb{T}^mMinus[c]] & \Leftarrow \text{otherwise} \end{cases}
    ParsePol[c] =  $\begin{cases} \langle \langle c, \langle \rangle \rangle \rangle & \Leftarrow \mathbb{T}^m$ IsNumber[c] \\ \text{ParseTerm}[c] & \Leftarrow \text{otherwise} \end{cases}
    ParseTerm[ $\mathbb{T}^m$ Minus[c]] =  $\begin{cases} \langle \langle -c, \langle \rangle \rangle \rangle & \Leftarrow \mathbb{T}^m$ IsNumber[c] \\ -1 \cdot \text{ParseMon}[c] & \Leftarrow \text{otherwise} \end{cases}
    ParseTerm[ $\mathbb{T}^m$ Times[ $\mathbb{T}^m$ Times[ $\bar{m}$ ]]] =  $\text{ParseTerm}[\mathbb{T}^m$ Times[ $\bar{m}$ ]]
    ParseTerm[ $\mathbb{T}^m$ Times[ $\mathbb{T}^m$ Minus[c],  $\bar{m}$ ]] =  $\begin{cases} -c \cdot \text{ParseTerm}[\mathbb{T}^m$ Times[ $\bar{m}$ ]] & \Leftarrow \mathbb{T}^mIsNumber[c] \\ -1 \cdot \text{ParseTerm}[\mathbb{T}^mTimes[c,  $\bar{m}$ ]] & \Leftarrow \text{otherwise} \end{cases}
    ParseTerm[ $\mathbb{T}^m$ Times[c,  $\bar{m}$ ]] =  $\begin{cases} c \cdot \text{ParseTerm}[\mathbb{T}^m$ Times[ $\bar{m}$ ]] & \Leftarrow \mathbb{T}^mIsNumber[c] \\ \text{ParseMon}[\mathbb{T}^mTimes[c,  $\bar{m}$ ]] & \Leftarrow \text{otherwise} \end{cases}
    ParseTerm[c] =  $\begin{cases} \langle \langle c, \langle \rangle \rangle \rangle & \Leftarrow \mathbb{T}^m$ IsNumber[c] \\ \text{ParseMon}[c] & \Leftarrow \text{otherwise} \end{cases}
    ParseTerm[ $\mathbb{T}^m$ Times[]] = <>
    ParseMon[ $\mathbb{T}^m$ Times[ $\mathbb{T}^m$ Power[x, n],  $\mathbb{T}^m$ Power[E,  $\mathbb{T}^m$ Times[ $\lambda$ , x]]] =  $\langle \langle 1, \langle \langle "[]" \rangle, \langle n, \lambda \rangle \rangle \rangle$ 
    ParseMon[ $\mathbb{T}^m$ Times[ $\mathbb{T}^m$ Times[ $\mathbb{T}^m$ Power[x, n],  $\mathbb{T}^m$ Power[E,  $\mathbb{T}^m$ Times[ $\lambda$ , x]]]] =  $\langle \langle 1, \langle \langle "[]" \rangle, \langle n, \lambda \rangle \rangle \rangle \Big]$ 
    ParseMon[ $\mathbb{T}^m$ Times[ $\mathbb{T}^m$ Power[x, n],  $\mathbb{T}^m$ Power[E, x]]] =  $\langle \langle 1, \langle \langle "[]" \rangle, \langle n, 1 \rangle \rangle \rangle$ 
    ParseMon[ $\mathbb{T}^m$ Times[ $\mathbb{T}^m$ Times[ $\mathbb{T}^m$ Power[x, n],  $\mathbb{T}^m$ Power[E, x]]] =  $\langle \langle 1, \langle \langle "[]" \rangle, \langle n, 1 \rangle \rangle \rangle$ 
    ParseMon[ $\mathbb{T}^m$ Times[x,  $\mathbb{T}^m$ Power[E,  $\mathbb{T}^m$ Times[ $\lambda$ , x]]] =  $\langle \langle 1, \langle \langle "[]" \rangle, \langle 1, \lambda \rangle \rangle \rangle$ 
    ParseMon[ $\mathbb{T}^m$ Times[ $\mathbb{T}^m$ Times[x,  $\mathbb{T}^m$ Power[E,  $\mathbb{T}^m$ Times[ $\lambda$ , x]]] =  $\langle \langle 1, \langle \langle "[]" \rangle, \langle 1, \lambda \rangle \rangle \rangle$ 
    ParseMon[ $\mathbb{T}^m$ Power[E,  $\mathbb{T}^m$ Times[ $\lambda$ , x]]] =  $\langle \langle 1, \langle \langle "[]" \rangle, \langle 0, \lambda \rangle \rangle \rangle$ 
    ParseMon[ $\mathbb{T}^m$ Times[ $\mathbb{T}^m$ Power[E,  $\mathbb{T}^m$ Times[ $\lambda$ , x]]] =  $\langle \langle 1, \langle \langle "[]" \rangle, \langle 0, \lambda \rangle \rangle \rangle$ 
    ParseMon[ $\mathbb{T}^m$ Power[x, n]] =  $\langle \langle 1, \langle \langle "[]" \rangle, \langle n, 0 \rangle \rangle \rangle$ 
    ParseMon[ $\mathbb{T}^m$ Times[ $\mathbb{T}^m$ Power[x, n]]] =  $\langle \langle 1, \langle \langle "[]" \rangle, \langle n, 0 \rangle \rangle \rangle$ 
    ParseMon[ $\mathbb{T}^m$ Power[E, x]] =  $\langle \langle 1, \langle \langle "[]" \rangle, \langle 0, 1 \rangle \rangle \rangle$ 
    ParseMon[ $\mathbb{T}^m$ Times[ $\mathbb{T}^m$ Power[E, x]]] =  $\langle \langle 1, \langle \langle "[]" \rangle, \langle 0, 1 \rangle \rangle \rangle$ 
    ParseMon[ $\mathbb{T}^m$ Power[E,  $\mathbb{T}^m$ Minus[x]]] =  $\langle \langle 1, \langle \langle "[]" \rangle, \langle 0, -1 \rangle \rangle \rangle$ 
    ParseMon[ $\mathbb{T}^m$ Times[ $\mathbb{T}^m$ Power[E,  $\mathbb{T}^m$ Minus[x]]] =  $\langle \langle 1, \langle \langle "[]" \rangle, \langle 0, -1 \rangle \rangle \rangle$ 
    ParseMon[x] =  $\langle \langle 1, \langle \langle "[]" \rangle, \langle 1, 0 \rangle \rangle \rangle$ 
    ParseMon[ $\mathbb{T}^m$ Times[x]] =  $\langle \langle 1, \langle \langle "[]" \rangle, \langle 1, 0 \rangle \rangle \rangle$ 
    ParseMon[ $\mathbb{T}^m$ Floor[c]] =  $\langle \langle 1, \langle \langle "[]" \rangle, \langle c \rangle \rangle \rangle$ 
    ParseMon[ $\mathbb{T}^m$ Times[ $\mathbb{T}^m$ Floor[c]]] =  $\langle \langle 1, \langle \langle "[]" \rangle, \langle c \rangle \rangle \rangle$ 

```

```

ParseMonP[TD] = ⟨⟨1, ⟨"0"⟩⟩⟩
ParseMonP[TMTimes[TD]] = ⟨⟨1, ⟨"0"⟩⟩⟩
ParseMonP[TMPower[m, 0]] = ⟨⟩
ParseMonP[TMPower[m, 1]] = ParseMonP[m]
ParseMonP[TMPower[m, n]] = ParseMonP[m] *F ParseMonP[TMPower[m, n - MMAAuxF[] 1]]
ParseMonP[TMTimes[TMPower[m, n]]] = ParseMonP[TMPower[m, n]]
ParseMonP[TMTimes[TMTimes[m̄]]] = ParseMonP[TMTimes[m̄]]
ParseMonP[TMTimes[m, n, n̄]] = ParseMonP[m] *F ParseMonP[TMTimes[n, n̄]]
ParseMonP[c] = ⟨⟨c, ⟨⟩⟩⟩
ParseMonP[] = ⟨⟩

```

Definition["Parsing", any[F],

Parsing[F] = Functor[P, any[c, m, n, r, λ, m̄, n̄],

s = ⟨⟩

ParsePol<sub>P</sub>[] = ⟨⟩

ParsePol<sub>P</sub><sup>[TMTuple[m̄]]</sup> = ⟨ParsePol<sub>P</sub>[⟨m̄⟩<sub>i</sub> |<sub>i=1,...,|m̄|</sub>⟩⟩

ParsePol<sub>P</sub><sup>[TMPlus[]]</sup> = ⟨⟩

ParsePol<sub>P</sub><sup>[TMPlus[m, n̄]]</sup> = ParsePol<sub>P</sub>[m] \*<sub>P</sub> ParsePol<sub>P</sub><sup>[TMPlus[n̄]]</sup>

ParsePol<sub>P</sub><sup>[TMTimes[TMMinus[c], m̄]]</sup> = ParseTerm<sub>P</sub><sup>[TMTimes[TMMinus[c], m̄]]</sup>

ParsePol<sub>P</sub><sup>[TMTimes[c, m̄]]</sup> = ParseTerm<sub>P</sub><sup>[TMTimes[c, m̄]]</sup>

ParsePol<sub>P</sub><sup>[TMMinus[TMTimes[TMMinus[c], m̄]]]</sup> = ParseTerm<sub>P</sub><sup>[TMTimes[c, m̄]]</sup>

ParsePol<sub>P</sub><sup>[TMMinus[TMTimes[c, m̄]]]</sup> = ParseTerm<sub>P</sub><sup>[TMTimes[TMMinus[c], m̄]]</sup>

ParsePol<sub>P</sub><sup>[TMMinus[m, n]]</sup> = ParsePol<sub>P</sub>[m] \*<sub>P</sub> ParsePol<sub>P</sub><sup>[TMMinus[n]]</sup>

ParsePol<sub>P</sub><sup>[TMMinus[m, n, r, m̄]]</sup> = ParsePol<sub>P</sub>[m] \*<sub>P</sub> ParsePol<sub>P</sub><sup>[TMMinus[TMMinus[n], r, m̄]]</sup>

ParsePol<sub>P</sub><sup>[TMMinus[c]]</sup> = { ⟨⟨-c, ⟨⟩⟩⟩ ← TMIsNumber[c]  
ParseTerm<sub>P</sub><sup>[TMMinus[c]]</sup> ← otherwise

ParsePol<sub>P</sub>[c] = { ⟨⟨c, ⟨⟩⟩⟩ ← TMIsNumber[c]  
ParseTerm<sub>P</sub>[c] ← otherwise

ParseTerm<sub>P</sub><sup>[TMMinus[c]]</sup> = { ⟨⟨-c, ⟨⟩⟩⟩ ← TMIsNumber[c]  
-1 \*<sub>F</sub> ParseMon<sub>P</sub>[c] ← otherwise

ParseTerm<sub>P</sub><sup>[TMTimes[TMTimes[m̄]]]</sup> = ParseTerm<sub>P</sub><sup>[TMTimes[m̄]]</sup>

ParseTerm<sub>P</sub><sup>[TMTimes[TMMinus[c], m̄]]</sup> = { -c \*<sub>F</sub> ParseTerm<sub>P</sub><sup>[TMTimes[m̄]]</sup> ← TMIsNumber[c]  
-1 \*<sub>F</sub> ParseTerm<sub>P</sub><sup>[TMTimes[c, m̄]]</sup> ← otherwise

ParseTerm<sub>P</sub><sup>[TMTimes[c, m̄]]</sup> = { c \*<sub>F</sub> ParseTerm<sub>P</sub><sup>[TMTimes[m̄]]</sup> ← TMIsNumber[c]  
ParseMon<sub>P</sub><sup>[TMTimes[c, m̄]]</sup> ← otherwise

ParseTerm<sub>P</sub>[c] = { ⟨⟨c, ⟨⟩⟩⟩ ← TMIsNumber[c]  
ParseMon<sub>P</sub>[c] ← otherwise

ParseTerm<sub>P</sub><sup>[TMTimes[]]</sup> = ⟨⟩

ParseMon<sub>P</sub><sup>[TMTimes[TMPower[x, n], TMPower[E, TMTimes[λ, x]]]]</sup> = ⟨⟨1, ⟨n, λ⟩⟩⟩

ParseMon<sub>P</sub><sup>[TMTimes[TMTimes[TMPower[x, n], TMPower[E, TMTimes[λ, x]]]]]</sup> = ⟨⟨1, ⟨n, λ⟩⟩⟩

ParseMon<sub>P</sub><sup>[TMTimes[TMPower[x, n], TMPower[E, x]]]</sup> = ⟨⟨1, ⟨n, 1⟩⟩⟩

ParseMon<sub>P</sub><sup>[TMTimes[TMTimes[TMPower[x, n], TMPower[E, x]]]]</sup> = ⟨⟨1, ⟨n, 1⟩⟩⟩

ParseMon<sub>P</sub><sup>[TMTimes[x, TMPower[E, TMTimes[λ, x]]]]</sup> = ⟨⟨1, ⟨1, λ⟩⟩⟩

ParseMon<sub>P</sub><sup>[TMTimes[TMTimes[x, TMPower[E, TMTimes[λ, x]]]]]</sup> = ⟨⟨1, ⟨1, λ⟩⟩⟩

ParseMon<sub>P</sub><sup>[TMPower[E, TMTimes[λ, x]]]</sup> = ⟨⟨1, ⟨0, λ⟩⟩⟩

]]

```

ParseMon[Power[E, Times[x, x]]] = <<1, {0, λ}>>
ParseMon[Power[x, n]] = <<1, {n, 0}>>
ParseMon[Times[Power[x, n]]] = <<1, {n, 0}>>
ParseMon[Power[E, x]] = <<1, {0, 1}>>
ParseMon[Times[Power[E, x]]] = <<1, {0, 1}>>
ParseMon[Power[E, Minus[x]]] = <<1, {0, -1}>>
ParseMon[Times[Power[E, Minus[x]]]] = <<1, {0, -1}>>
ParseMon[x] = <<1, {1, 0}>>
ParseMon[Times[x]] = <<1, {1, 0}>>
ParseMon[Power[m, 0]] = <>
ParseMon[Power[m, 1]] = ParseMon[m]
ParseMon[Power[m, n]] = ParseMon[m] * ParseMon[Power[m, n - 1]]
ParseMon[Times[Power[m, n]]] = ParseMon[Power[m, n]]
ParseMon[Times[Times[m̄]]] = ParseMon[Times[m̄]]
ParseMon[Times[m, n, n̄]] = ParseMon[m] * ParseMon[Times[n, n̄]]
ParseMon[c] = <<c, {}>>
ParseMon[] = <>

```

Definition["ParsingF", any[ $\mathcal{F}$ ],

ParsingF[ $\mathcal{F}$ ] = Functor[P, any[c, m, n, r, λ, m̄, n̄],

s = <>

ParsePolF[] = <>

ParsePolF[0] = <>

ParsePolF[Tuple[m̄]] = {ParsePolF[<m̄><sub>i</sub>] | i=1,...,|<m̄>}

ParsePolF[Plus[]] = <>

ParsePolF[Plus[m, n̄]] = ParsePolF[m] \* ParsePolF[Plus[n̄]]

ParsePolF[Times[Minus[c], m̄]] = ParseTermF[Times[Minus[c], m̄]]

ParsePolF[Times[c, m̄]] = ParseTermF[Times[c, m̄]]

ParsePolF[Minus[Times[Minus[c], m̄]]] = ParseTermF[Times[c, m̄]]

ParsePolF[Minus[Times[c, m̄]]] = ParseTermF[Times[Minus[c], m̄]]

ParsePolF[Minus[m, n]] = ParsePolF[m] \* ParsePolF[Minus[n]]

ParsePolF[Minus[m, n, r, m̄]] = ParsePolF[m] \* ParsePolF[Minus[Minus[n], r, m̄]]

ParsePolF[Minus[c]] = { <<-c, {}>> ← IsNumber[c]  
ParseTermF[Minus[c]] ← otherwise

ParsePolF[c] = { <<c, {}>> ← IsNumber[c]  
ParseTermF[c] ← otherwise

ParseTermF[Minus[c]] = { <<-c, {}>> ← IsNumber[c]  
-1 \* ParseMonF[c] ← otherwise

ParseTermF[Times[Times[m̄]]] = ParseTermF[Times[m̄]]

ParseTermF[Times[Minus[c], m̄]] = { -c \* ParseTermF[Times[m̄]] ← IsNumber[c]  
-1 \* ParseTermF[Times[c, m̄]] ← otherwise

ParseTermF[Times[c, m̄]] = { c \* ParseTermF[Times[m̄]] ← IsNumber[c]  
ParseMonF[Times[c, m̄]] ← otherwise

$$\text{ParseTermF}[c] = \begin{cases} \langle \langle c, \langle \rangle \rangle & \Leftarrow \text{IsNumber}[c] \\ \text{ParseMonF}[c] & \Leftarrow \text{otherwise} \end{cases}$$

$$\text{ParseTermF}[\text{Times}[]] = \langle \rangle$$

$$\text{ParseMonF}[\text{Times}[\text{Power}[x, m], \text{Power}[E, \text{Times}[\lambda, x]]]] = \langle \langle 1, \langle m, \lambda \rangle \rangle \rangle$$

$$\text{ParseMonF}[\text{Times}[\text{Power}[E, \text{Times}[\lambda, x]], \text{Power}[x, m]]] = \langle \langle 1, \langle m, \lambda \rangle \rangle \rangle$$

$$\text{ParseMonF}[\text{Times}[\text{Power}[x, m], \text{Power}[E, x]]] = \langle \langle 1, \langle m, 1 \rangle \rangle \rangle$$

$$\text{ParseMonF}[\text{Times}[x, \text{Power}[E, \text{Times}[\lambda, x]]]] = \langle \langle 1, \langle 1, \lambda \rangle \rangle \rangle$$

$$\text{ParseMonF}[\text{Times}[\text{Times}[\text{Power}[x, m], \text{Power}[E, \text{Times}[\lambda, x]]]]] = \langle \langle 1, \langle m, \lambda \rangle \rangle \rangle$$

$$\text{ParseMonF}[\text{Times}[\text{Times}[\text{Power}[E, \text{Times}[\lambda, x]], \text{Power}[x, m]]]] = \langle \langle 1, \langle m, \lambda \rangle \rangle \rangle$$

$$\text{ParseMonF}[\text{Times}[\text{Times}[\text{Power}[x, m], \text{Power}[E, x]]]] = \langle \langle 1, \langle m, 1 \rangle \rangle \rangle$$

$$\text{ParseMonF}[\text{Times}[\text{Times}[x, \text{Power}[E, \text{Times}[\lambda, x]]]]] = \langle \langle 1, \langle 1, \lambda \rangle \rangle \rangle$$

$$\text{ParseMonF}[\text{Power}[x, m]] = \langle \langle 1, \langle m, 0 \rangle \rangle \rangle$$

$$\text{ParseMonF}[\text{Times}[\text{Power}[x, m]]] = \langle \langle 1, \langle m, 0 \rangle \rangle \rangle$$

$$\text{ParseMonF}[\text{Power}[E, \text{Times}[\lambda, x]]] = \langle \langle 1, \langle 0, \lambda \rangle \rangle \rangle$$

$$\text{ParseMonF}[\text{Times}[\text{Power}[E, \text{Times}[\lambda, x]]]] = \langle \langle 1, \langle 0, \lambda \rangle \rangle \rangle$$

$$\text{ParseMonF}[x] = \langle \langle 1, \langle 1, 0 \rangle \rangle \rangle$$

$$\text{ParseMonF}[\text{Times}[x]] = \langle \langle 1, \langle 1, 0 \rangle \rangle \rangle$$

$$\text{ParseMonF}[\text{Power}[E, x]] = \langle \langle 1, \langle 0, 1 \rangle \rangle \rangle$$

$$\text{ParseMonF}[\text{Times}[\text{Power}[E, x]]] = \langle \langle 1, \langle 0, 1 \rangle \rangle \rangle$$

$$\text{ParseMonF}[] = \langle \rangle$$

Definition["Parsing Diff Polys", any[ $\mathcal{F}$ ],

ParsingDP[ $\mathcal{F}$ ] = Functor[P, any[c, k, l, m, n, r,  $\lambda$ ,  $\bar{c}$ ,  $\bar{d}$ ,  $\bar{m}$ ,  $\bar{n}$ ],

s =  $\langle \rangle$

ParsePolF[] =  $\langle \rangle$

ParsePolF[0] =  $\langle \rangle$

ParsePolF[Tuple[ $\bar{m}$ ]] =  $\left( \text{ParsePolF}[\langle \bar{m} \rangle_i] \mid_{i=1, \dots, |\langle \bar{m} \rangle|} \right)$

ParsePolF[Plus[]] =  $\langle \rangle$

ParsePolF[Plus[m,  $\bar{n}$ ]] = ParsePolF[m]  $\times$  ParsePolF[Plus[ $\bar{n}$ ]]

ParsePolF[Times[Minus[c],  $\bar{m}$ ]] = ParseTermF[Times[Minus[c],  $\bar{m}$ ]]

ParsePolF[Times[c,  $\bar{m}$ ]] = ParseTermF[Times[c,  $\bar{m}$ ]]

ParsePolF[Minus[Times[Minus[c],  $\bar{m}$ ]]] = ParseTermF[Times[c,  $\bar{m}$ ]]

ParsePolF[Minus[Times[c,  $\bar{m}$ ]]] = ParseTermF[Times[Minus[c],  $\bar{m}$ ]]

ParsePolF[Minus[m, n]] = ParsePolF[m]  $\times$  ParsePolF[Minus[n]]

ParsePolF[Minus[m, n, r,  $\bar{m}$ ]] = ParsePolF[m]  $\times$  ParsePolF[Minus[Minus[n], r,  $\bar{m}$ ]]

ParsePolF[Minus[c]] =  $\begin{cases} \langle \langle -c, \langle 0, 0 \rangle \rangle & \Leftarrow \text{IsNumber}[c] \\ \text{ParseTermF}[Minus[c]] & \Leftarrow \text{otherwise} \end{cases}$

ParsePolF[c] =  $\begin{cases} \langle \langle c, \langle 0, 0 \rangle \rangle & \Leftarrow \text{IsNumber}[c] \\ \text{ParseTermF}[c] & \Leftarrow \text{otherwise} \end{cases}$

ParseTermF[Minus[c]] =  $\begin{cases} \langle \langle -c, \langle 0, 0 \rangle \rangle & \Leftarrow \text{IsNumber}[c] \\ -1_{\mathcal{F}} \cdot \text{ParseMonF}[c] & \Leftarrow \text{otherwise} \end{cases}$

ParseTermF[Times[Times[ $\bar{m}$ ]]] = ParseTermF[Times[ $\bar{m}$ ]]

$$\text{ParseTermF}_{\text{P}}[\text{Times}[\text{Minus}[c], \bar{m}]] = \begin{cases} -c \cdot \text{ParseTermF}_{\text{P}}[\text{Times}[\bar{m}]] & \Leftarrow \text{IsNumber}[c] \\ -1 \cdot \text{ParseTermF}_{\text{P}}[\text{Times}[c, \bar{m}]] & \Leftarrow \text{otherwise} \end{cases}$$

$$\text{ParseTermF}_{\text{P}}[\text{Times}[c, \bar{m}]] = \begin{cases} c \cdot \text{ParseTermF}_{\text{P}}[\text{Times}[\bar{m}]] & \Leftarrow \text{IsNumber}[c] \\ \text{ParseMonF}_{\text{P}}[\text{Times}[c, \bar{m}]] & \Leftarrow \text{otherwise} \end{cases}$$

$$\text{ParseTermF}_{\text{P}}[c] = \begin{cases} \langle \langle c, \langle 0, 0 \rangle \rangle & \Leftarrow \text{IsNumber}[c] \\ \text{ParseMonF}_{\text{P}}[c] & \Leftarrow \text{otherwise} \end{cases}$$

$$\text{ParseTermF}_{\text{P}}[\text{Times}[]] = \langle \rangle$$

$$\text{ParseMonF}_{\text{P}}[\text{Times}[\text{Power}[x, m], \text{Power}[E, \text{Times}[\lambda, x]]]] = \langle \langle 1, \langle m, \lambda \rangle \rangle \rangle$$

$$\text{ParseMonF}_{\text{P}}[\text{Times}[\text{Power}[E, \text{Times}[\lambda, x]], \text{Power}[x, m]]] = \langle \langle 1, \langle m, \lambda \rangle \rangle \rangle$$

$$\text{ParseMonF}_{\text{P}}[\text{Times}[\text{Power}[x, m], \text{Power}[E, x]]] = \langle \langle 1, \langle m, 1 \rangle \rangle \rangle$$

$$\text{ParseMonF}_{\text{P}}[\text{Times}[x, \text{Power}[E, x]]] = \langle \langle 1, \langle 1, 1 \rangle \rangle \rangle$$

$$\text{ParseMonF}_{\text{P}}[\text{Times}[x, \text{Power}[E, \text{Times}[\lambda, x]]]] = \langle \langle 1, \langle 1, \lambda \rangle \rangle \rangle$$

$$\text{ParseMonF}_{\text{P}}[\text{Times}[\text{Times}[\text{Power}[x, m], \text{Power}[E, \text{Times}[\lambda, x]]]]] = \langle \langle 1, \langle m, \lambda \rangle \rangle \rangle$$

$$\text{ParseMonF}_{\text{P}}[\text{Times}[\text{Times}[\text{Power}[E, \text{Times}[\lambda, x]], \text{Power}[x, m]]]] = \langle \langle 1, \langle m, \lambda \rangle \rangle \rangle$$

$$\text{ParseMonF}_{\text{P}}[\text{Times}[\text{Times}[\text{Power}[x, m], \text{Power}[E, x]]]] = \langle \langle 1, \langle m, 1 \rangle \rangle \rangle$$

$$\text{ParseMonF}_{\text{P}}[\text{Times}[\text{Times}[x, \text{Power}[E, \text{Times}[\lambda, x]]]]] = \langle \langle 1, \langle 1, \lambda \rangle \rangle \rangle$$

$$\text{ParseMonF}_{\text{P}}[\text{Power}[x, m]] = \langle \langle 1, \langle m, 0 \rangle \rangle \rangle$$

$$\text{ParseMonF}_{\text{P}}[\text{Times}[\text{Power}[x, m]]] = \langle \langle 1, \langle m, 0 \rangle \rangle \rangle$$

$$\text{ParseMonF}_{\text{P}}[\text{Times}[c, \text{Power}[x, m]]] = \langle \langle c, \langle m, 0 \rangle \rangle \rangle$$

$$\text{ParseMonF}_{\text{P}}[\text{Power}[E, \text{Times}[\lambda, x]]] = \langle \langle 1, \langle 0, \lambda \rangle \rangle \rangle$$

$$\text{ParseMonF}_{\text{P}}[\text{Times}[\text{Power}[E, \text{Times}[\lambda, x]]]] = \langle \langle 1, \langle 0, \lambda \rangle \rangle \rangle$$

$$\text{ParseMonF}_{\text{P}}[x] = \langle \langle 1, \langle 1, 0 \rangle \rangle \rangle$$

$$\text{ParseMonF}_{\text{P}}[E] = \langle \langle 1, \langle 0, 1 \rangle \rangle \rangle$$

$$\text{ParseMonF}_{\text{P}}[\text{Times}[x]] = \langle \langle 1, \langle 1, 0 \rangle \rangle \rangle$$

$$\text{ParseMonF}_{\text{P}}[\text{Power}[E, x]] = \langle \langle 1, \langle 0, 1 \rangle \rangle \rangle$$

$$\text{ParseMonF}_{\text{P}}[\text{Times}[\text{Power}[E, x]]] = \langle \langle 1, \langle 0, 1 \rangle \rangle \rangle$$

$$\text{ParseMonF}_{\text{P}}[\text{Times}[c]] = \text{ParsePolF}_{\text{P}}[c]$$

$$\text{ParseMonF}_{\text{P}}[c] = \langle \langle c, \langle 0, 0 \rangle \rangle \rangle$$

$$\text{ParseMonF}_{\text{P}}[] = \langle \rangle$$

$$\text{ParseDP}_{\text{P}}[] = \langle \rangle$$

$$\text{ParseDP}_{\text{P}}[\text{Plus}[]] = \langle \rangle$$

$$\text{ParseDP}_{\text{P}}[\text{Plus}[m, \bar{n}]] = \text{ParseMonDP}_{\text{P}}[m] \times \text{ParseDP}_{\text{P}}[\text{Plus}[\bar{n}]]$$

$$\text{ParseDP}_{\text{P}}[\text{Power}[u, k]] = \left\langle \left\langle \frac{1}{\bar{r}}, k \right\rangle \right\rangle$$

$$\text{ParseDP}_{\text{P}}[\text{Times}[\bar{c}, \text{Power}[u, k]]] = \text{ParseMonDP}_{\text{P}}[\text{Times}[\bar{c}, \text{Power}[u, k]]]$$

$$\text{ParseMonDP}_{\text{P}}[\text{Times}[\bar{c}, \text{Power}[u, k]]] = \text{where} \left[ p = \text{ParsePolF}_{\text{P}}[\text{Times}[\bar{c}]], \right. \\ \left. \langle \langle p, k \rangle \rangle \right]$$

$$\text{ParseMonDP}_{\text{P}}[\text{Power}[u, k]] = \left\langle \left\langle \frac{1}{\bar{r}}, k \right\rangle \right\rangle$$

$$\text{ParseIDP}_{\text{P}}[] = \langle \rangle$$

$$\text{ParseIDP}_{\text{P}}[\text{Plus}[]] = \langle \rangle$$

$$\text{ParseIDP}_{\text{P}}[\text{Plus}[m, \bar{n}]] = \text{ParseIDP}_{\text{P}}[m] \times \text{ParseIDP}_{\text{P}}[\text{Plus}[\bar{n}]]$$

$$\text{ParseIDP}_{\text{P}}[\text{Times}[\text{Minus}[c], \bar{m}]] = \text{ParseMonIDP}_{\text{P}}[\text{Times}[\text{Minus}[c], \bar{m}]]$$

$$\text{ParseIDP}_{\text{P}}[\text{Minus}[\text{Times}[\text{Minus}[c], \bar{m}]]] = \text{ParseMonIDP}_{\text{P}}[\text{Times}[c, \bar{m}]]$$

$$\text{ParseIDP}_{\text{P}}[\text{Minus}[\text{Times}[c, \bar{m}]]] = \text{ParseMonIDP}_{\text{P}}[\text{Times}[\text{Minus}[c], \bar{m}]]$$

$$\begin{aligned}
\text{ParseIDP}[\text{TMMinus}[m, n]] &= \text{ParseIDP}[m] \times \text{ParseIDP}[\text{TMMinus}[n]] \\
\text{ParseIDP}[\text{TMMinus}[m, n, r, \bar{m}]] &= \text{ParseIDP}[m] \times \text{ParseIDP}[\text{TMMinus}[\text{TMMinus}[n], r, \bar{m}]] \\
\text{ParseIDP}[\text{TMMinus}[c]] &= \begin{cases} \left\langle \left\langle \text{ParsePolF}[-c], \langle 0, 0 \rangle \right\rangle \right\rangle \Leftarrow \in_{\mathcal{F}}[c] \\ \text{ParseMonIDP}[\text{TMMinus}[c]] \Leftarrow \text{otherwise} \end{cases} \\
\text{ParseMonIDP}[\text{TMMinus}[c]] &= \begin{cases} \left\langle \left\langle \text{ParsePolF}[-c], \langle 0, 0 \rangle \right\rangle \right\rangle \Leftarrow \in_{\mathcal{F}}[c] \\ \left( \begin{array}{c} 1 \\ \text{IntDiffPolys}[\mathcal{F}, K] \end{array} \right) \cdot \text{IntDiffPolys}[\mathcal{F}, K] \text{ParseMonIDP}[c] \Leftarrow \text{otherwise} \end{cases} \\
\text{ParseIDP}[\text{TMTimes}[c, \bar{d}]] &= \text{ParseMonIDP}[\text{TMTimes}[c, \bar{d}]] \\
\text{ParseIDP}[\text{TMPower}[u, k]] &= \left\langle \left\langle \frac{1}{\mathcal{F}}, \langle \langle \rangle, k, \langle \rangle \rangle \right\rangle \right\rangle \\
\text{ParseIDP}[u] &= \left\langle \left\langle \frac{1}{\mathcal{F}}, \langle \langle \rangle, \langle 1 \rangle, \langle \rangle \rangle \right\rangle \right\rangle \\
\text{ParseIDP}[\text{TMPower}[u[0], k]] &= \left\langle \left\langle \frac{1}{\mathcal{F}}, \langle k, \langle \rangle, \langle \rangle \rangle \right\rangle \right\rangle \\
\text{ParseIDP}[u[0]] &= \left\langle \left\langle \frac{1}{\mathcal{F}}, \langle \langle 1 \rangle, \langle \rangle, \langle \rangle \rangle \right\rangle \right\rangle \\
\text{ParseIDP}[\text{TMTimes}[c, \text{TMPower}[u, k]]] &= \text{ParseMonIDP}[\text{TMTimes}[c, \text{TMPower}[u, k]]] \\
\text{ParseIDP}[c] &= \left\langle \left\langle \text{ParsePolF}[c], \langle \langle \rangle, \langle \rangle, \langle \rangle \rangle \right\rangle \right\rangle \\
\text{ParseMonIDP}[\text{TMMinus}[c]] &= \begin{cases} \left\langle \left\langle \text{ParsePolF}[-c], \langle 0, 0 \rangle \right\rangle \right\rangle \Leftarrow \in_{\mathcal{F}}[c] \\ \left( \begin{array}{c} - \\ \text{IntDiffPolys}[\mathcal{F}, K] \end{array} \right) \text{ParseMonIDP}[c] \Leftarrow \text{otherwise} \end{cases} \\
\text{ParseMonIDP}[\text{TMTimes}[\text{TMMinus}[c], \bar{m}]] &= \begin{cases} \text{ParsePolF}[-c] \cdot \text{IntDiffPolys}[\mathcal{F}, K] \text{ParseMonIDP}[\text{TMTimes}[\bar{m}]] \Leftarrow \in_{\mathcal{F}}[c] \\ \left( \begin{array}{c} - \\ \text{IntDiffPolys}[\mathcal{F}, K] \end{array} \right) \text{ParseMonIDP}[\text{TMTimes}[c, \bar{m}]] \Leftarrow \text{otherwise} \end{cases} \\
\text{ParseMonIDP}[\text{TMPower}[u, k]] &= \left\langle \left\langle \frac{1}{\mathcal{F}}, k \right\rangle \right\rangle \\
\text{ParseMonIDP}[\text{TMTimes}[c, \text{TMPower}[u[0], k], \text{TMPower}[u, 1], \bar{d}]] &= \text{where } [p = \text{ParsePolF}[c], \\ & \left\langle \left\langle p, \langle k, 1, \text{ParseFunc}[\text{TMTimes}[\bar{d}]] \rangle \right\rangle \right\rangle \\
\text{ParseMonIDP}[\text{TMTimes}[\text{TMPower}[u[0], k], \text{TMPower}[u, 1], \bar{d}]] &= \left\langle \left\langle \frac{1}{\mathcal{F}}, \langle k, 1, \text{ParseFunc}[\text{TMTimes}[\bar{d}]] \rangle \right\rangle \right\rangle \\
\text{ParseMonIDP}[\text{TMTimes}[c, \text{TMPower}[u[0], k], "f", \bar{d}]] &= \text{where } [p = \text{ParsePolF}[c], \\ & \left\langle \left\langle p, \langle k, \langle \rangle, \text{ParseFunc}[\text{TMTimes}["f", \bar{d}]] \rangle \right\rangle \right\rangle \\
\text{ParseMonIDP}[\text{TMTimes}[\text{TMPower}[u[0], k], "f", \bar{d}]] &= \left\langle \left\langle \frac{1}{\mathcal{F}}, \langle k, \langle \rangle, \text{ParseFunc}[\text{TMTimes}["f", \bar{d}]] \rangle \right\rangle \right\rangle \\
\text{ParseMonIDP}[\text{TMTimes}[c, \text{TMPower}[u, 1], "f", \bar{d}]] &= \text{where } [p = \text{ParsePolF}[c], \\ & \left\langle \left\langle p, \langle \langle \rangle, 1, \text{ParseFunc}[\text{TMTimes}["f", \bar{d}]] \rangle \right\rangle \right\rangle \\
\text{ParseMonIDP}[\text{TMTimes}[\text{TMPower}[u, 1], "f", \bar{d}]] &= \left\langle \left\langle \frac{1}{\mathcal{F}}, \langle \langle \rangle, 1, \text{ParseFunc}[\text{TMTimes}["f", \bar{d}]] \rangle \right\rangle \right\rangle \\
\text{ParseMonIDP}[\text{TMTimes}[c, \text{TMPower}[u[0], k]]] &= \text{where } [p = \text{ParsePolF}[c], \\ & \left\langle \left\langle p, \langle k, \langle \rangle, \langle \rangle \rangle \right\rangle \right\rangle \\
\text{ParseMonIDP}[\text{TMTimes}[\text{TMPower}[u[0], k], \bar{c}]] &= \left\langle \left\langle \frac{1}{\mathcal{F}}, \langle k, \langle \rangle, \text{ParseFunc}[\text{TMTimes}[\bar{c}]] \rangle \right\rangle \right\rangle \\
\text{ParseMonIDP}[\text{TMTimes}[c, u[0]]] &= \text{where } [p = \text{ParsePolF}[c], \\ & \left\langle \left\langle p, \langle \langle 1 \rangle, \langle \rangle, \langle \rangle \rangle \right\rangle \right\rangle \\
\text{ParseMonIDP}[u[0]] &= \left\langle \left\langle \frac{1}{\mathcal{F}}, \langle \langle 1 \rangle, \langle \rangle, \langle \rangle \rangle \right\rangle \right\rangle \\
\text{ParseMonIDP}[\text{TMTimes}[c, \text{TMPower}[u, k]]] &= \text{where } [p = \text{ParsePolF}[c], \\ & \left\langle \left\langle p, \langle \langle \rangle, k, \langle \rangle \rangle \right\rangle \right\rangle
\end{aligned}$$

Definition["ParsingIDO", any $[\mathcal{F}]$ ,

ParsingIDO $[\mathcal{F}] = \text{Functor}[P, \text{any}[c, m, n, r, \lambda, \bar{m}, \bar{n}]$ ,

```

s = ⟨⟩
ParsePolP[] = ⟨⟩
ParsePolP[TMTuple[m̄]] = ⟨ParsePolP[⟨m̄i⟩i=1,...,|⟨m̄⟩|⟩⟩
ParsePolP[TMPlus[]] = ⟨⟩
ParsePolP[TMPlus[m, n̄]] = ParsePolP[m] ⊗ ParsePolP[TMPlus[n̄]]
ParsePolP[TMTimes[TMMinus[c], m̄]] = ParseTermP[TMTimes[TMMinus[c], m̄]]
ParsePolP[TMTimes[c, m̄]] = ParseTermP[TMTimes[c, m̄]]
ParsePolP[TMMinus[TMTimes[TMMinus[c], m̄]]] = ParseTermP[TMTimes[c, m̄]]
ParsePolP[TMMinus[TMTimes[c, m̄]]] = ParseTermP[TMTimes[TMMinus[c], m̄]]
ParsePolP[TMMinus[m, n]] = ParsePolP[m] ⊗ ParsePolP[TMMinus[n]]
ParsePolP[TMMinus[m, n, r, m̄]] = ParsePolP[m] ⊗ ParsePolP[TMMinus[TMMinus[n], r, m̄]]
ParsePolP[TMMinus[c]] = {⟨⟨-c, ⟨⟩⟩⟩ ⊆ TMIsNumber[c]
ParseTermP[TMMinus[c]] ⊆ otherwise
ParsePolP[c] = {⟨⟨c, ⟨⟩⟩⟩ ⊆ TMIsNumber[c]
ParseTermP[c] ⊆ otherwise
ParseTermP[TMMinus[c]] = {⟨⟨-c, ⟨⟩⟩⟩ ⊆ TMIsNumber[c]
ParseTermP[TMMinus[c]] = {-1 f ParseMonP[c] ⊆ otherwise
ParseTermP[TMTimes[TMTimes[m̄]]] = ParseTermP[TMTimes[m̄]]
ParseTermP[TMTimes[TMMinus[c], m̄]] = { -c f ParseTermP[TMTimes[m̄]] ⊆ TMIsNumber[c]
-1 f ParseTermP[TMTimes[c, m̄]] ⊆ otherwise
ParseTermP[TMTimes[c, m̄]] = { c f ParseTermP[TMTimes[m̄]] ⊆ TMIsNumber[c]
ParseMonP[TMTimes[c, m̄]] ⊆ otherwise
ParseTermP[c] = {⟨⟨c, ⟨⟩⟩⟩ ⊆ TMIsNumber[c]
ParseMonP[c] ⊆ otherwise
ParseTermP[TMTimes[]] = ⟨⟩
ParseMonP[TMTimes[TMPower[u, n], TMPower[v, λ]]] = ⟨⟨1, ⟨⟨"[]", ⟨⟨⟨⟨1 f, ⟨⟩, n, ⟨⟩⟩⟩⟩, ⟨⟩, ⟨⟩, ⟨⟩⟩⟩⟩, ⟨"[]",
ParseMonP[TMTimes[TMTimes[TMPower[u, n], TMPower[v, λ]]]] = ⟨⟨1, ⟨⟨"[]", ⟨⟨⟨⟨1 f, ⟨⟩, n, ⟨⟩⟩⟩⟩, ⟨⟩, ⟨⟩, ⟨⟩⟩⟩⟩
ParseMonP[TMPower[u, n]] = ⟨⟨1, ⟨⟨"[]", ⟨⟨⟨⟨1 f, ⟨⟩, n, ⟨⟩⟩⟩⟩, ⟨⟩, ⟨⟩, ⟨⟩⟩⟩⟩
ParseMonP[TMPower[u[0], λ]] = ⟨⟨1, ⟨⟨"[]", ⟨⟨⟨⟨1 f, ⟨λ, ⟨⟩, ⟨⟩⟩⟩⟩, ⟨⟩, ⟨⟩, ⟨⟩⟩⟩⟩
ParseMonP[TMFloor[c]] = ⟨⟨1, ⟨⟨"[]", c⟩⟩⟩
ParseMonP[TMTimes[TMFloor[c]]] = ⟨⟨1, ⟨⟨"[]", c⟩⟩⟩
ParseMonP[TME] = ⟨⟨1, ⟨⟨"[]", 0⟩⟩⟩
ParseMonP[TMTimes[TME]] = ⟨⟨1, ⟨⟨"[]", 0⟩⟩⟩
ParseMonP[TM∂] = ⟨⟨1, ⟨"∂"⟩⟩⟩
ParseMonP[TMTimes[TM∂]] = ⟨⟨1, ⟨"∂"⟩⟩⟩
ParseMonP[TM∫] = ⟨⟨1, ⟨"∫"⟩⟩⟩
ParseMonP[TMTimes[TM∫]] = ⟨⟨1, ⟨"∫"⟩⟩⟩
ParseMonP[TMPower[m, 0]] = ⟨⟩
ParseMonP[TMPower[m, 1]] = ParseMonP[m]
ParseMonP[TMPower[m, n]] = ParseMonP[m] * ParseMonP[TMPower[m, nMmaAuxF[] - 1]]
ParseMonP[TMTimes[TMPower[m, n]]] = ParseMonP[TMPower[m, n]]
ParseMonP[TMTimes[TMTimes[m̄]]] = ParseMonP[TMTimes[m̄]]
ParseMonP[TMTimes[m, n, n̄]] = ParseMonP[m] * ParseMonP[TMTimes[n, n̄]]
ParseMonP[c] = ⟨⟨c, ⟨⟩⟩⟩

```

```

ParseMon[] = <>

```

```

Definition["ParsingRules", any[ $\mathcal{F}$ ],
ParsingRules[ $\mathcal{F}$ ] = Functor[P, any[c, m, n, r,  $\lambda$ ,  $\bar{m}$ ,  $\bar{n}$ ],

s = <>
ParsePol[ $\mathbb{T}^{\text{Plus}}$ [ $\mathbb{T}^{\text{Times}}$ [ $\mathbb{T}^{\text{Minus}}$ [ $\mathbb{T}^{\text{Power}}$ [u,  $\mathbb{T}^{\text{Tuple}}$ [1]],  $\mathbb{T}^{\text{Power}}$ [v,  $\mathbb{T}^{\text{Tuple}}$ [1]],  $\mathbb{T}^{\text{CenterDot}}$ [ $\mathbb{T}^{\text{Power}}$ [u,  $\mathbb{T}^{\text{Tuple}}$ [1]]],
ParsePol[ $\mathbb{T}^{\text{Plus}}$ [ $\mathbb{T}^{\text{Minus}}$ ["E"],  $\mathbb{T}^{\text{Times}}$ ["E", "E"]] = <<1, <<"["], 0>, <<"["], 0>>>, <-1, <<"["], 0>>>
ParsePol[ $\mathbb{T}^{\text{Minus}}$ [ $\mathbb{T}^{\text{Times}}$ ["E",  $\mathbb{T}^{\text{Power}}$ [u,  $\mathbb{T}^{\text{Tuple}}$ [1]],  $\mathbb{T}^{\text{Times}}$ [ $\mathbb{T}^{\text{Power}}$ [u[0],  $\mathbb{T}^{\text{Tuple}}$ [1]], "E"]] = <<1, <<"["]'
ParsePol[ $\mathbb{T}^{\text{Minus}}$ [ $\mathbb{T}^{\text{Times}}$ [" $\partial$ ",  $\mathbb{T}^{\text{Power}}$ [u,  $\mathbb{T}^{\text{Tuple}}$ [1]],  $\mathbb{T}^{\text{Power}}$ [u,  $\mathbb{T}^{\text{Tuple}}$ [0, 1]],  $\mathbb{T}^{\text{Times}}$ [ $\mathbb{T}^{\text{Power}}$ [u,  $\mathbb{T}^{\text{Tuple}}$ [1]]],
ParsePol[ $\mathbb{T}^{\text{Times}}$ [" $\partial$ ", "E"]] = <<1, <<" $\partial$ ", <<"["], 0>>>
ParsePol[ $\mathbb{T}^{\text{Plus}}$ [ $\mathbb{T}^{\text{Minus}}$ [1],  $\mathbb{T}^{\text{Times}}$ [" $\partial$ ", " $\int$ "]] = <<1, <<" $\partial$ ", " $\int$ ">>, <-1, <>>
ParsePol[ $\mathbb{T}^{\text{Minus}}$ [ $\mathbb{T}^{\text{Plus}}$ [ $\mathbb{T}^{\text{Times}}$ [" $\int$ ",  $\mathbb{T}^{\text{Power}}$ [u,  $\mathbb{T}^{\text{Tuple}}$ [1]], " $\int$ "],  $\mathbb{T}^{\text{Times}}$ [" $\int$ ", " $\int$ ",  $\mathbb{T}^{\text{Power}}$ [u,  $\mathbb{T}^{\text{Tuple}}$ [1]]],
ParsePol[ $\mathbb{T}^{\text{Plus}}$ [ $\mathbb{T}^{\text{Minus}}$ [ $\mathbb{T}^{\text{Plus}}$ [ $\mathbb{T}^{\text{Times}}$ [" $\int$ ",  $\mathbb{T}^{\text{Power}}$ [u,  $\mathbb{T}^{\text{Tuple}}$ [0, 1]],  $\mathbb{T}^{\text{Times}}$ [" $\int$ ",  $\mathbb{T}^{\text{Power}}$ [u,  $\mathbb{T}^{\text{Tuple}}$ [1]], " $\partial$ "],
ParsePol[ $\mathbb{T}^{\text{Minus}}$ [ $\mathbb{T}^{\text{Times}}$ [" $\int$ ",  $\mathbb{T}^{\text{Power}}$ [u,  $\mathbb{T}^{\text{Tuple}}$ [1]], "E"],  $\mathbb{T}^{\text{Times}}$ [" $\int$ ",  $\mathbb{T}^{\text{Power}}$ [u,  $\mathbb{T}^{\text{Tuple}}$ [1]], "E"]] = <<1,
ParseRules[ $\langle \bar{m} \rangle$ ] = {ParsePol[ $\langle \bar{m} \rangle_i$ ] |  $i=1, \dots, |\langle \bar{m} \rangle|$ }
ParsePol[ $\mathbb{T}^{\text{Minus}}$ [ $\mathbb{T}^{\text{Minus}}$ [ $\mathbb{T}^{\text{Times}}$ [ $\mathbb{T}^{\text{CenterDot}}$ [" $\int$ ",  $\mathbb{T}^{\text{Power}}$ [u,  $\mathbb{T}^{\text{Tuple}}$ [1]],  $\mathbb{T}^{\text{CenterDot}}$ [" $\int$ ",  $\mathbb{T}^{\text{Power}}$ [v,  $\mathbb{T}^{\text{Tuple}}$ [1]]],
ParsePol[ $\mathbb{T}^{\text{Plus}}$ [ $\mathbb{T}^{\text{Minus}}$ [ $\mathbb{T}^{\text{Times}}$ [" $\int$ ",  $\mathbb{T}^{\text{Power}}$ [v,  $\mathbb{T}^{\text{Tuple}}$ [1]], " $\int$ "],  $\mathbb{T}^{\text{Times}}$ [ $\mathbb{T}^{\text{CenterDot}}$ [" $\int$ ",  $\mathbb{T}^{\text{Power}}$ [v,  $\mathbb{T}^{\text{Tuple}}$ [1]]],
ParsePol[ $\mathbb{T}^{\text{Plus}}$ [ $\mathbb{T}^{\text{Minus}}$ [ $\mathbb{T}^{\text{Times}}$ [" $\int$ ",  $\mathbb{T}^{\text{Power}}$ [u,  $\mathbb{T}^{\text{Tuple}}$ [1]], " $\int$ "],  $\mathbb{T}^{\text{Times}}$ [ $\mathbb{T}^{\text{CenterDot}}$ [" $\int$ ",  $\mathbb{T}^{\text{Power}}$ [u,  $\mathbb{T}^{\text{Tuple}}$ [1]]],
ParsePol[ $\mathbb{T}^{\text{Plus}}$ [ $\mathbb{T}^{\text{Minus}}$ [ $\mathbb{T}^{\text{Times}}$ [" $\int$ ",  $\mathbb{T}^{\text{Power}}$ [u,  $\mathbb{T}^{\text{Tuple}}$ [1]], " $\partial$ "],  $\mathbb{T}^{\text{Power}}$ [u,  $\mathbb{T}^{\text{Tuple}}$ [1]],  $\mathbb{T}^{\text{Plus}}$ [ $\mathbb{T}^{\text{Times}}$ [" $\int$ ",
ParsePol[ $\mathbb{T}^{\text{Plus}}$ [ $\mathbb{T}^{\text{Minus}}$ [ $\mathbb{T}^{\text{Plus}}$ [ $\mathbb{T}^{\text{Times}}$ [" $\int$ ",  $\mathbb{T}^{\text{Power}}$ [u,  $\mathbb{T}^{\text{Tuple}}$ [1]], " $\partial$ "],  $\mathbb{T}^{\text{Times}}$ [" $\int$ ",  $\mathbb{T}^{\text{Power}}$ [u,  $\mathbb{T}^{\text{Tuple}}$ [0, 1]]],
]]

```

```

ParseP

```

```

AddFreshNameKeywords[{ "ParseP", "ParseF", "ParseDP", "ParseDPMon", "ParseFunc", "ParseIDPMon",
"ParseIDP", "ParseDP", "Parse", "ComputeP", "FormatP", "GreensOp", "Mult", "Fact", "MultF",
"GreensFct", "Act", "AddIDP", "MultIDP", "MinIDP", "DiffIDP", "IntIDP", "ExpBasis", "ParseIDO",
"ParseRules", "ParseRule", "SpolRed", "Spol", "Spol1", "ReducePol", "RedPol", "GB"}];

```

```

Clear[ParseP, ParseF, ParseDP, ParseDPMon, Parse, ComputeP, FormatP,
GreensOp, Mult, Fact, MultF, GreensFct, Act, ParseFunc, ParseMonIDP, ParseIDP];

```

## Theory IDP

```
Definition["Integro-Differential Polynomials",
          Exp = ExpPolys[K]
          K = ReductionField[C]
          F = DiffPolys[Exp, K]
          I = IntDiffOp[Exp, K]
          Gr = GreenSystem[Exp, K]
          IDP = IntDiffPolys[Exp, K]
          a = Parsing[I]
          P = GB[FreeIntDiffOp[IntDiffPolys[IntDiffPolys[Exp, K], K], K]]
          p = PrettyPrintingIDP[F]
          gs = GreenSystem[IntDiffPolys[IntDiffPolys[Exp, K], K], K]
          G = GB[FreeIntDiffOp[IntDiffPolys[IntDiffPolys[Exp, K], K], K]]
        ]
```

```

Theory["Integro-Differential Polynomials",
  Definition["Natural Numbers"]
    Definition["N"]
  Definition["Rational Number Field"]
    Definition["Q"]
  Definition["Reduction Field"]
  Definition["Complex Number Field"]
    Definition["C"]
  Definition["Complex Numbers"]
    Definition["CM"]
  Definition["Tuples Monoid"]
  Definition["Cartesian Product"]
  Definition["Cartesian Power"]
    Definition["Free Module"]
  Definition["Monoid Algebra"]
    Definition["Exp Polys"]
    Definition["Mma Functions"]
  Definition["Differential Polynomials"]
  Definition["Integro-Differential Polynomials"]
    Definition["Term Monoid for IDP"]
  Definition["Integro-Differential Polys"]
    Definition["Finite Chain"]
    Definition["Word Monoid"]
  Definition["Word Monoid with Degree Ordering"]
    Definition["Groebner Extension"]
  Definition["Theorema general functions"]
    Definition["Basis"]
    Definition["Double Basis"]
  Definition["PrettyPrinting"]
    Definition["↦"]
  Definition["Green System"]
  Definition["Groebner Basis Normal Form"]
    Definition["FreeIntDiffOp"]
  Definition["Basis for Characters"]
  Definition["Quotient Algebra"]
    Definition["IntDiffOp"]
  Definition["Parsing Exp"]
    Definition["ParsingF"]
    Definition["ParsingIDO"]
  Definition["Parsing Diff Polys"]
    Definition["ParsingRules"]
  Definition["Integro-Differential Polynomials"]
]

```

```

Built-in["Integro-Differential Polynomials",
  Built-in["Numbers"]
  Built-in["Quantifiers"]
  Built-in["Tuples"]
  Built-in["Connectives"]
  Built-in["Sets"]
  Built-in["Operators"]
  Built-in["IsComplex"]
  Built-in["ComplexLt"]
  Built-in["Mult"]
  Built-in["Minus"]
  Built-in["IsNumber"]
  Built-in["MmaSimplify"]
  Built-in["FuncOrder"]
  Built-in["FuncQ"]
  Built-in["MmaDiff"]
  Built-in["MmaInt"]
  Built-in["ToMma"]
  Built-in["MmaPlus"]
  Built-in["MmaMinus"]
  Built-in["MmaTimes"]
  Built-in["MmaSubs"]
  Built-in["TmaTRFM"]
  Built-in["TmaSplit"]
]

```

```
Use[Theory["Integro-Differential Polynomials"]]
```

```

ParseP[p___] := Compute[ParsePol[p],
  built-in → {Built-in["Tuples"], Built-in["Quantifiers"], Built-in["Connectives"],
    Built-in["Mult"], Built-in["Minus"], Built-in["IsNumber"], Built-in["MmaSimplify"]}];

```

```

ParseF[f___] := Compute[ParsePolF[f],
  ParsingDP[Exp]
  built-in → {Built-in["Tuples"], Built-in["Quantifiers"], Built-in["Connectives"],
    Built-in["Mult"], Built-in["Minus"], Built-in["IsNumber"], Built-in["MmaSimplify"]}];

```

```

ParseDP[f___] := Compute[ParseDP[f],
  ParsingDP[Exp]
  built-in → {Built-in["Tuples"], Built-in["Quantifiers"], Built-in["Connectives"],
    Built-in["Mult"], Built-in["Minus"], Built-in["IsNumber"], Built-in["MmaSimplify"]}];

```

```

ParseDPMon[f___] := Compute[ParseMonDP[f],
  ParsingDP[Exp]
  built-in → {Built-in["Tuples"], Built-in["Quantifiers"], Built-in["Connectives"],
    Built-in["Mult"], Built-in["Minus"], Built-in["IsNumber"], Built-in["MmaSimplify"]}];

```

```

ParseFunc[f___] := Compute[ParseFunc[f],
  ParsingDP[Exp]
  built-in → {Built-in["Tuples"], Built-in["Quantifiers"], Built-in["Connectives"],
    Built-in["Mult"], Built-in["Minus"], Built-in["IsNumber"], Built-in["MmaSimplify"]}];

```

```
ParseIDPMon[f___] := Compute[ParseMonIDP[f],
  ParsingDP[Exp],
  built-in → <Built-in["Tuples"], Built-in["Quantifiers"], Built-in["Connectives"],
    Built-in["Mult"], Built-in["Minus"], Built-in["IsNumber"], Built-in["MmaSimplify"]>>;
```

```
ParseIDP[f___] := Compute[ParseIDP[f],
  ParsingDP[Exp],
  built-in → <Built-in["Tuples"], Built-in["Quantifiers"], Built-in["Connectives"],
    Built-in["Mult"], Built-in["Minus"], Built-in["IsNumber"], Built-in["MmaSimplify"]>>;
```

```
ParseIDO[f___] := Compute[ParsePol[f],
  ParsingIDO[Exp],
  built-in → <Built-in["Tuples"], Built-in["Quantifiers"], Built-in["Connectives"],
    Built-in["Mult"], Built-in["Minus"], Built-in["IsNumber"], Built-in["MmaSimplify"]>>;
```

```
ParseRules[f___] := Compute[ParseRules[f],
  ParsingRules[Exp],
  built-in → <Built-in["Tuples"], Built-in["Quantifiers"], Built-in["Connectives"],
    Built-in["Mult"], Built-in["Minus"], Built-in["IsNumber"], Built-in["MmaSimplify"]>>;
```

```
ParseRule[f___] := Compute[ParsePol[f],
  ParsingRules[Exp],
  built-in → <Built-in["Tuples"], Built-in["Quantifiers"], Built-in["Connectives"],
    Built-in["Mult"], Built-in["Minus"], Built-in["IsNumber"], Built-in["MmaSimplify"]>>;
```

```
AddIDP[p_, q_] := ComputeP[ParseIDP[p] + ParseIDP[q]] // PrintIDP;
```

```
MinIDP[p_, q_] := ComputeP[ParseIDP[p] IntDiffPolys[Exp,K] ParseIDP[q]] // PrintIDP;
```

```
MultIDP[p_, q_] := ComputeP[ParseIDP[p] * ParseIDP[q]] // PrintIDP;
```

```
DiffIDP[p_] := ComputeP[∂IntDiffPolys[Exp,K] ParseIDP[p]] // PrintIDP;
```

```
IntIDP[p_] := ComputeP[∫IntDiffPolys[Exp,K] ParseIDP[p]] // PrintIDP;
```

```
ExpBasis[p_] := ComputeP[exp-basisIDP[ParseIDP[p]]] // PrintIDP;
```

```
AddFreshNameKeywords[{"PrintDP", "PrintIDP", "PrintP", "Print", "PrintRules"}];
```

```
PrintDP[x_] := Compute[PrettyPrint[x], built-in → Built-in["Integro-Differential Polynomials"]];
```

```
PrintIDP[p_] := Compute[PrettyPrintIntDiffPolys[Exp,K][p], built-in → Built-in["Integro-Differential Polynomials"]];
```

```
PrintP[p_] := Compute[
  PrettyPrint
  PrettyPrinting[FreeIntDiffOp[IntDiffPolys[IntDiffPolys[Exp,K],Exp],K],K]
  [p],
  built-in → Built-in["Integro-Differential Polynomials"]];
```

```
PrintRules[r_] := Compute[
  PrettyPrintRules
  PrettyPrinting[FreeIntDiffOp[IntDiffPolys[IntDiffPolys[Exp,K],Exp],K],K]
  [r],
  built-in → Built-in["Integro-Differential Polynomials"]];
```

```
ComputeP[p_] := Compute[p, built-in → Built-in["Integro-Differential Polynomials"]];
```

```
Parse[p_] := ParseP[p /. mytimes → mTimes];
```

```
Mult[p_, q_] := PrintIDP[ComputeP[Parse[p] * Parse[q]]];
```

```
Act[g_, f_] := FormatP[ComputeP[Parse[g] ⊙F Parse[f]]] /.
  {mTimes → Times, mPlus → Plus, mMinus → Minus, mPower → Power};
```

```
SpolRed[f_, g_] := PrintP[spolredG[ParseRules[f], g]];
```

```
SpolRed[f_] := PrintP[spolredG[ParseRules[f],  $\mathcal{G}$ ]];
```

```
Spol[f_, g_] := PrintP[spolG[ParseRule[f], ParseRule[g]]];
```

```
ReducePol[f_] := PrintP[tredG[ParseRule[f],  $\mathcal{G}$ ]];
```

```
GB[s_] := PrintP[GbG[ParseRules[s],  $\mathcal{G}$ ]];
```

#### Example Computations IDP

```
Use[Theory["Integro-Differential Polynomials"]]
```

```
MultIDP[u[0]^(1,2) ∫ x u^(1,1) ∫ (3 x^2) u^(0,2), 3 u[0]^(2,3) u^(3,1) ∫ (4 x) u^(1,0,1)]
```

```
3 u[0]^(3,3) u^(5,1) ∫ x u^(1,1) ∫ 4 x u^(1,0,1) ∫ 3 x^2 u^(0,2) +
3 u[0]^(3,3) u^(5,1) ∫ x u^(1,1) ∫ 3 x^2 u^(0,2) ∫ 4 x u^(1,0,1) + 3 u[0]^(3,3) u^(5,1) ∫ 4 x u^(1,0,1) ∫ x u^(1,1) ∫ 3 x^2 u^(0,2)
```

$$\text{DiffIDP}\left[3 u[0]^{(1,2,3)} u^{(4,5)} \int u^{(5,6)} - (2 x^3 E^{4x} + x^3 - x) u^{(5,6)} - (3 E^x + x) u[0]^{(1,2,3)} u^{(4,3)} \int u^{(5,6)} - (2 x^3 E^{4x} - x^3) u^{(5,6)} + (3 E^x + x) u[0]^{(1,2,3)} u^{(4,2)} \int u^{(5,5)} - (3 E^x) u[0]^{(1,2,3)} u^{(4,1)} \int u^{(5,2)}\right]$$

$$\begin{aligned} & -12 E^x u[0]^{(1,2,3)} u^{(3,2)} \int -x^3 - 2 x^3 E^{4x} u^{(5,2)} + (4 x + 12 E^x) u[0]^{(1,2,3)} u^{(3,3)} \int -x^3 - 2 x^3 E^{4x} u^{(5,5)} + \\ & (-4 x - 12 E^x) u[0]^{(1,2,3)} u^{(3,4)} \int -x^3 - 2 x^3 E^{4x} u^{(5,6)} + 12 u[0]^{(1,2,3)} u^{(3,6)} \int -x + x^3 - 2 x^3 E^{4x} u^{(5,6)} - \\ & 3 E^x u[0]^{(1,2,3)} u^{(4,1)} \int -x^3 - 2 x^3 E^{4x} u^{(5,2)} + (1 + 3 E^x) u[0]^{(1,2,3)} u^{(4,2)} \int -x^3 - 2 x^3 E^{4x} u^{(5,5)} + \\ & (-1 - 3 E^x) u[0]^{(1,2,3)} u^{(4,3)} \int -x^3 - 2 x^3 E^{4x} u^{(5,6)} + (3 x^3 E^x + 6 x^3 E^{5x}) u[0]^{(1,2,3)} u^{(9,3)} + \\ & (-x^4 - 3 x^3 E^x - 2 x^4 E^{4x} - 6 x^3 E^{5x}) u[0]^{(1,2,3)} u^{(9,7)} + (x^4 + 3 x^3 E^x + 2 x^4 E^{4x} + 6 x^3 E^{5x}) u[0]^{(1,2,3)} u^{(9,9)} + \\ & (-3 x + 3 x^3 - 6 x^3 E^{4x}) u[0]^{(1,2,3)} u^{(9,11)} - 3 E^x u[0]^{(1,2,3)} u^{(4,0,1)} \int -x^3 - 2 x^3 E^{4x} u^{(5,2)} + \\ & (2 x + 6 E^x) u[0]^{(1,2,3)} u^{(4,1,1)} \int -x^3 - 2 x^3 E^{4x} u^{(5,5)} + \\ & (-3 x - 9 E^x) u[0]^{(1,2,3)} u^{(4,2,1)} \int -x^3 - 2 x^3 E^{4x} u^{(5,6)} + \\ & 15 u[0]^{(1,2,3)} u^{(4,4,1)} \int -x + x^3 - 2 x^3 E^{4x} u^{(5,6)} \end{aligned}$$

$$\text{IntIDP}\left[3 u[0]^{(1,2,3)} u^{(4,5,6,1)} \int u^{(5,6)} - (2 x^3 E^{4x} + x^3 - x) u^{(5,6)} - (3 E^x + x) u[0]^{(1,2,3)} u^{(4,3)} \int u^{(5,6)} - (2 x^3 E^{4x} - x^3) u^{(5,6)} + (3 E^x + x) u[0]^{(1,2,3)} u^{(4,2)} \int u^{(5,5)} - (3 E^x) u[0]^{(1,2,3)} u^{(4,1)} \int u^{(5,2)}\right]$$

$$\begin{aligned} & -\frac{1}{7} u[0]^{(1,2,3)} \int 12 u^{(3,6,7)} \int -x + x^3 - 2 x^3 E^{4x} u^{(5,6)} - \frac{1}{7} u[0]^{(1,2,3)} \int 15 u^{(4,4,8)} \int -x + x^3 - 2 x^3 E^{4x} u^{(5,6)} - \\ & \frac{1}{5} u[0]^{(1,2,3)} \int -3 E^x u^{(5)} \int -x^3 - 2 x^3 E^{4x} u^{(5,2)} + u[0]^{(1,2,3)} \int -x - 3 E^x u^{(4,3)} \int -x^3 - 2 x^3 E^{4x} u^{(5,6)} + \\ & u[0]^{(1,2,3)} \int x + 3 E^x u^{(4,2)} \int -x^3 - 2 x^3 E^{4x} u^{(5,5)} - \frac{1}{7} u[0]^{(1,2,3)} \int -3 x + 3 x^3 - 6 x^3 E^{4x} u^{(9,11,7)} - \\ & \frac{1}{5} u[0]^{(1,2,3)} \int 3 x^3 E^x + 6 x^3 E^{5x} u^{(10,2)} - \frac{3}{5} E^x u[0]^{(1,2,3)} u^{(5)} \int -x^3 - 2 x^3 E^{4x} u^{(5,2)} + \\ & \frac{3}{7} u[0]^{(1,2,3)} u^{(4,5,7)} \int -x + x^3 - 2 x^3 E^{4x} u^{(5,6)} \end{aligned}$$

$$\text{ExpBasis}\left[3 u[0]^{(1,2,3)} u^{(4,5)} \int u^{(5,6)} - (2 x^3 E^{4x} + x^3 - x) u^{(5,6)} - (3 E^x + x) u[0]^{(1,2,3)} u^{(4,3)} \int u^{(5,6)} - (2 x^3 E^{4x} - x^3) u^{(5,6)} + (3 E^x + x) u[0]^{(1,2,3)} u^{(4,2)} \int u^{(5,5)} - (3 E^x) u[0]^{(1,2,3)} u^{(4,1)} \int u^{(5,2)}\right]$$

$$\begin{aligned} & 3 E^x u[0]^{(1,2,3)} u^{(4,1)} \int x^3 u^{(5,2)} + 6 E^x u[0]^{(1,2,3)} u^{(4,1)} \int x^3 E^{4x} u^{(5,2)} + \\ & (-x - 3 E^x) u[0]^{(1,2,3)} u^{(4,2)} \int x^3 u^{(5,5)} + (-2 x - 6 E^x) u[0]^{(1,2,3)} u^{(4,2)} \int x^3 E^{4x} u^{(5,5)} + \\ & (x + 3 E^x) u[0]^{(1,2,3)} u^{(4,3)} \int x^3 u^{(5,6)} + (2 x + 6 E^x) u[0]^{(1,2,3)} u^{(4,3)} \int x^3 E^{4x} u^{(5,6)} - 3 u[0]^{(1,2,3)} u^{(4,5)} \int x u^{(5,6)} + \\ & 3 u[0]^{(1,2,3)} u^{(4,5)} \int x^3 u^{(5,6)} - 6 u[0]^{(1,2,3)} u^{(4,5)} \int x^3 E^{4x} u^{(5,6)} \end{aligned}$$