Formal Semantics of Programming Languages Exercise 2 (December 1)

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The exercise is to be submitted by **December 1** (hard deadline)

- 1. as a single PDF file sent to me per email, or
- 2. as a paper report (cover page with full name and Matrikelnummer, pages stapled) handed out to me in class.

1 Expressions with Side effects

Take the following language of commands C, expressions E, numerals N, and identifiers I:

$$\begin{array}{l} C ::= I := E \mid C_1; C_2 \mid \textbf{if} \ (E_1 = E_2) \ C \\ E ::= I \mid N \mid E_1 + E_2 \mid \textbf{exec} \ C \ \textbf{result} \ E \end{array}$$

The **exec** expression executes C and then returns the result of the evaluation of E. Correspondingly, the evaluation of an expression may alter the store.

Define an operational semantics for this language with configurations of the form $\langle C, s \rangle \to s'$ (command C evaluated in store s yields store s') and $\langle E, s \rangle \to \langle n, s' \rangle$ (expression E evaluated in store s yields number n and store s').

2 A New Language Construct

A language designer proposes a new looping construct **entangle** with the following property:

$$[$$
entangle (B) $C]$ $]$ $=$ $[$ if B $\{C;$ entangle (B) $C;$ $C\}]$

- 1. Show the sequence of "unfoldings" of the construct (0-3 unfoldings).
- 2. Define a denotational semantics for the construct.
- 3. Prove that your semantics satisfies the property stated above.
- 4. Sketch a possible implementation of the construct on a computer.