Automated Reasoning Systems Resolution Theorem Proving: Prover9

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Prover9

- Automated theorem prover from the Argonne group
- Successor of Otter
- Author: Bill McCune
- Theory: First-order logic with equality
- Inference rules: Based on resolution and paramodulation

- Main Applications: In abstract algebra and formal logic
- Implemented in C.
- Web page: http://www.cs.unm.edu/~mccune/prover9/

Running Prover9

- 1. Prepare the input file containing the logical specification of a conjecture and the search parameters.
- 2. Issue a command that runs Prover9 on the input file and produces an output file.
- 3. Look at the output.
- 4. maybe run Prover9 again with different search parameters.

How to Run Prover9

From the command line:

- > Run: prover9 -f inputfile or prover9 -f inputfile > outputfile or prover9 < inputfile or prover9 < inputfile > outputfile
- > There can be several input files, e.g., prover9 -f file1 ... filen > outputfile
- Time limit can be specified from the command line: prover9 -t 10 -f inputfile > outputfile
 Using GUI.

- Ordinary symbols: made from the characters a-z, A-Z, 0-9, \$, and _.
- Special symbols: made from the special characters: {+-*/\^<>= `~?@& | ! #';.
- Quoted symbols: any string enclosed in double quotes.
- Special characters can not be mixed with the others to make symbols, unless quoted symbols are formed.

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Meaning	Connective	Example
negation	-	(-p)
disjunction	I	(p q r)
conjunction	&	(p & q & r)
implication	->	(p -> q)
backward implic.	<-	(p <- q)
equivalence	<->	(p <-> q)
universal quant.	all	(all x all y p(x,y))
existential quant.	exists	(exists x exists y p(x,y))
true	\$T	
false	\$F	
equality	=	(a = b)
negated equality	! =	(a != b)

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- Symbols are assigned precedences.
- Infix, prefix, and postfix declarations are used.
- All that helps to avoid excessive use of parentheses.
- Details:

http://www.cs.unm.edu/~mccune/prover9/manual/2009-11A/syntax.html#declarations

For instance, terms involving +, * and binary – can be written in infix notation: a+b, a*b, a-b.

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- Prover9 input file: A sequence of lists and commands.
- Lists (of formulas or clauses) can be used to declare, for instance, which formulas or clauses are assumptions and which ones are goals.
- Commands can be used, for instance, to set certain options.
- The order is largely irrelevant, except for some special cases.

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Example

assign(max_seconds, 30).

% Is Socrates mortal?

```
formulas(assumptions).
    all x (man(x) -> mortal(x)).
    man(socrates).
end of list.
```

```
formulas(goals).
    mortal(socrates).
end_of_list.
```

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Example

```
assign(max_seconds, 30).
formulas(usable).
   all x (man(x) -> mortal(x)).
   man(socrates).
end_of_list.
```

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```
formulas(sos).
    -mortal(socrates).
end_of_list.
```

Example

```
assign(max_seconds, 30).
formulas(sos).
   all x (man(x) -> mortal(x)).
   man(socrates).
   -mortal(socrates).
end_of_list.
```

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Example

assign(max_seconds, 30).

% Is Socrates mortal?

```
clauses(assumptions).
   - man(x) | mortal(x).
   man(socrates).
end of list.
```

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```
formulas(goals).
    mortal(socrates).
end_of_list.
```

Distinction Between Clauses and Formulas

- Formulas can have any of the logic connectives, and all variables are explicitly quantified. Formulas go into lists that start with formulas (list_name).
- Clauses are simple disjunctions in which all variables are implicitly universally quantified. Clauses go into lists that start with clauses (list_name).
- Clauses without variables are also formulas, so they can go into either kind of list. (An exception: clauses can have attributes, and formulas cannot.)
- Because variables in clauses are not explicitly quantified, a rule is needed for distinguishing variables from constants in clauses (see terms below). No such rule is needed for formulas.
- Prover9's inference rules operate on clauses. If formulas are input, Prover9 immediately translates them into clauses.

Mail Loop: Given Clause Algorithm

Operates on the sos and usable lists. While the sos list is not empty:

- 1. Select a given clause from sos and move it to the usable list;
- Infer new clauses using the inference rules in effect; each new clause must have the given clause as one of its parents and members of the usable list as its other parents;
- 3. Process each new clause;
- 4. Append new clauses that pass the retention tests to the sos list.

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end of while loop.

Example

Barbers paradox:

- There is a town with just one male barber.
- Every man in the town keeps himself clean-shaven.
- ► Some shave themselves, some are shaved by the barber.
- The barber shaves all and only those men in town who do not shave themselves.

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Does the barber shave himself?

Example

Barbers paradox:

- There is a town with just one male barber.
- ► Every man in the town keeps himself clean-shaven.
- Some shave themselves, some are shaved by the barber.
- The barber shaves all and only those men in town who do not shave themselves.
- Does the barber shave himself?

```
formulas(goals).
    exists x (barber(x) &
        (all y (-shaves(y,y) <-> shaves(x,y)))).
end_of_list.
```

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Run Prover9 on Examples Example

Assumption : $\forall x \forall y \ Q(x) \Rightarrow P(y, y).$ Assumption : $\forall x \forall y \ P(x, b) \lor Q(f(y, x)).$ Goal : $\exists x \ P(a, x).$

Example

Assumption : $\forall x \forall y \forall z \ (x * y) * z = x * (y * z).$ Assumption : $\forall x x * 1 = x.$ Assumption : $\forall x x * x^{-1} = 1.$ Assumption : $\forall x x * x = 1.$ Goal : $\forall x \forall y x * y = y * x.$